Real-Time Systems

Hard Real-Time Multiprocessor Scheduling

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Outline

- Introduction
- Terminology, Notation and Assumptions
- Anomalies + Impossibility Results
- Partitioned Scheduling (no migration)
- Global Scheduling (task- / job-level migration)
    - G-FP
    - G-EDF
- Optimal MP Scheduling
- MP – Resource Access Protocols
- Open Research Issues

Lessons Learned From UP

- **Liu-Layland Criterion for fixed task priority algorithms:**
  - can schedule any workload up to a utilization of $U_{RMS}(n) = n\sqrt[n]{2} - 1 \leq 0.693$
  - scheduling of workloads with higher utilization not guaranteed
- **fixed job priority algorithms are optimal: (e.g., EDF)**
- **there are optimal greedy algorithms**
  - with a single measure characterizing the "importance" of a job (e.g., time to deadline, laxity, …)
- **all preemptive FTP, FJP algorithms are predictable**
  - response times cannot increase when decreasing execution times
- **all preemptive FTP algorithms and EDF are sustainable**
  - no period / deadline anomalies
- **simultaneous release is critical instant**
- **response times depend on <u>set</u> but not on order of high-priority tasks**

Taxonomy of Multiprocessor Scheduling

Two problems to solve:

- Priority Problem: When to run a given job of the workload?
  - fixed task priority (e.g., RMS, …)
  - fixed job priority (e.g., EDF, …)
  - dynamic job priority (e.g., LST, …)

- Allocation Problem: Where to run this job?
  - no migration — partitioned
  - task-level migration (no migration of running jobs) — global
  - job-level migration (migrate also running jobs) — global

Taxonomy of Multiprocessor Scheduling

Two problems to solve:

- Priority Problem: When to run a given job of the workload?
  - fixed task p
  - fixed job pri
  - dynamic job

Preemption Costs – UP / MP

- Allocation Problem: Where to run this job?
  - no migration
  - task-level m
  - job-level mi

partitioned

global

Migration Costs – MP

Preemption Costs

- direct costs:
  - timer / device interrupt
  - save register state
  - manipulate ready list
    - UP: no synchronization required
  - load register state of next thread
- indirect costs
  - cache evictions between two consecutive runs
  - TLB refills after evictions / shootdown

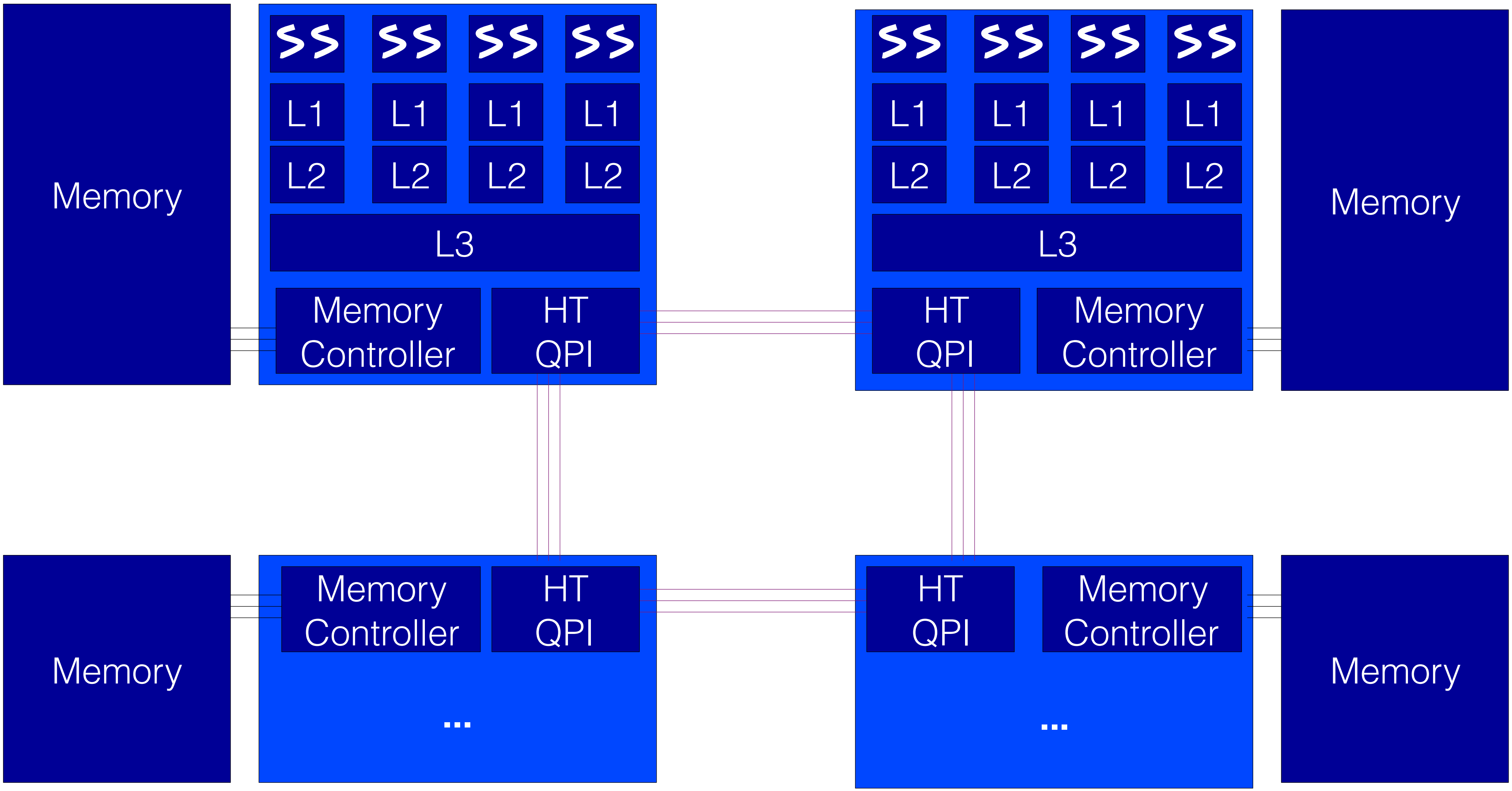
Migration Costs

- job-level migration
  - migration of running job implies preemption at source CPU
- task-level migration
  - job is already preempted

- direct costs
  - manipulate remote / global ready list
    - synchronization
  - fetch register state
- indirect costs
  - fetch active cache working set from remote cache
  - load remaining data from remote memory

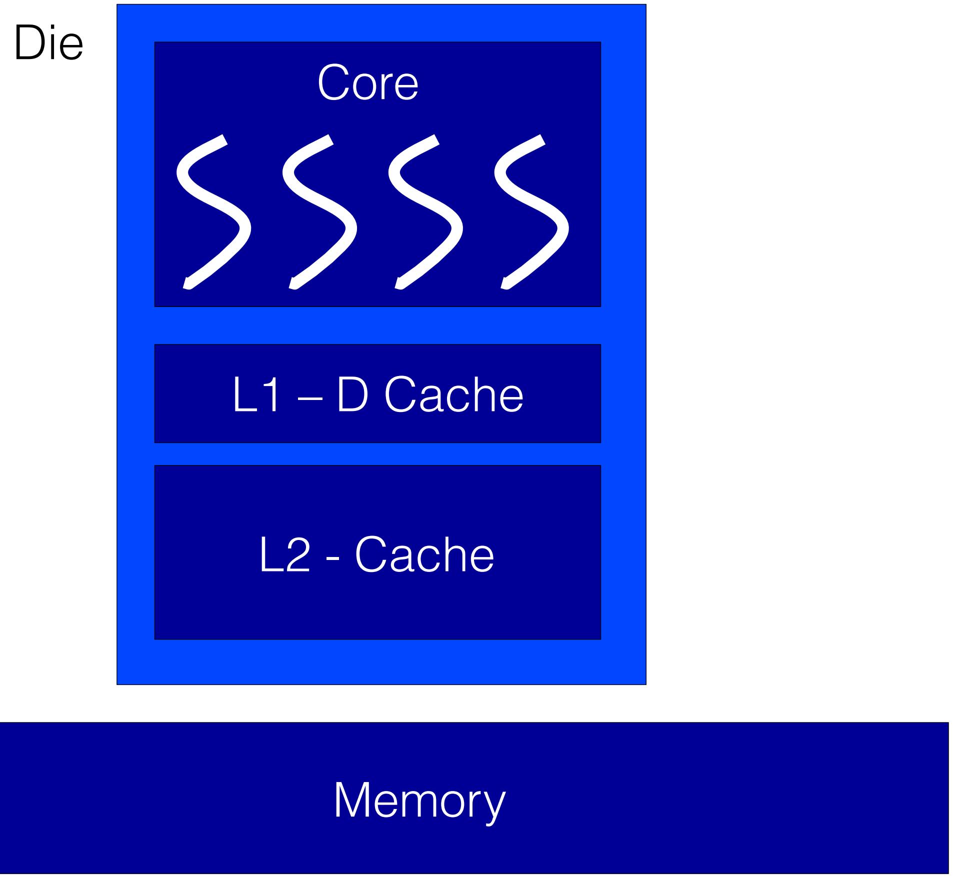
Multiprocessor Architectures

- AMD Opteron / Intel Core Duo: SMT + multi core + ccNuma



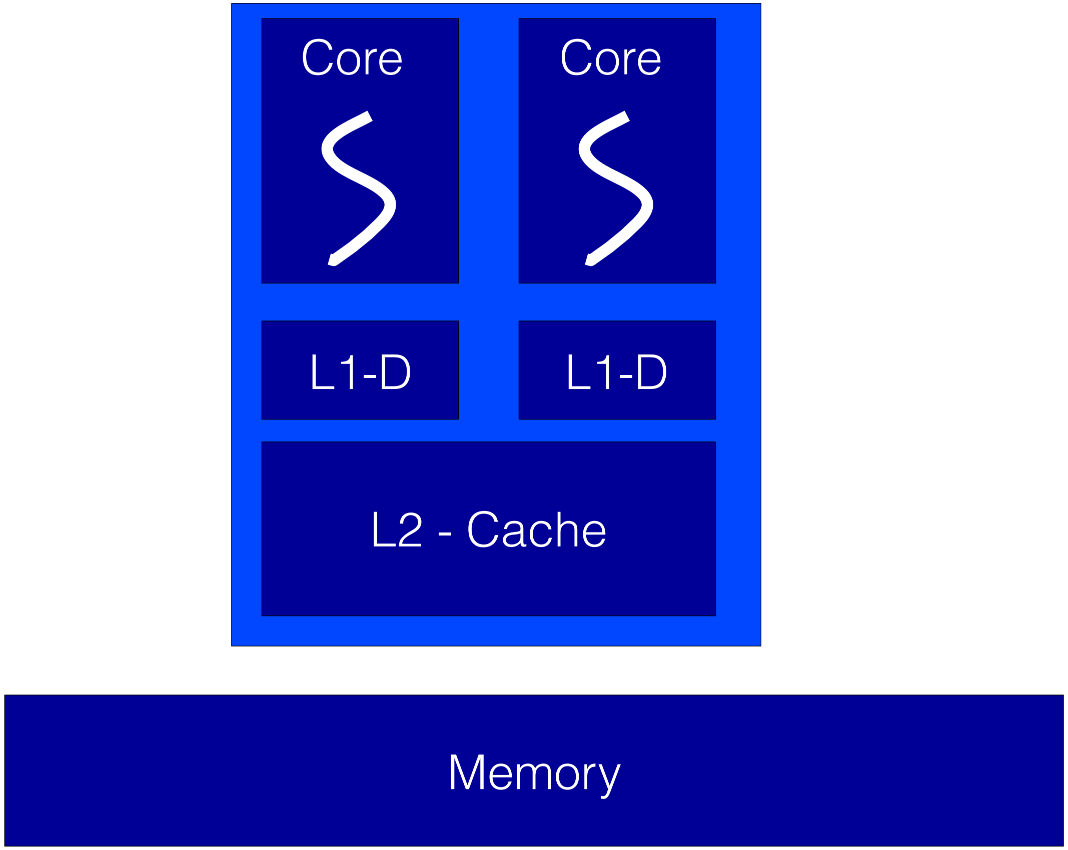
Multiprocessor Architectures

- Symmetric Multi-Threaded (SMT) Processors
  - operating system multiplexes n SW threads on m HW threads
  - caches + pipeline is shared => no indirect migration costs



Multiprocessor Architectures

- Multi-Core Processors
  - operating system multiplexes n SW threads on m cores
  - timing of last level cache dominates migration costs



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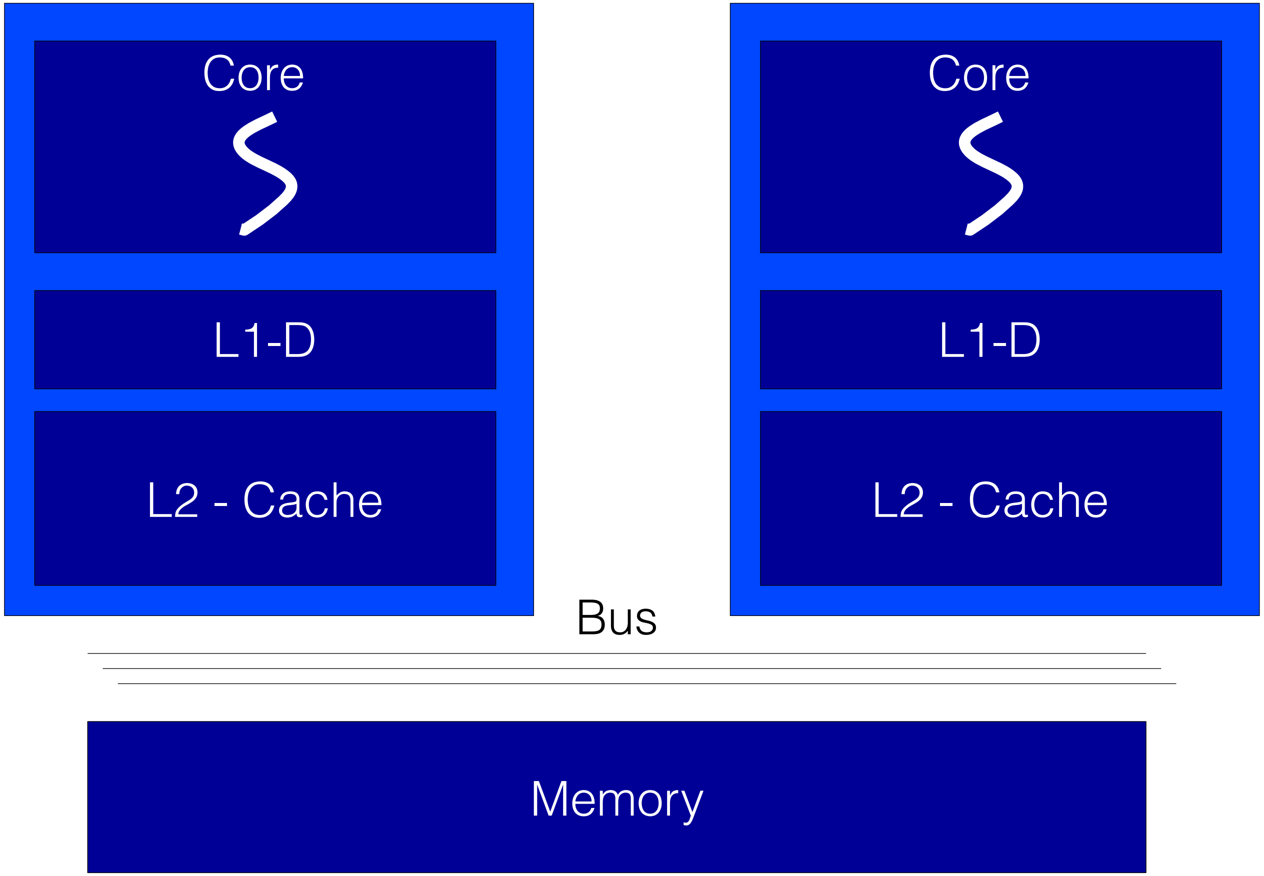
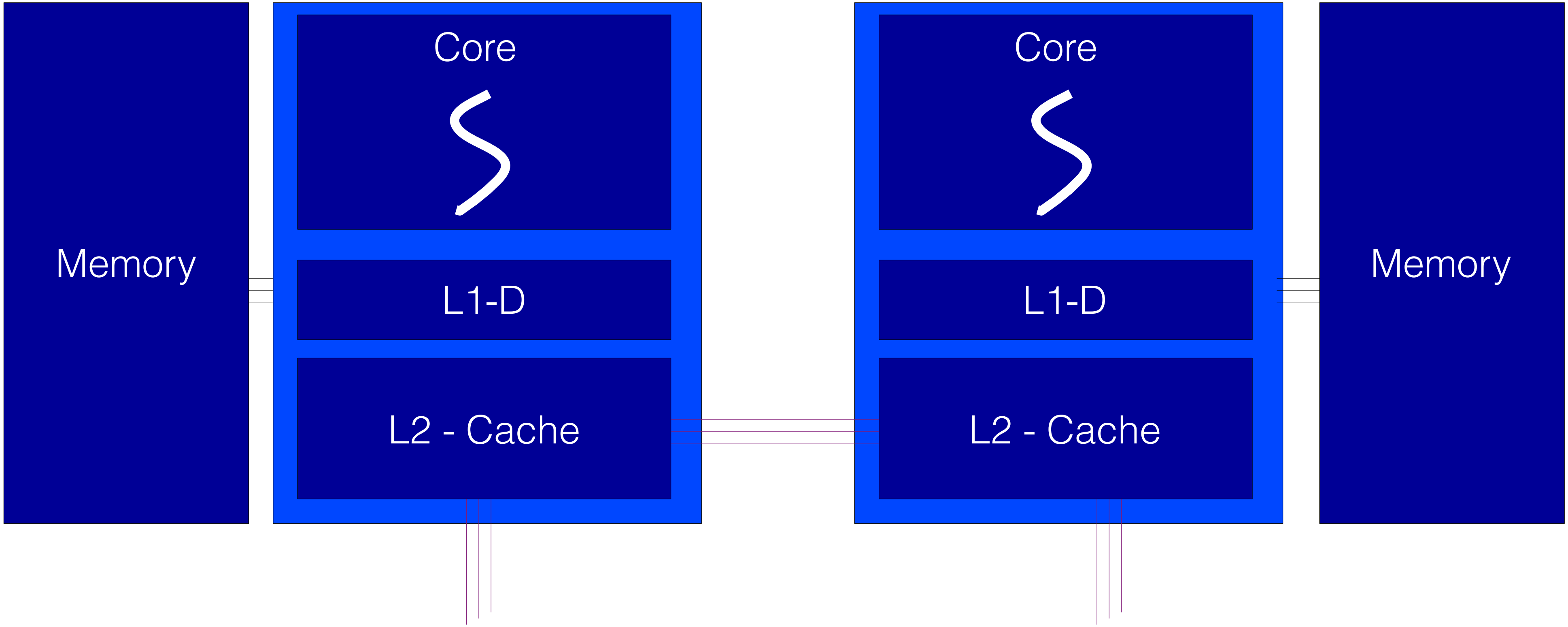


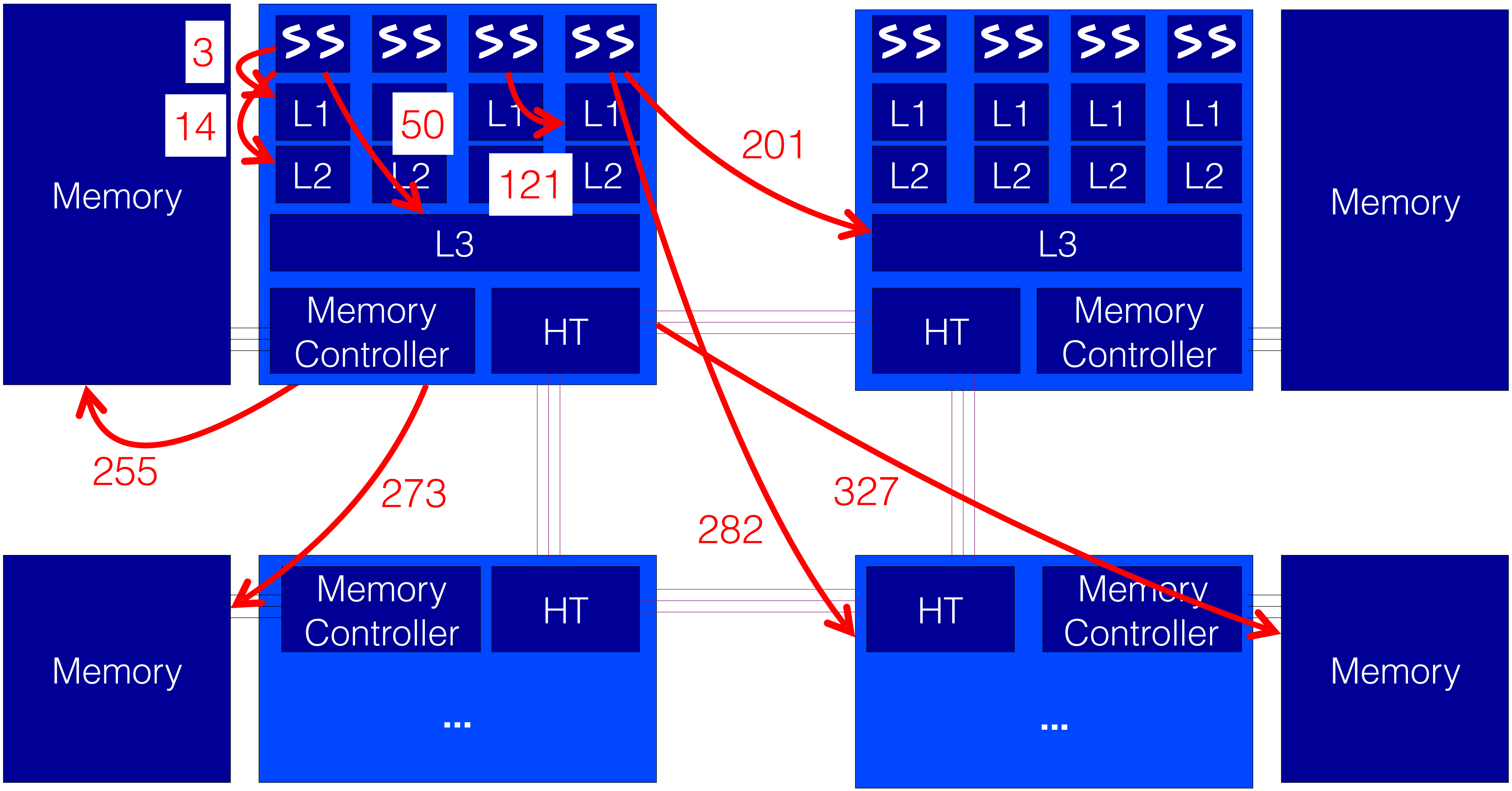
Multiprocessor Architectures

- (cache coherent) NUMA
  - like SMP
  - non-uniform memory access: fetch from remote memory



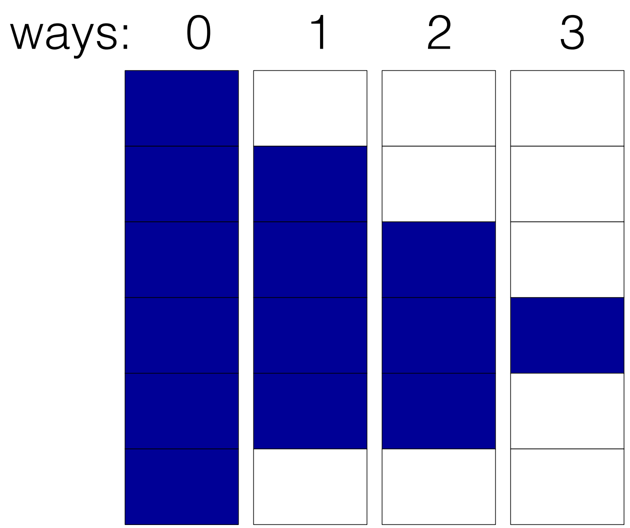
Multiprocessor Architectures

- AMD Opteron [Corey: OSDI '08]



Migration Costs

- Active Cache Working Set
  - cachelines a thread would access again if it would run
  - varies over time
  - ages out after preemption



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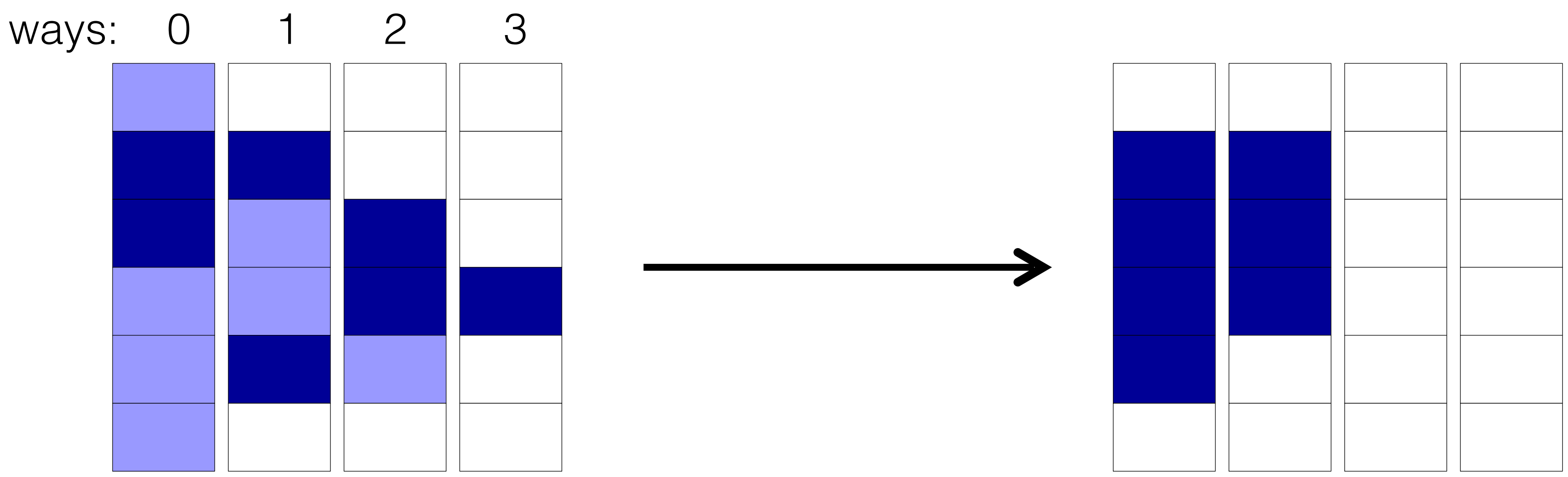

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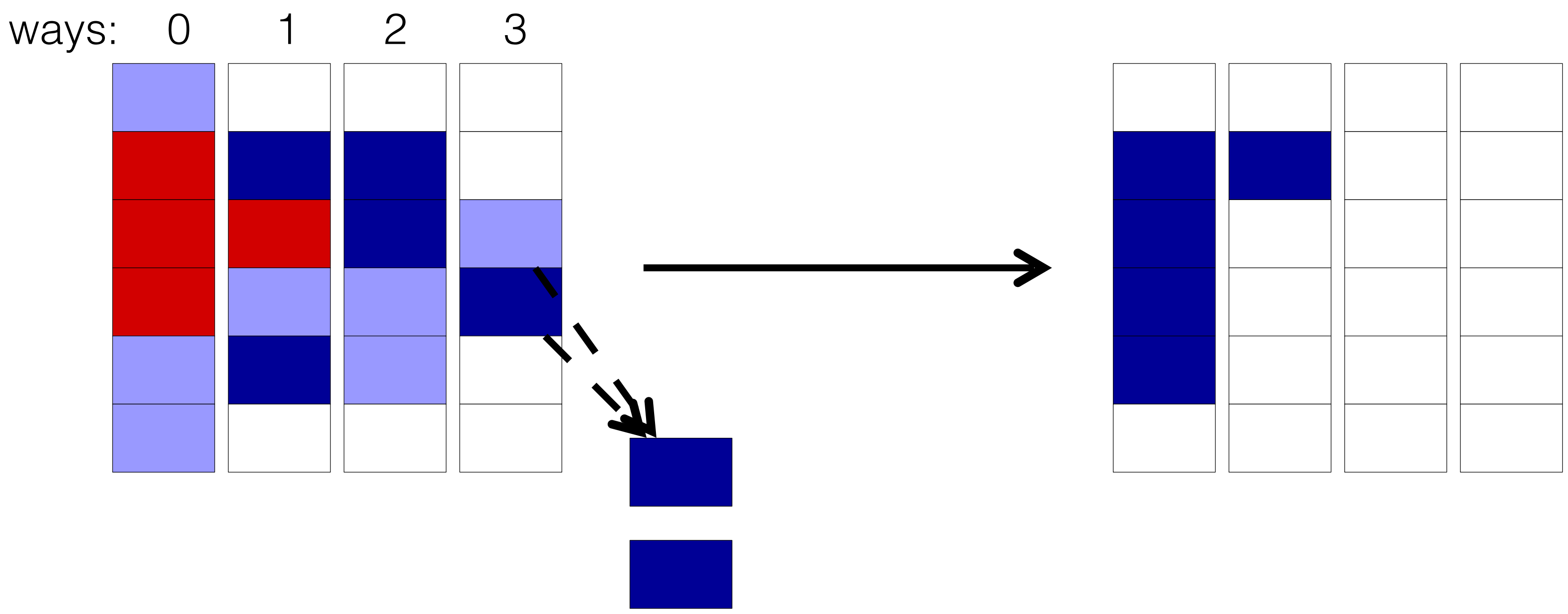
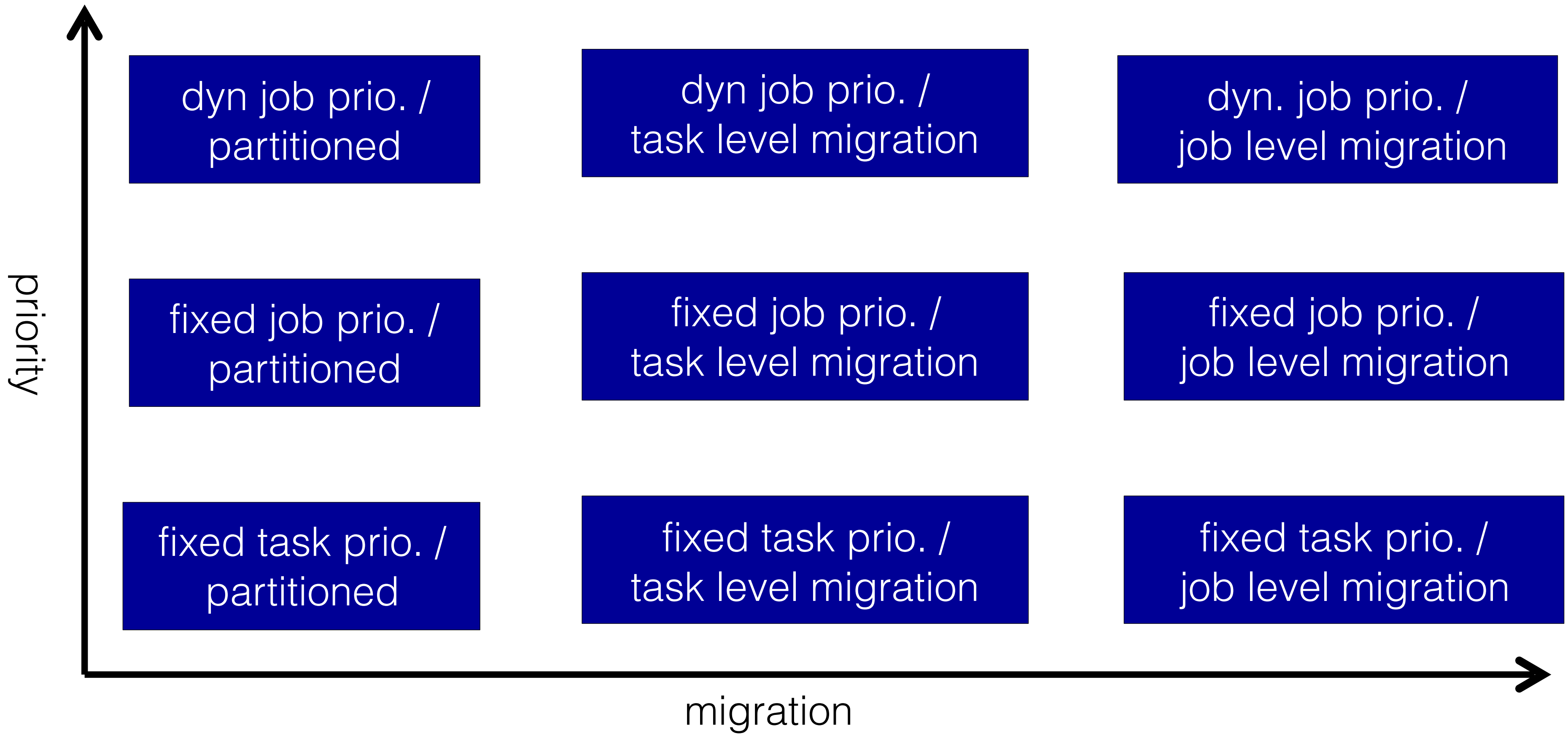


Migration Costs

- Summary
  - migration costs are highly architecture dependent
  - non-trivial to predict
  - may cause a significant delay when a thread resumes execution
- Assumption for the remainder of this lecture:
  - zero preemption and migration costs / attributed to WCET

Design Space of MP Scheduling



Design Space of MP Scheduling

Partitioned	Global	
dyn job prio. / partitioned	dyn job prio. / task level migration	dyn. job prio. / job level migration
fixed job prio. / partitioned	fixed job prio. / task level migration	fixed job prio. / job level migration
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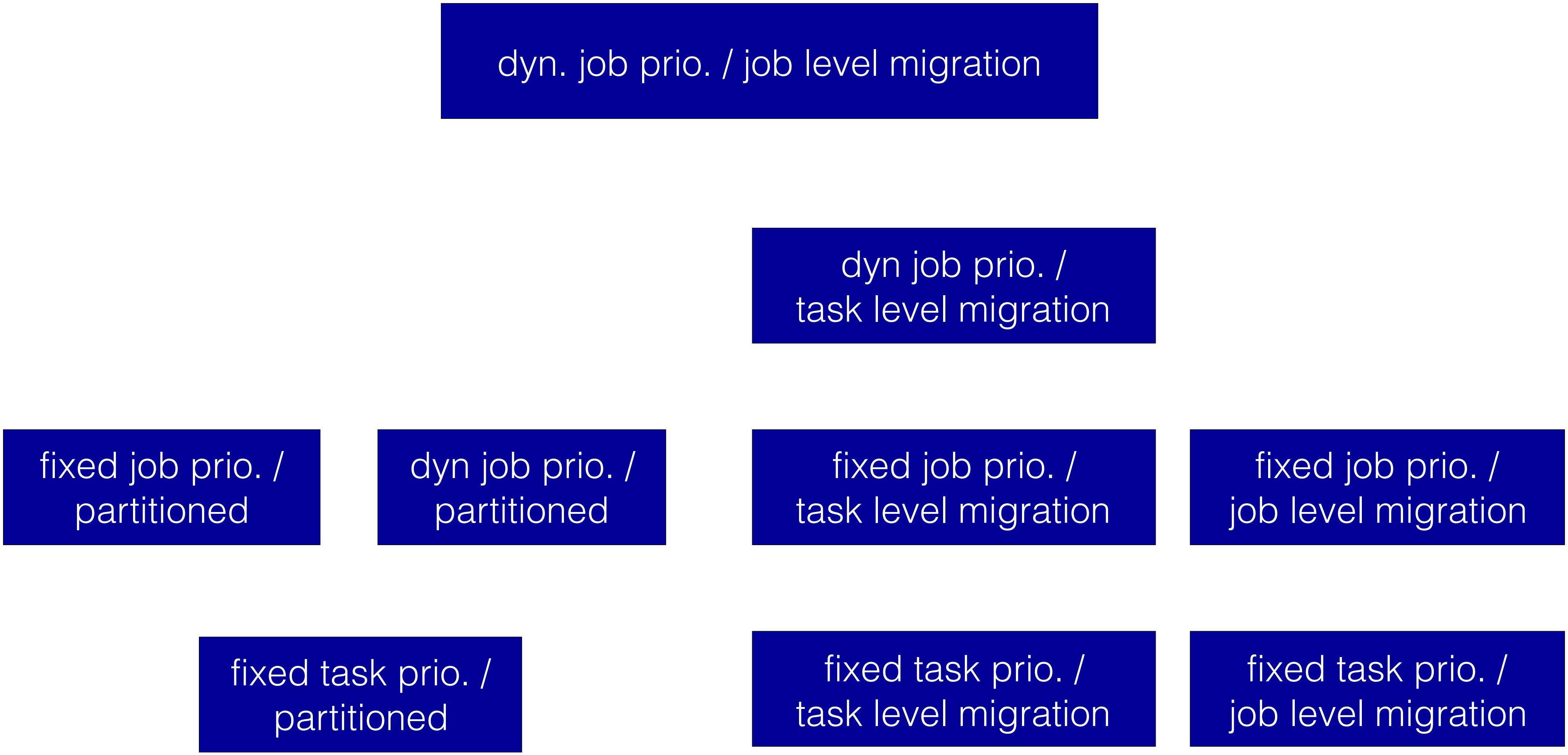
Design Space of MP Scheduling

Relative ordering between classes of scheduling algorithms



Later in this lecture

Outline

- Introduction
- Terminology, Notation and Assumptions
- Anomalies + Impossibility Results
- Partitioned Scheduling
- Global (Task-Lvl migration) Scheduling
  - G-FTP (e.g., G-RMS)
  - G-EDF
- Optimal MP Scheduling
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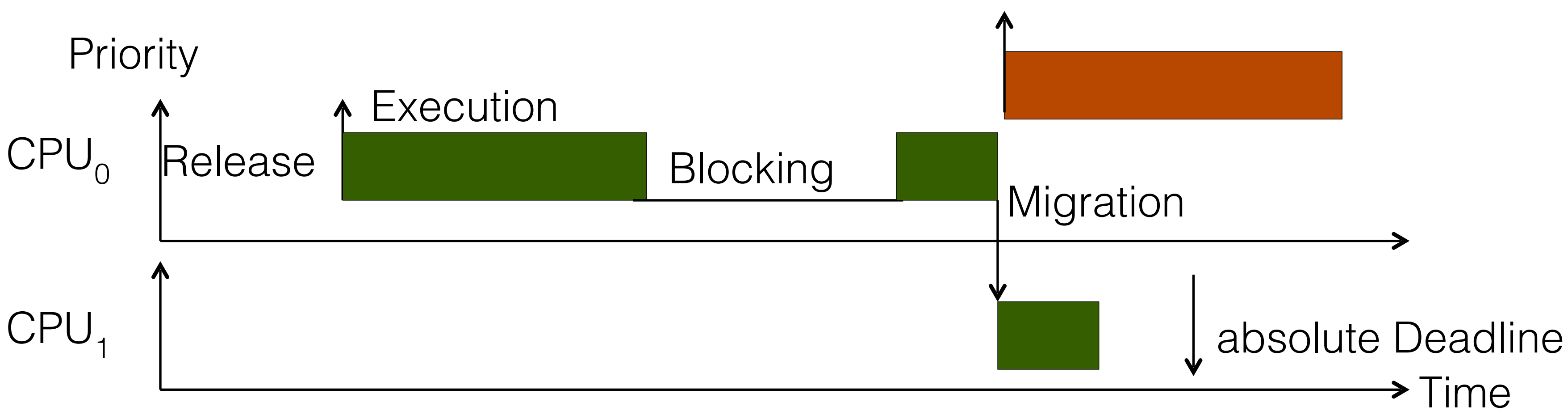
Terminology, Notation and Assumptions

- Periodic Tasks
  - Task $t_i = (P_i, D_i, C_i)$ P_i = const.
- Sporadic Tasks
  - P_i = Minimal Interrelease Time
- Deadlines
  - implicit deadline: $D_i = P_i$ (relative deadline = period)
  - constrained: $D_i \leq P_i$ (relative deadline < period)
  - arbitrary (deadline may be after period end)
- Utilization $U_i = \frac{C_i}{P_i}$
- Density $\partial_i = \frac{C_i}{min(D_i, P_i)}$

Terminology, Notation and Assumptions

- Assumptions for the remainder of this lecture
  - independent tasks
  - fully preemptible / migratable (negligible costs)
  - unlimited number of priorities
  - tasks are single threaded: a job can utilize only 1 CPU at a time
  - jobs do not block (shared resources later in this lecture)
  - pictures show schedules for 2 CPUs



Terminology: P-DMS / G-RMS / G-EDF

- Scheduling Algorithms:

Deadline Monotonic Scheduling:
prio inverse proportional to deadline

Rate Monotonic Scheduling:
prio inverse proportional to period

Earliest Deadline First:
job prio. inverse proportional to deadline

- P-DMS

G-RMS / G-EDF

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- P-DMS

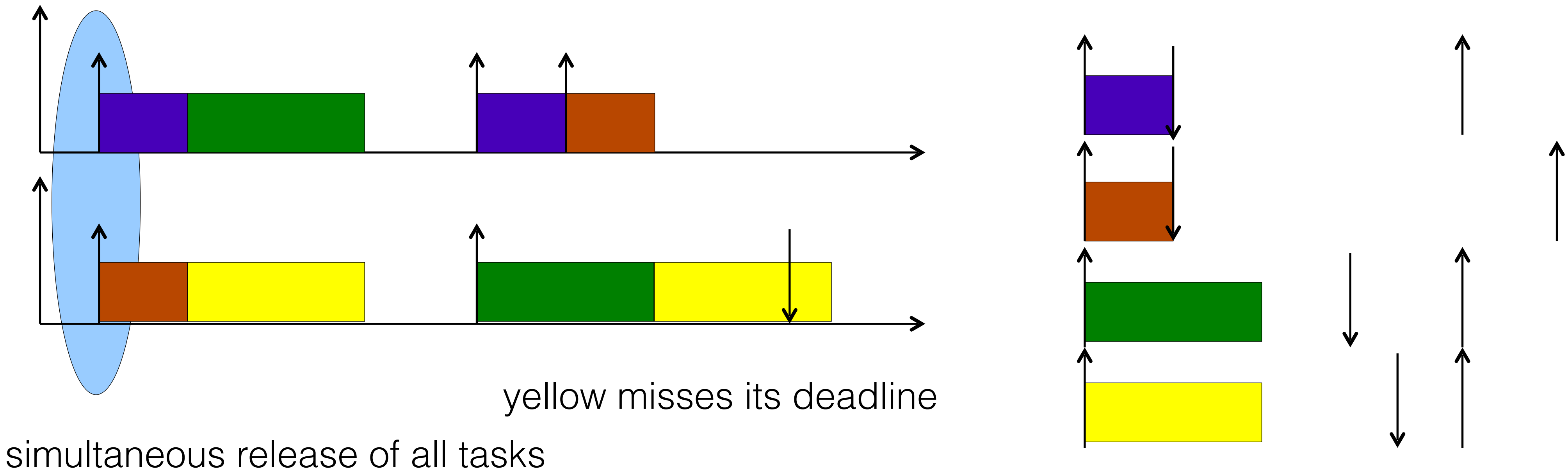
G-RMS / G-EDF

Global:
- threads may migrate to other CPUs
- scheduler picks thread from global ready queue
- accesses to ready queue must be synchronized

Partitioned:
- assign threads to processors
- scheduler picks threads from local (per CPU) ready queue
- no synchronization overhead for accessing the ready queue

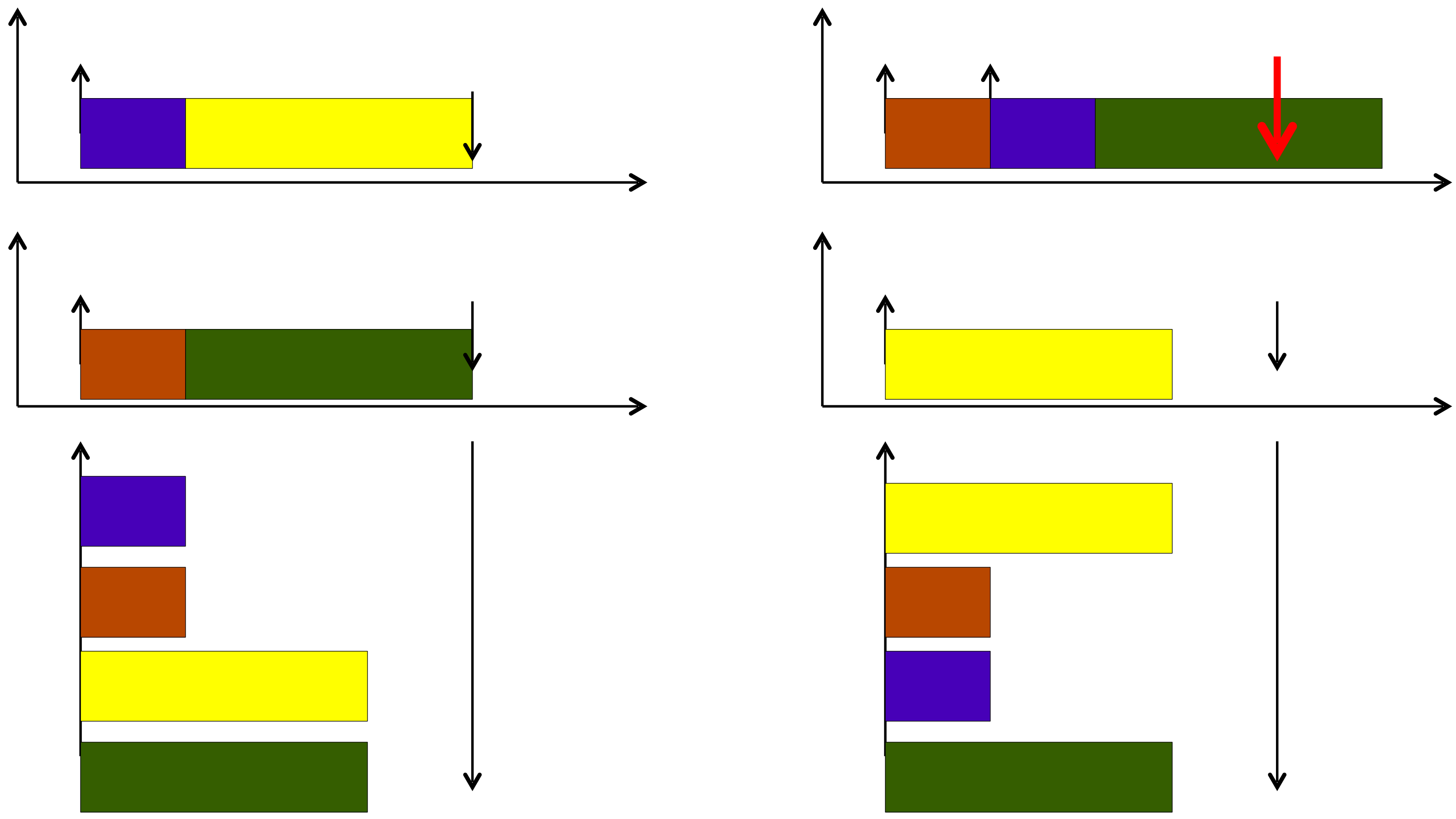
Anomalies

- Simultaneous Release is not Critical Instance [Lauzac '98]
  - longer response time in second period



Anomalies

- Response time (of green) depends not only on set of higher prioritized tasks but also on their relative priority ordering



Sustainability [Baruah '06]

- A schedulable workload remains schedulable if we
  - decrease the execution time of a task (predicability)
    - otherwise, WCET won't work as admission criterion
  - increase the minimal interrelease time (period) of a task
    - otherwise, more frequent recurrence is no safe approximation
  - increase the relative deadline of a task
    - otherwise, earlier deadline is no safe approximation

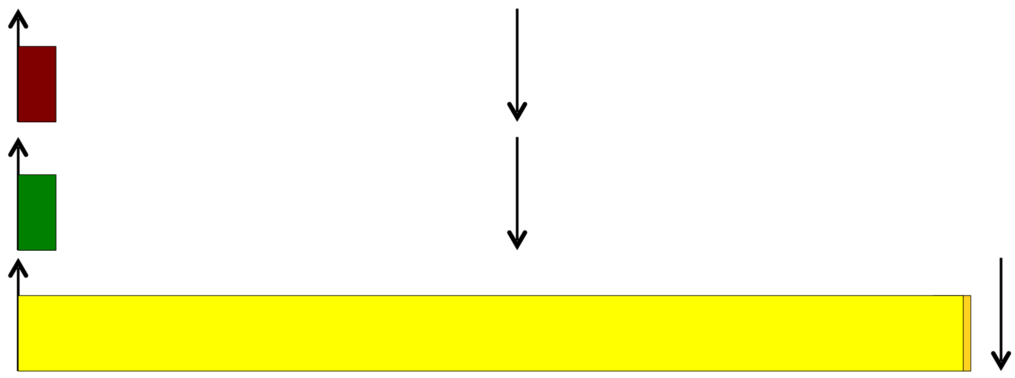
- G-FTP + G-EDF are not sustainable if #CPUs > 1

- all preemptive FJP / FTP algorithms are predictable

Fixed Job Priority Fixed Task Priority

Dhall Effect

- The utilization bound of Global EDF is as low as $U_{EDF} = 1 + \varepsilon$
  - m tasks with short periods and infinitesimal low U_i (e.g., $U_i = \varepsilon$)
  - 1 task with larger period and U_j close to 1 (e.g., $U_j > (2 - \varepsilon) / 2$)



- Dhall Effect does not manifest if $U_i < 41$ %

Dhall Effect

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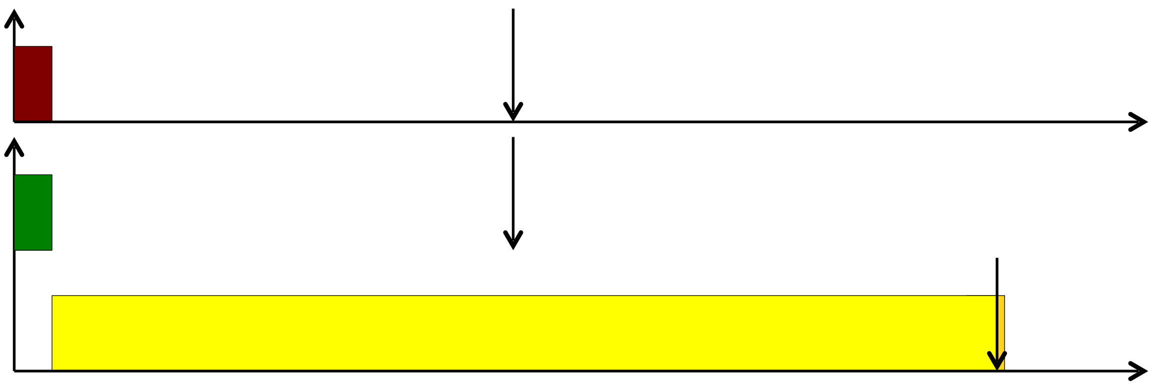

- Dhall Effect does not manifest if $U_i < 41\ \%$

some Impossibility Results

- [Hong '88]
  - No optimal online MP scheduling algorithm for arbitrary collections of jobs, unless all jobs have the same relative deadline.
- [Dertouzos '89]
  - Even if execution times are known precisely, clairvoyance for job arrivals is necessary for optimality.
- [Fisher '07]
  - No optimal online algorithm for sporadic tasksets with constrained or arbitrary deadlines.

Partitioned Scheduling

Partitioned		
dyn job prio. / partitioned	dyn job prio. / task level migration	dyn. job prio. / job level migration
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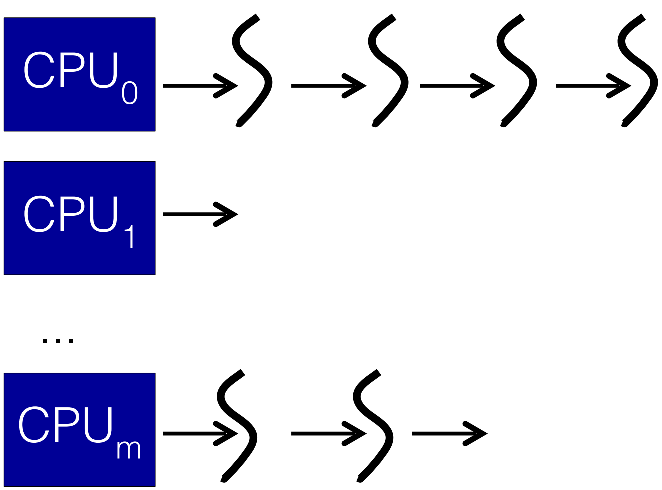
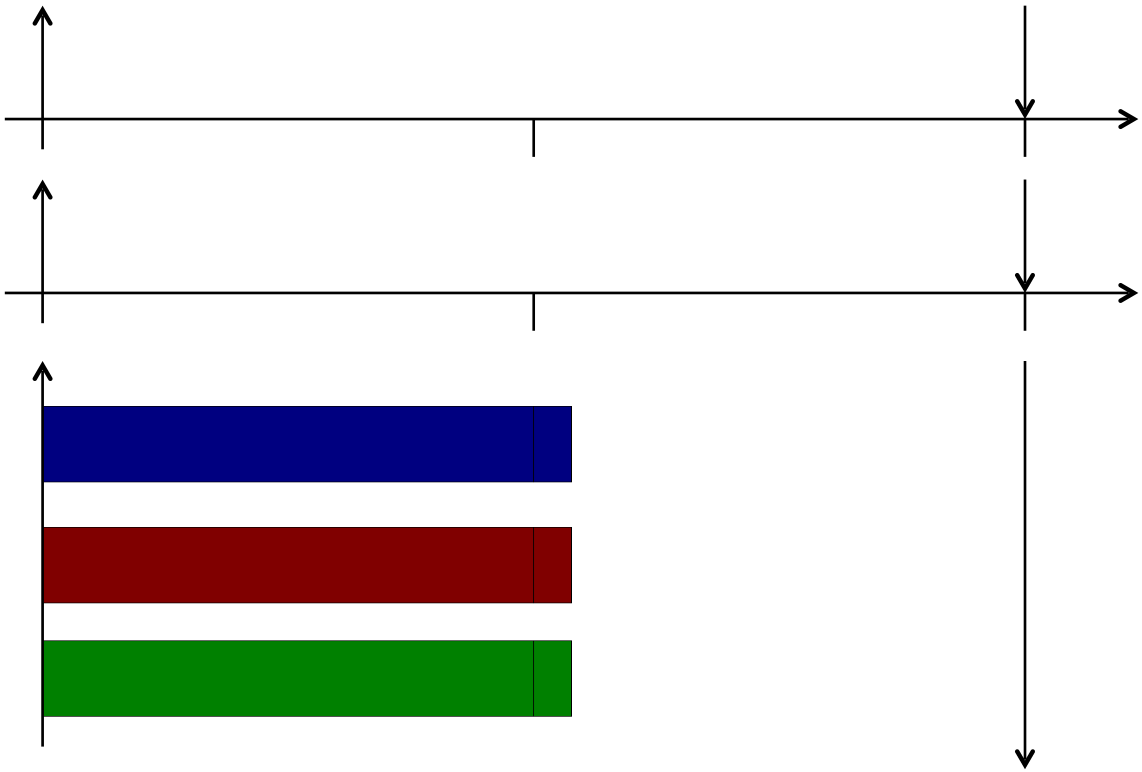
Partitioned Scheduling

- Split workload by allocating tasks to CPUs
- Run allocated task with UP scheduling algorithm
  - reap benefit of well known UP results
  - optimal task allocation is NP complete:
    - pack n tasks with density d_i on m CPUs with capacity $d_{max} = 1$
    - Bin-packing



Partitioned Scheduling

- Utilization bound for implicit deadline workloads [Anderson '01]



$$U_{opt} = \frac{m+1}{2}$$

No partitioning scheduling algorithm can produce a feasible schedule of m+1 tasks with execution time 1+e and period of 2 on m processors

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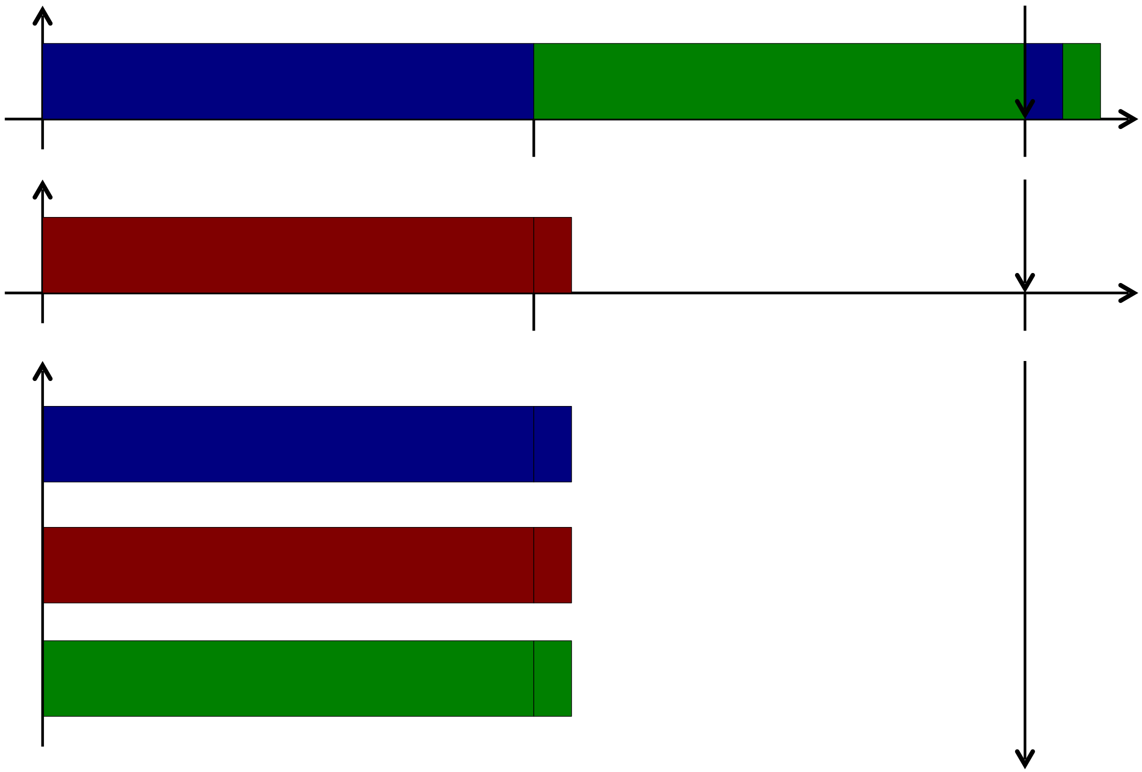

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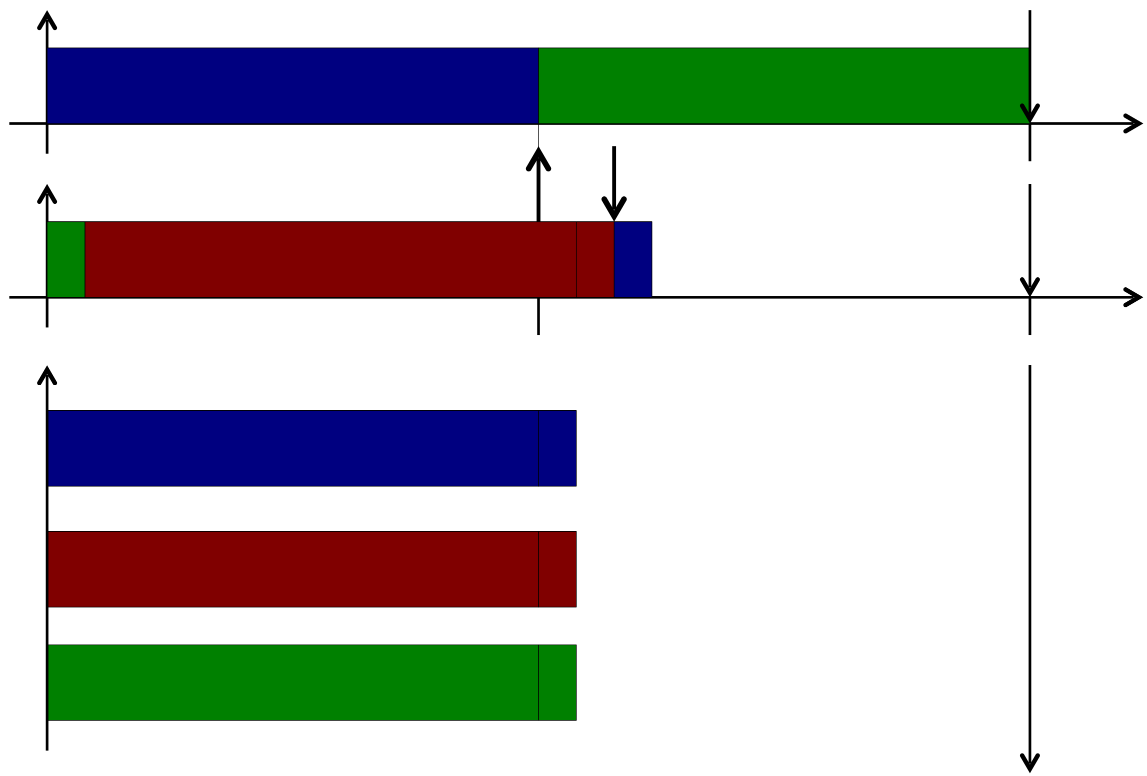

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Easy if blue and green can migrate to CPU_2

PDMS-HPTS-DS [Lakshmanan '09]

- Can we improve on Anderson's Utilization Bound?
  - by allowing a few jobs to migrate

- **PDMS** Partitioned Deadline Monotonic Scheduling
- **HPTS** Split Highest Priority Task
- **DS** Allocate according to Highest Density First

PDMS-HPTS-DS

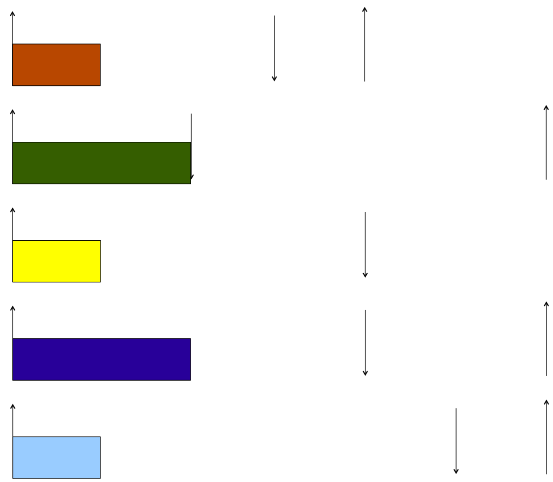
$\tau_1 = (4,3,1)$

$\tau_2 = (6,2,2)$

$\tau_3 = (4,4,1)$

$\tau_4 = (6,4,2)$

$\tau_5 = (6,5,1)$



$u_1 = 0.25 \quad \delta_1 = 1/3 = 0.33$

$u_2 = 0.33 \quad \delta_2 = 1$

$u_3 = 0.25 \quad \delta_3 = 1/4 = 0.25$

$u_4 = 0.33 \quad \delta_4 = 1/2 = 0.5$

$u_5 = 0.16 \quad \delta_5 = 1/5 = 0.2$

$u_{sum} = 1.33 \Rightarrow u_{sum} / 2 = 0.66\%$

CPU_0

CPU_1

PDMS-HPTS-DS

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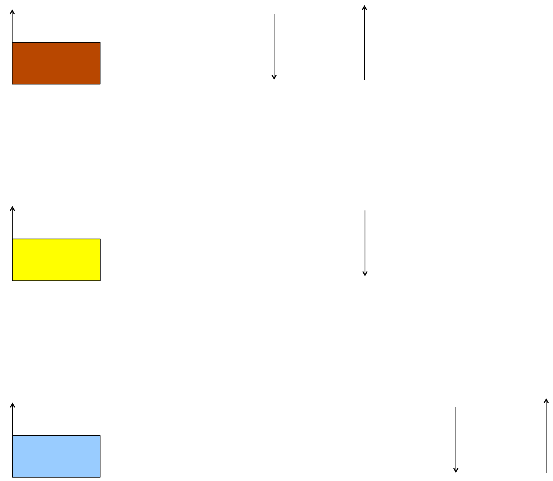

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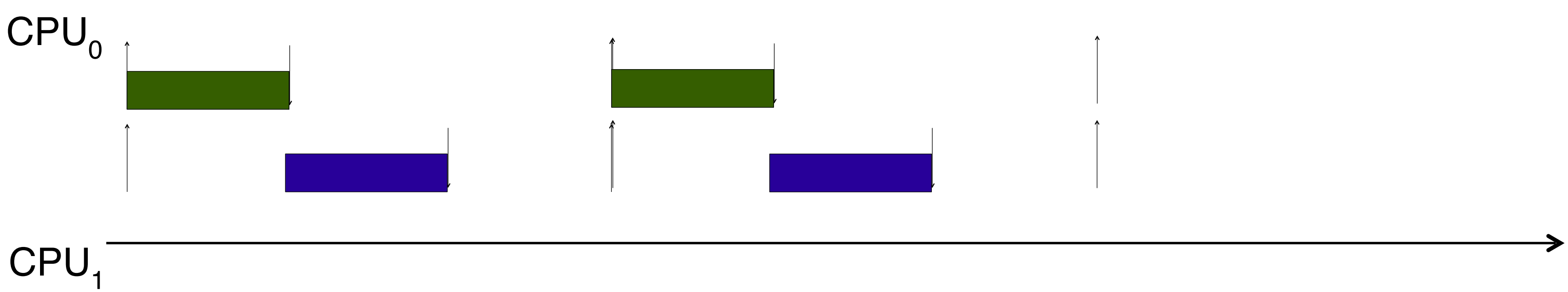

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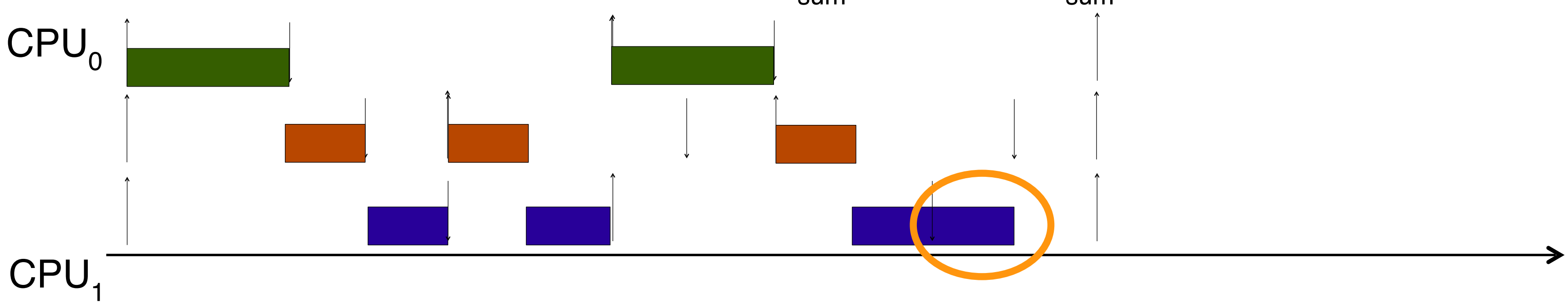

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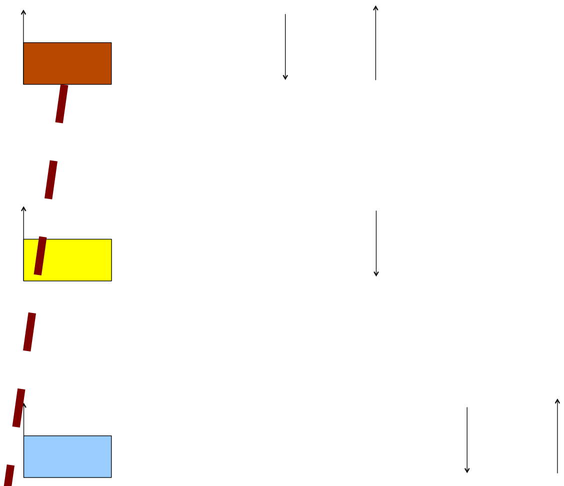

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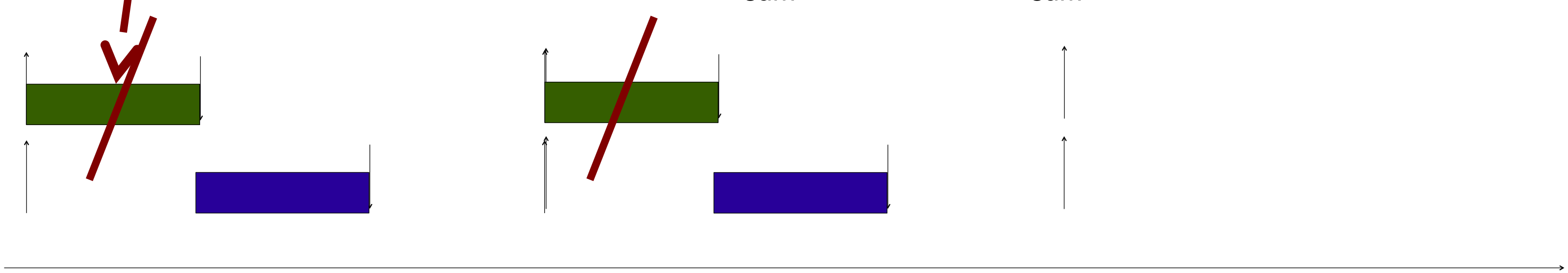

CPU_1

PDMS-HPTS-DS

Splitting HP Task into D'' – R' = D'' = C'
=> deadline of τ'' maximized

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τ_2'' = (5,1,1)
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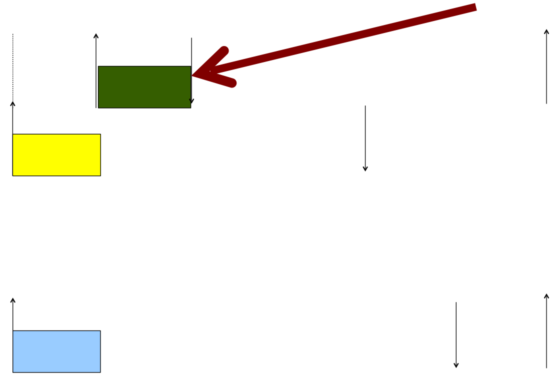

$u_2 = 0.33 \quad \delta_2 = 1$

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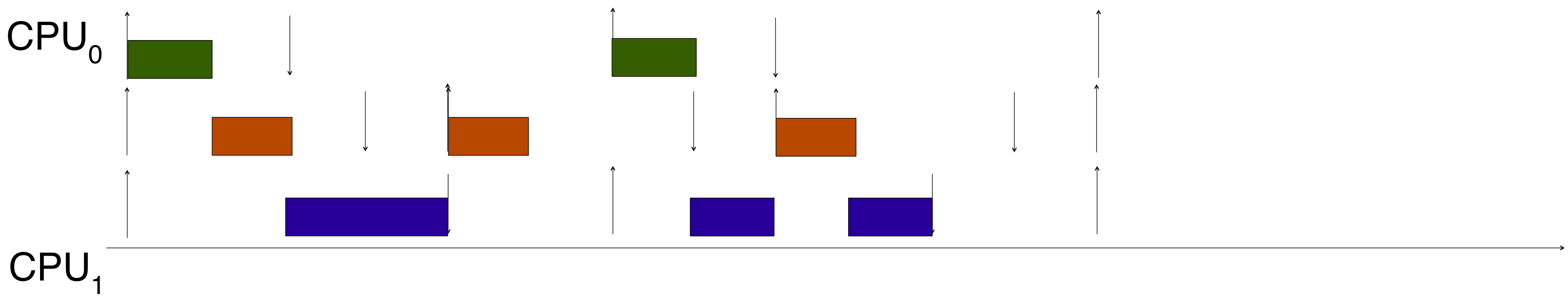

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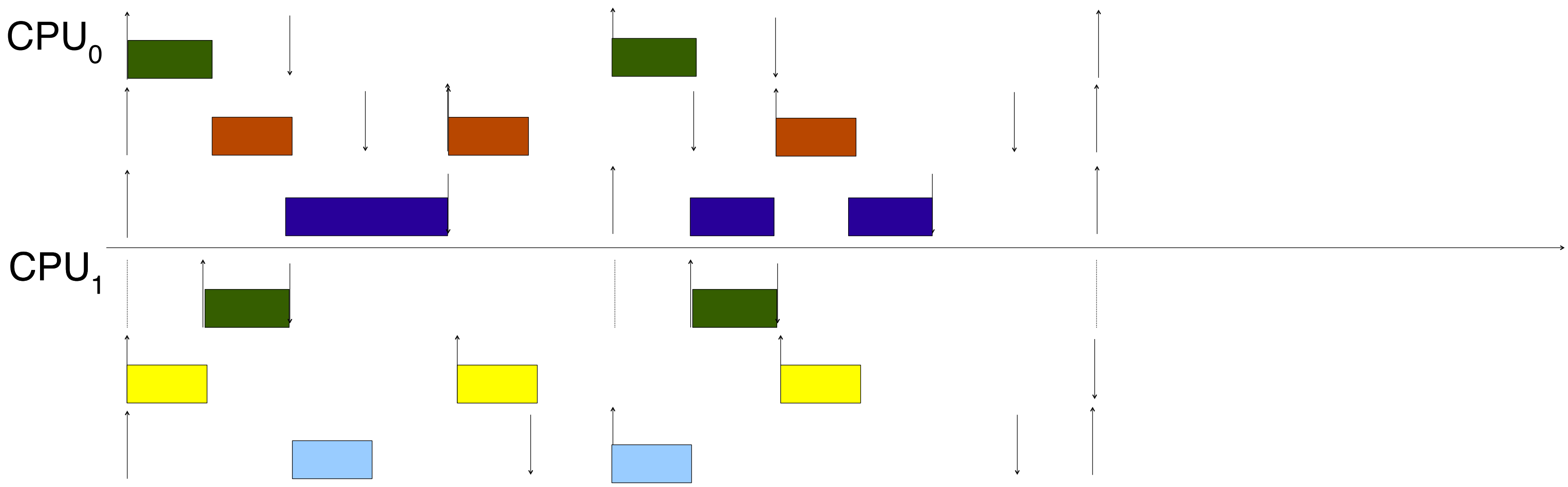
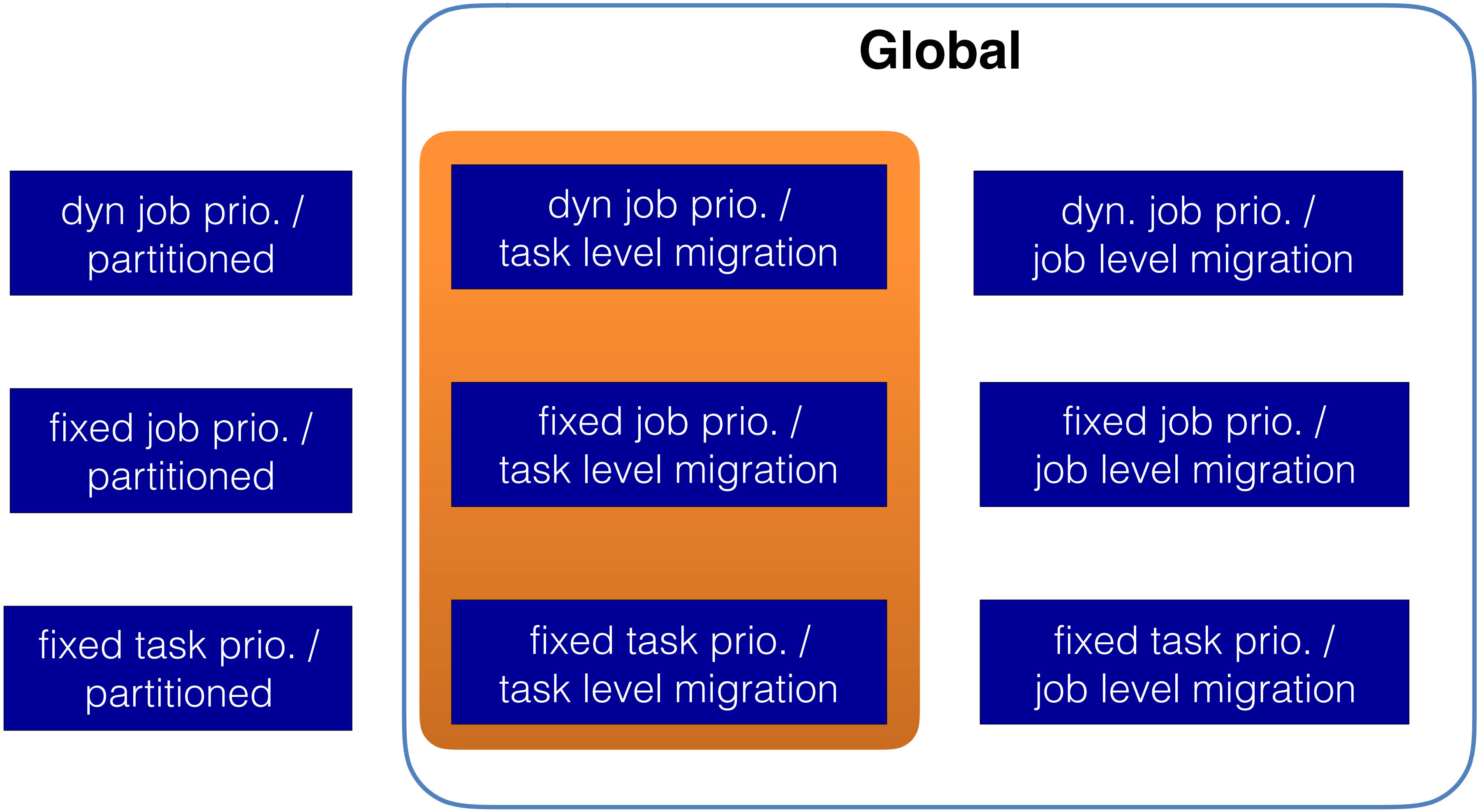


PDMS-HPTS-DS

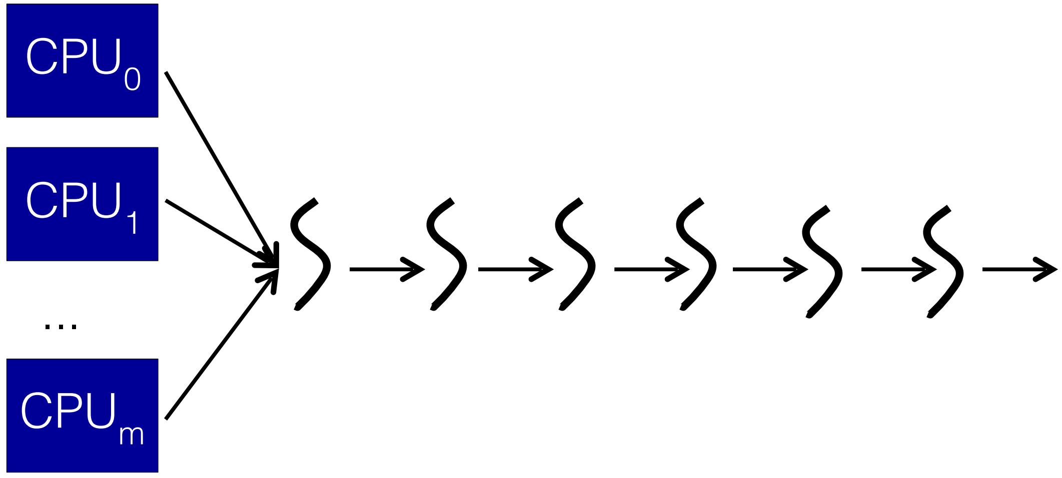
- $U_{PDMS\text{-}HPTS\text{-}DS}$ = 69.3 % if all tasks have a utilization $U_i < 41\%$

Global Scheduling



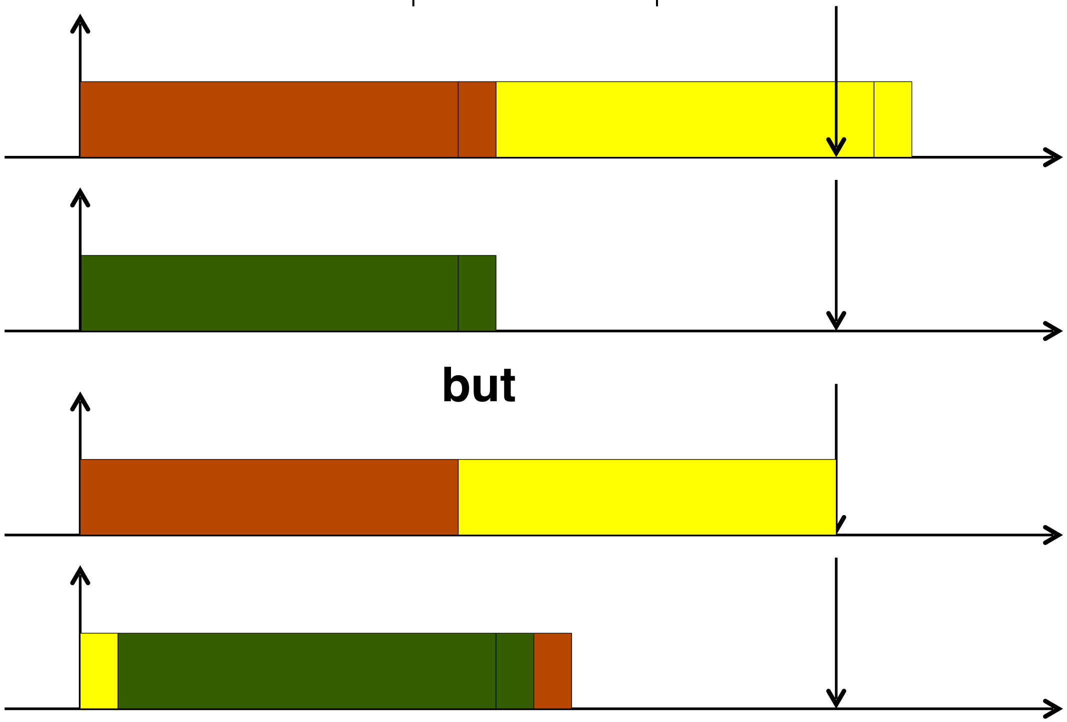
Global Scheduling

- Always pick the ready jobs of the m most “important” tasks
- A task may migrate
  - When a new job of a task is released it may receive a different CPU; once started, a job is no longer migrated
- No need for allocation / load balancing
  - Load balancing is automatic



Global Scheduling – Utilization Bound

- Utilization bound for global fixed-job priority algorithms
  - on m CPUs, G-FJP algorithms cannot schedule m+1 tasks with $C_i = 1 + e$, $P_i = 2$ $(e \rightarrow 0)$



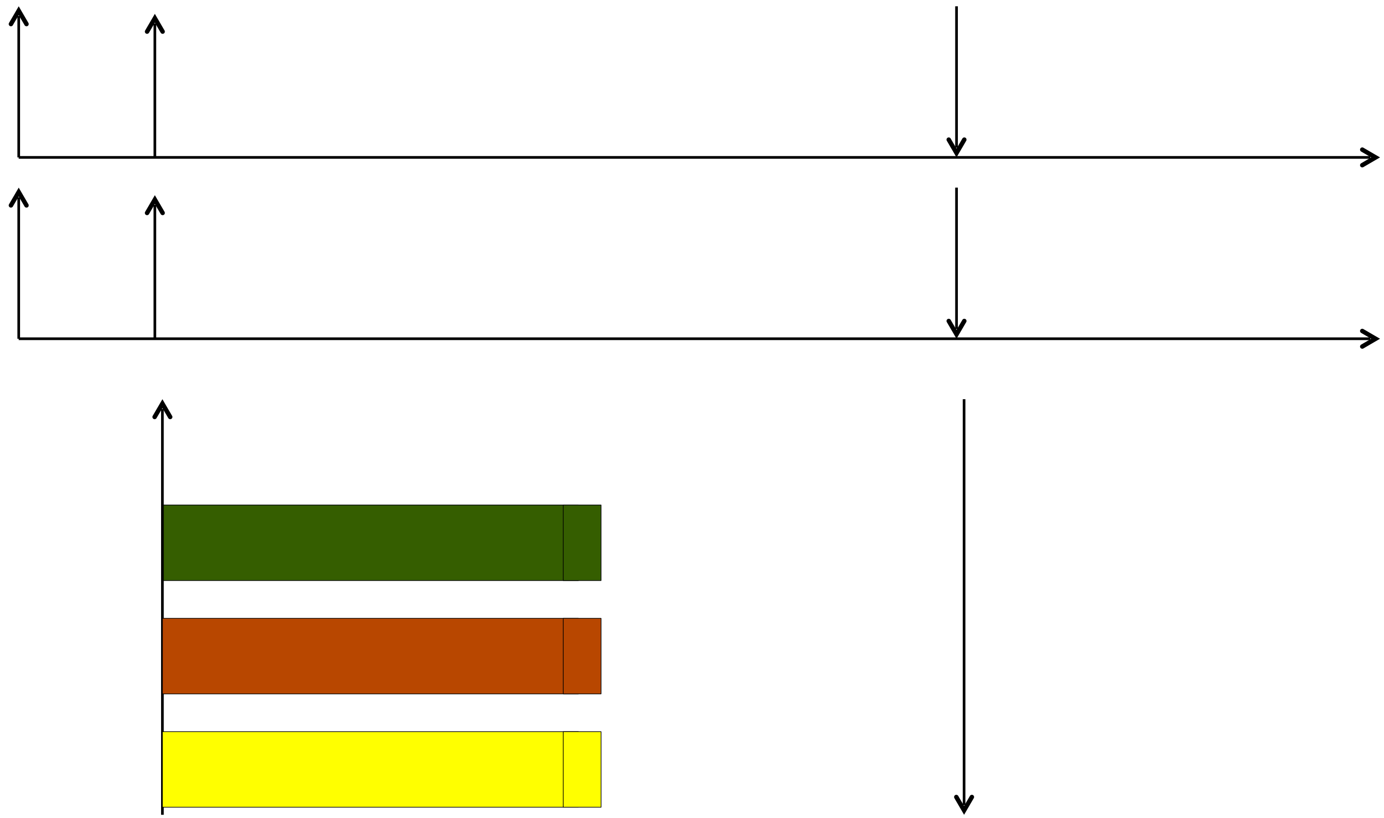
- $U_{OPT} = \lim_{e \rightarrow 0} \frac{(m+1)(1+e)}{2} = \frac{m+1}{2}$

Global Scheduling - Job Level Migration

dyn job prio. / partitioned	dyn job prio. / task level migration	dyn. job prio. / job level migration
fixed job prio. / partitioned	fixed job prio. / task level migration	fixed job prio. / job level migration
fixed task prio. / partitioned	fixed task prio. / task level migration	fixed task prio. / job level migration

PFAIR [Baruah '96]

- Divide timeline into equal length quanta
- At each quanta of length t, allocate tasks to processors such that the accumulated processor time is either $\lceil tu_i \rceil$ or $\lfloor tu_i \rfloor$



PFAIR [Baruah '96]

- Divide timeline into equal length quanta
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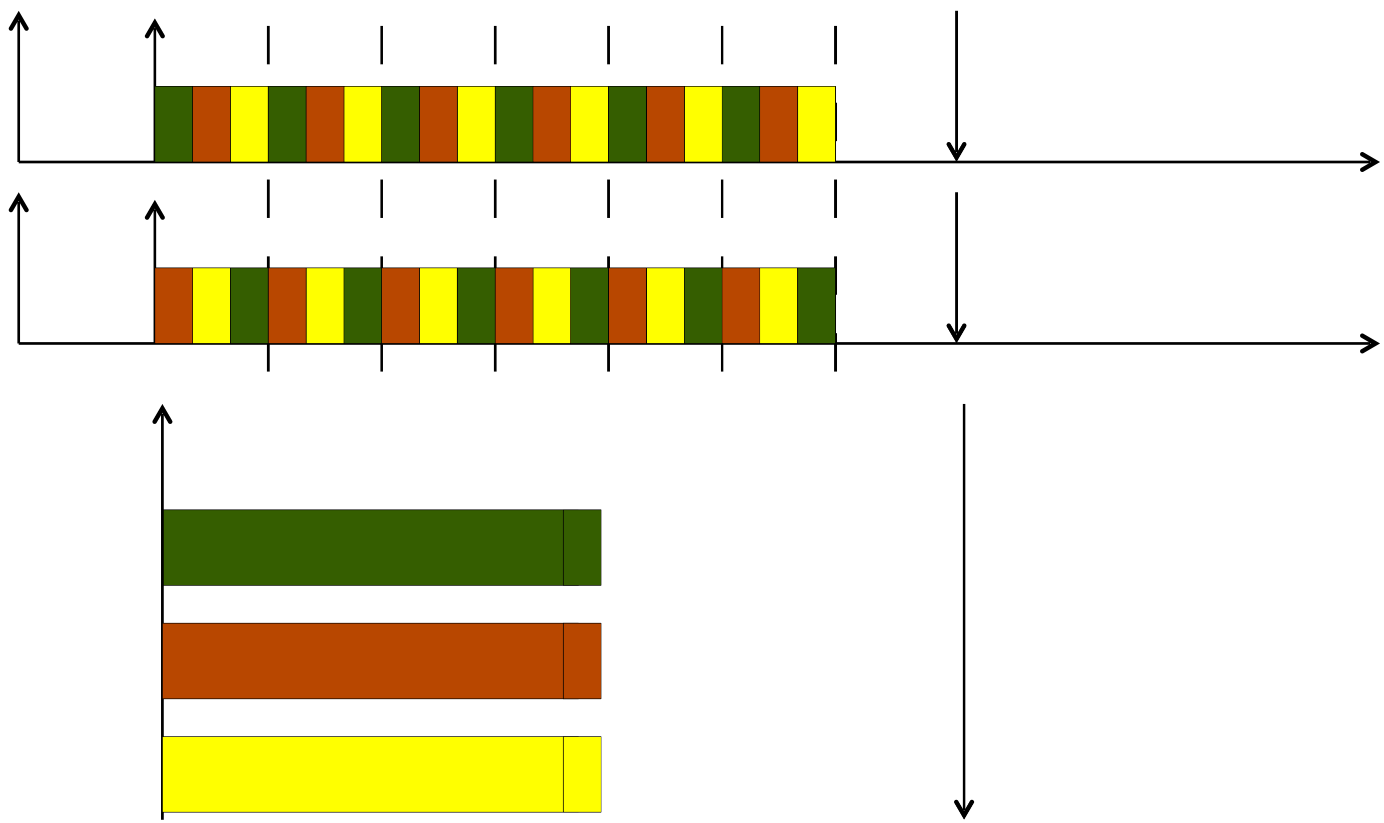

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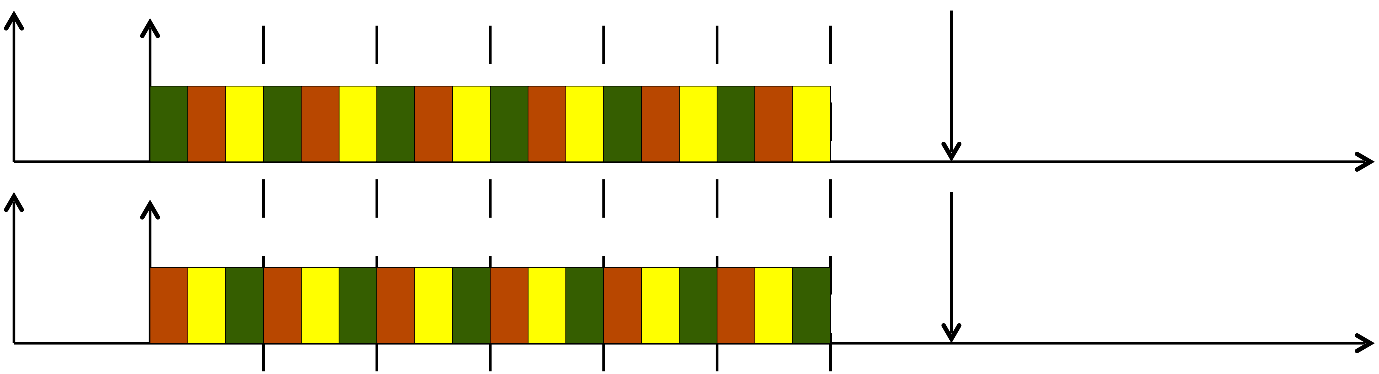


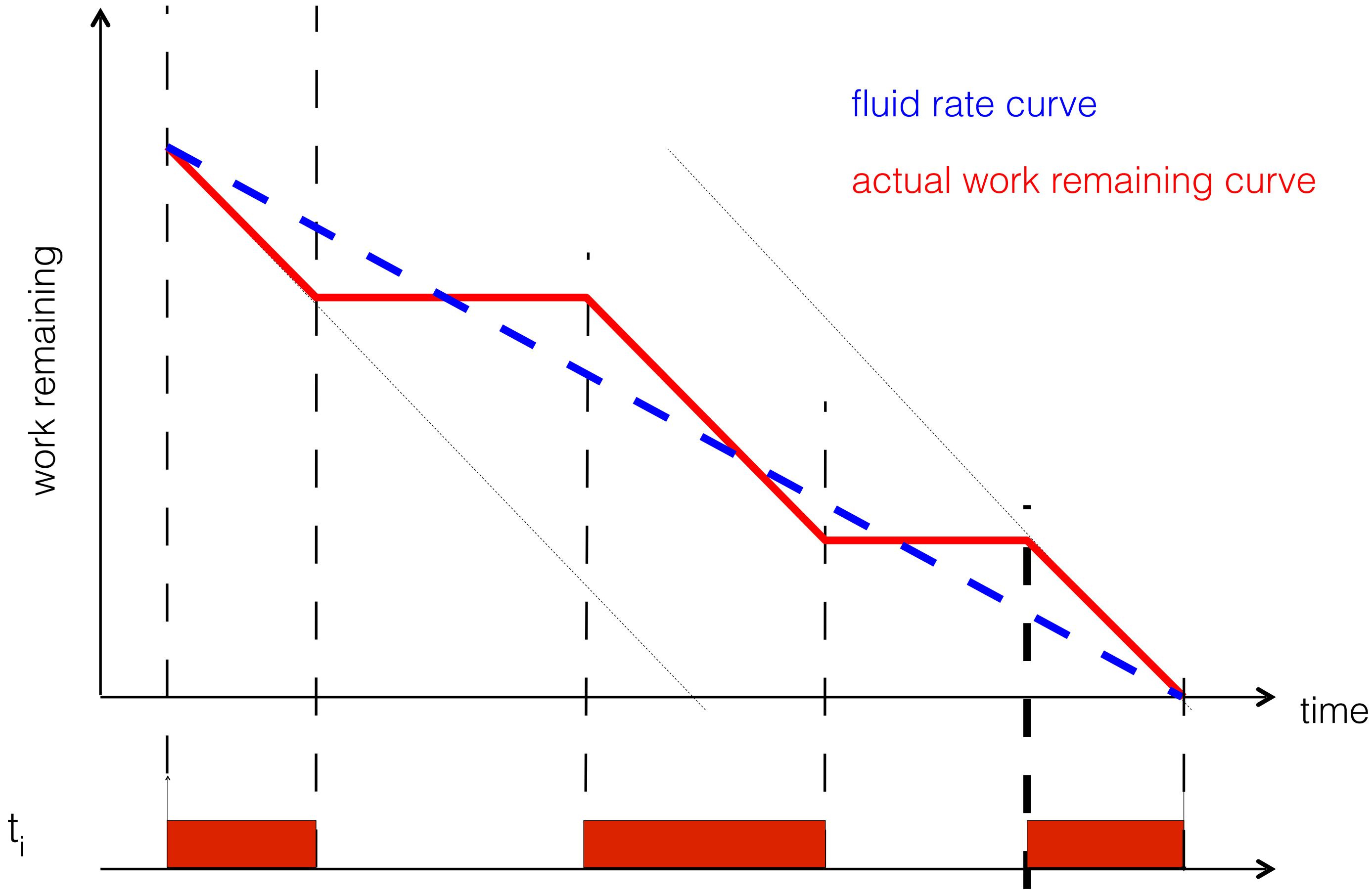
- PFAIR is optimal for periodic implicit deadline tasksets: $U_{OPT} = m$
- Very high preemption and migration costs

DP-Fair [Brandt '10]

- DP-Fair
  - Optimal scheduler for periodic implicit deadline tasksets with a minimal number of preemptions / migrations?
  - Recall [Hong '88]:
    - No optimal MP scheduling algorithm for arbitrary tasksets if not all tasks have the same relative deadline
  - deadline partitioning
    - any task's deadline becomes a deadline for all tasks
  - always run zero laxity jobs
    - laxity = time to deadline – remaining execution time
    - zero laxity => job may miss its deadline if it is not run
  - jobs that twine themselves around the *fluid rate curve* are somehow in good shape

DP-Fair [Brandt '10]

- Fluid Rate Curve



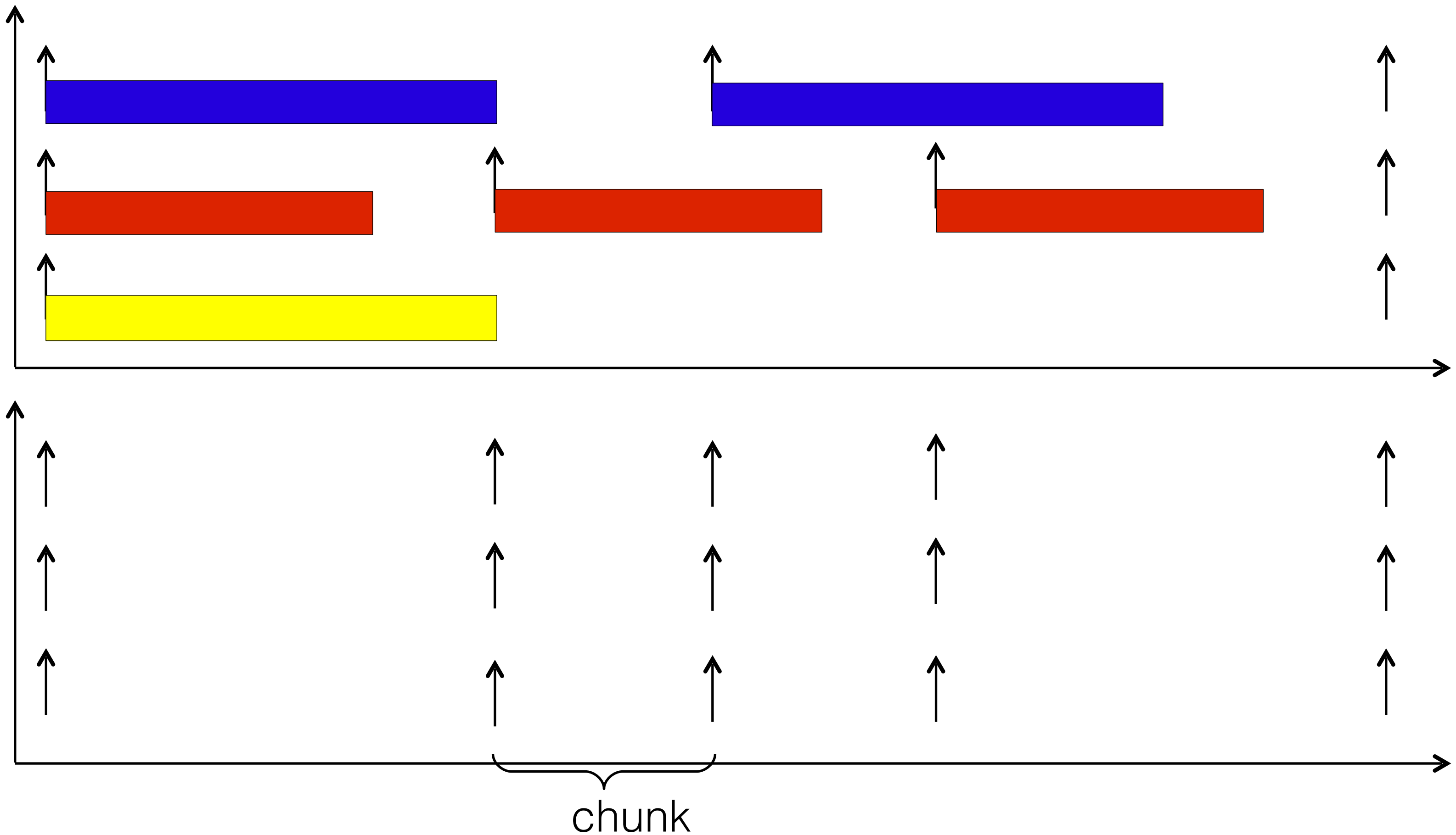
zero laxity event: no more time to run others

DP-Fair [Brandt '10]

- DP-Fair:
  - a family of optimal, deadline partitioning scheduling algorithms
  - Split timeline into chunks according to job deadlines
  - Allocate work to a chunk proportional to U_i
    - local execution time: $C^l_{i,j}=(t_{j+1} – t_j)\ U_i$
  - Rule 1: always run a job with zero local laxity
    - jobs with remaining local execution time = time to end of chunk
  - Rule 2: never run a job with no remaining local work
  - Rule 3: split up idle time proportional to length of chunk
    - allocate at most $(m – U_{sum})\ (t_{j+1} – t_j)$ idle time to chunk j

DP-Fair [Brandt '10]

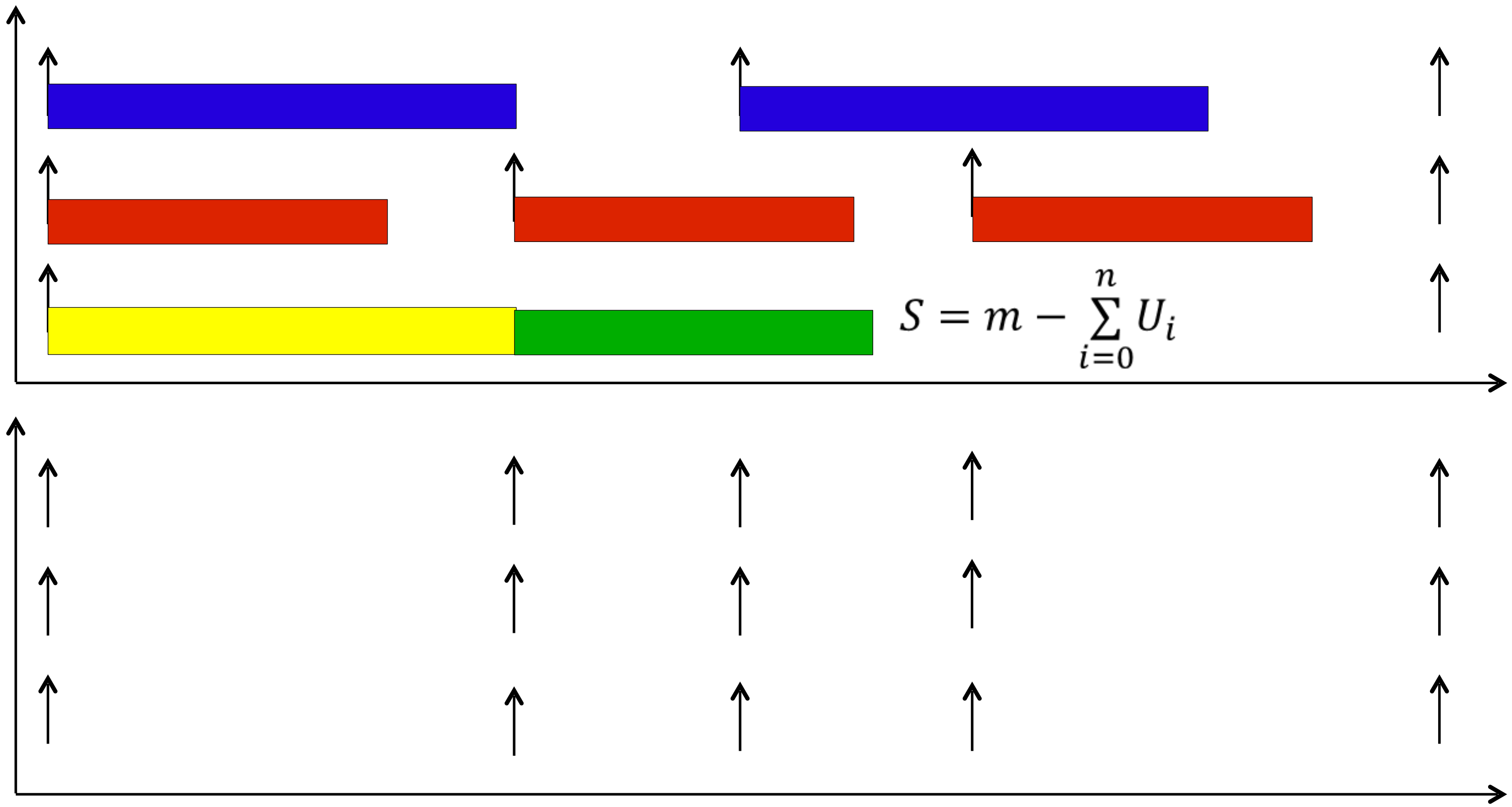
Deadline partitioning



Introduce additional releases / deadlines for all jobs whenever there is such an event for one job in the original schedule => chunks

DP-Fair [Brandt '10]

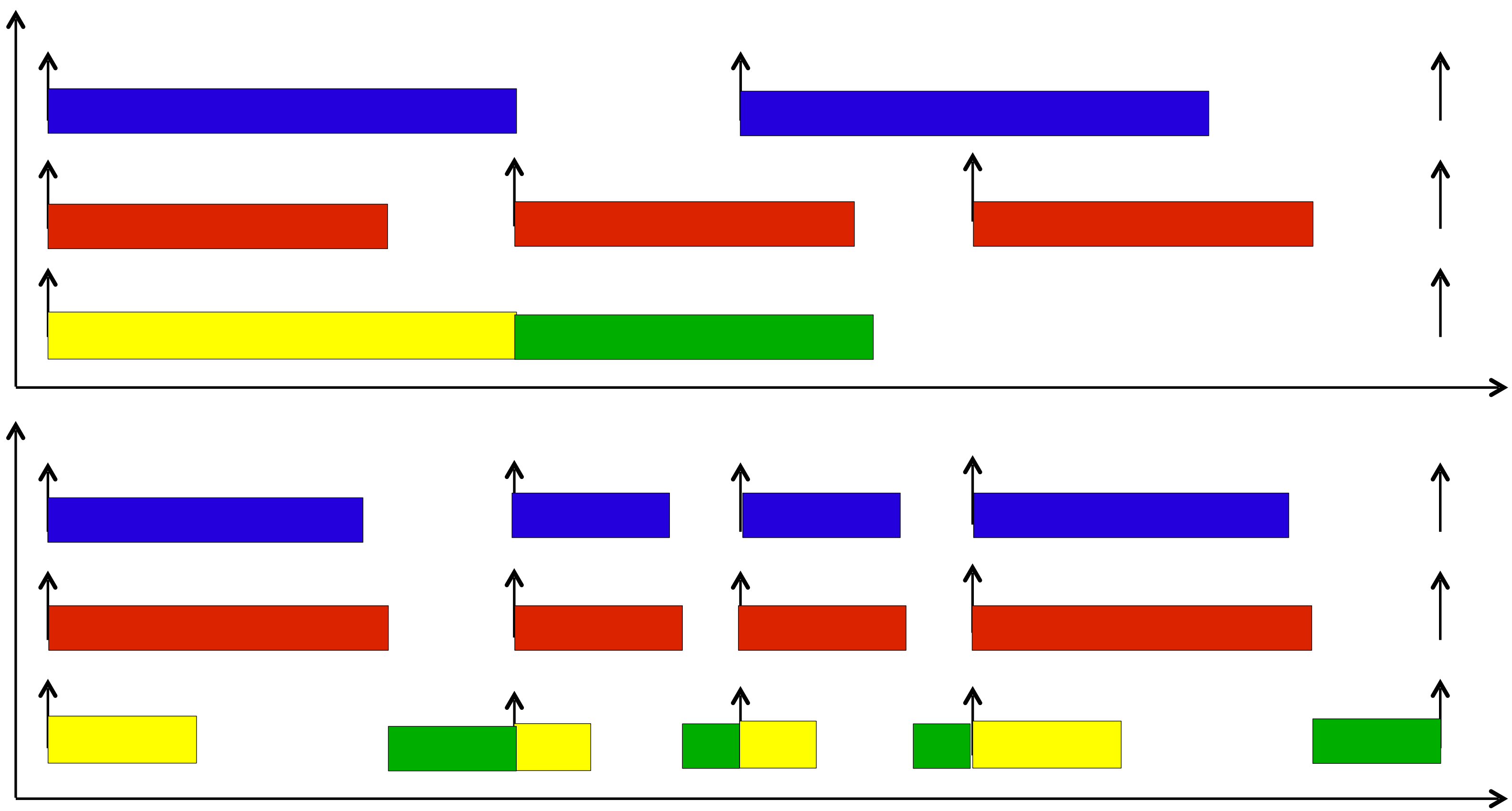
Idle time



Treat idle time as just another job to schedule.

DP-Fair [Brandt '10]

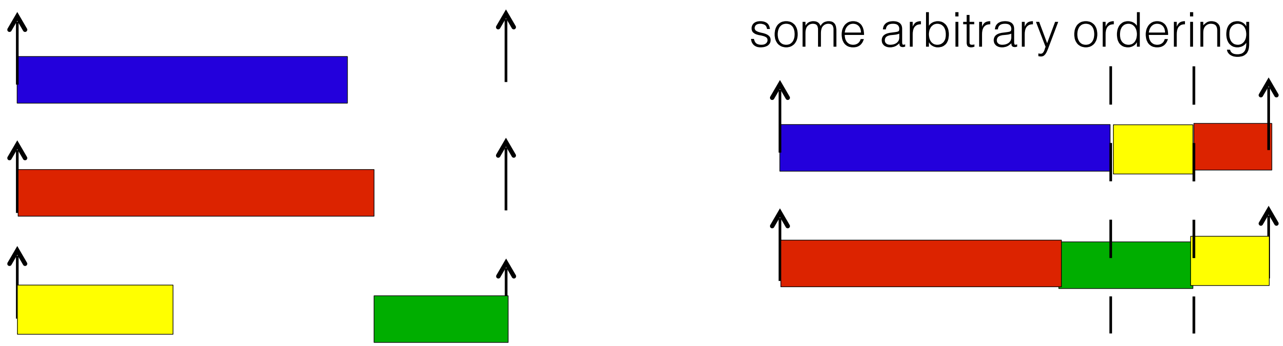
Allocate work proportional to U_i

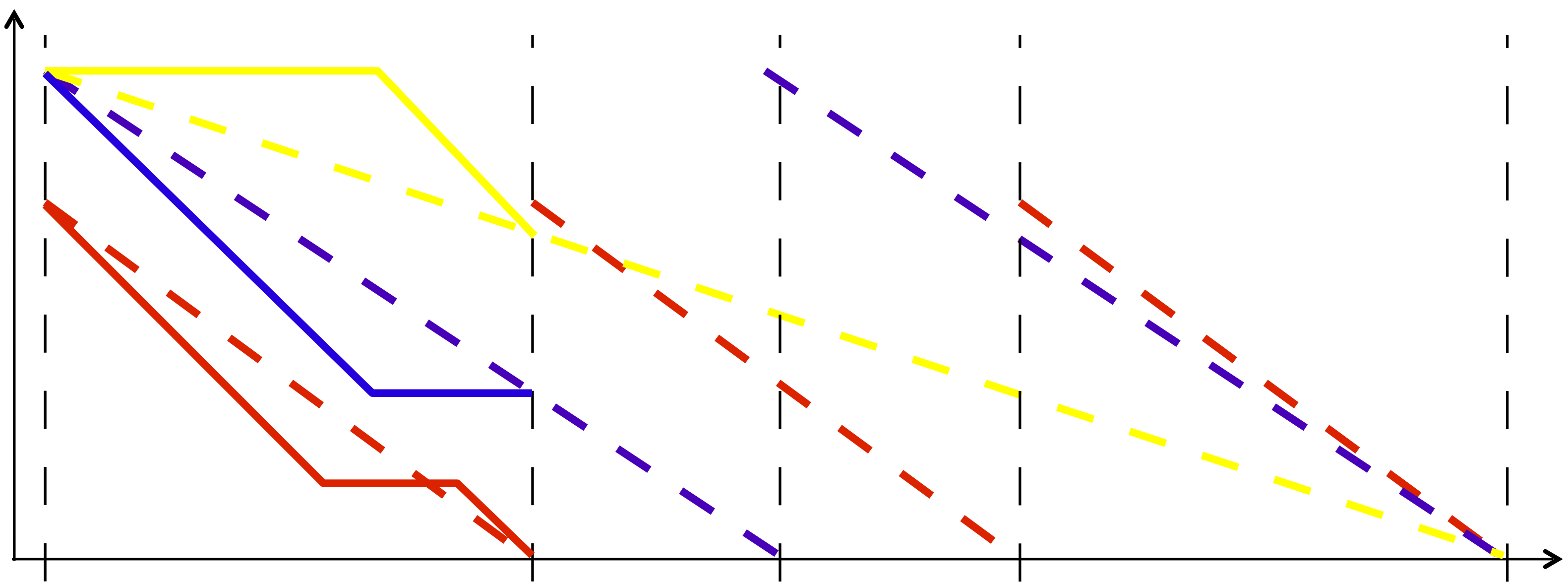


Allocate execution (and idle) time of a job proportionally to its utilization
=> amount of time that this job must run in a given chunk

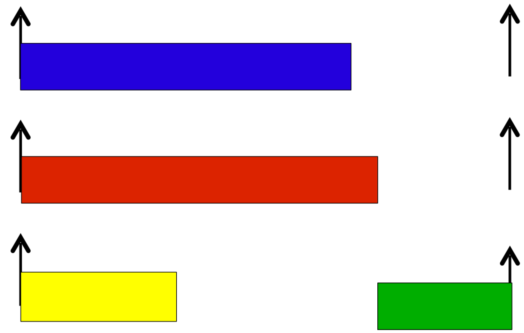
DP-Fair [Brandt '10]

Jobs hit their fluid rate curve at the end of each chunk



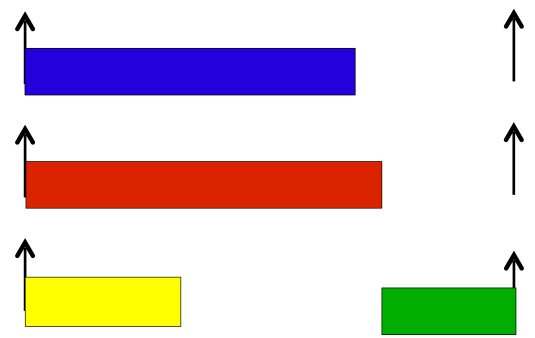


DP-Wrap – a DP-Fair Scheduler



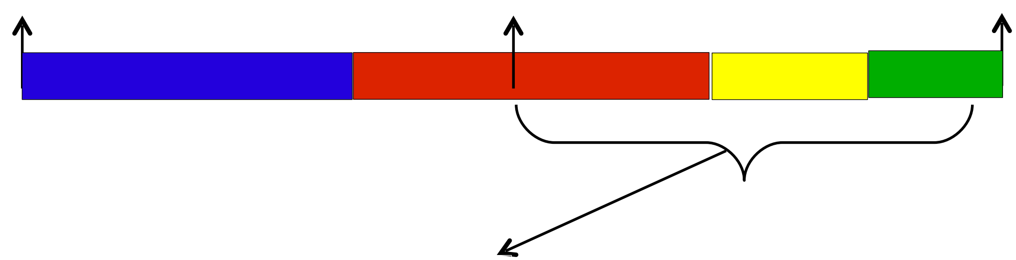
Work allocated to a chunk:
find a schedule

DP-Wrap – a DP-Fair Scheduler



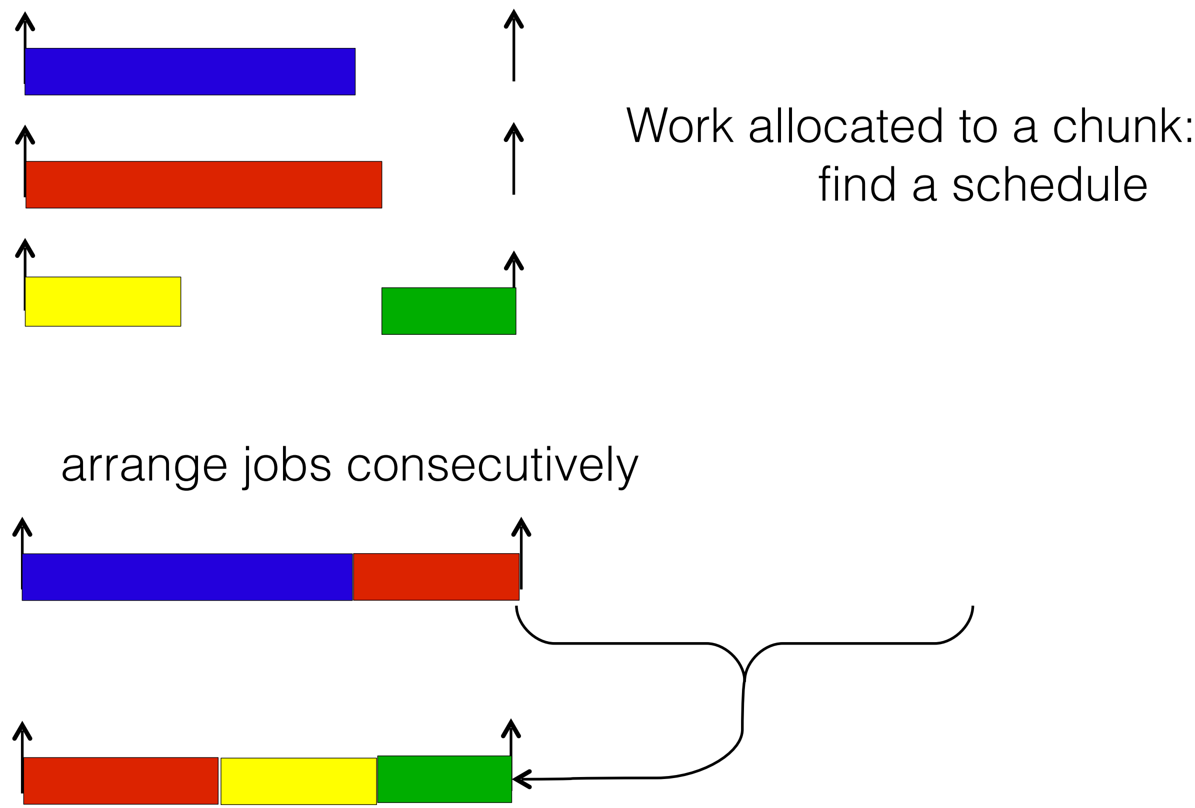
Work allocated to a chunk:
find a schedule

arrange jobs consecutively



wrap around to obtain schedule for 2nd CPU, ...

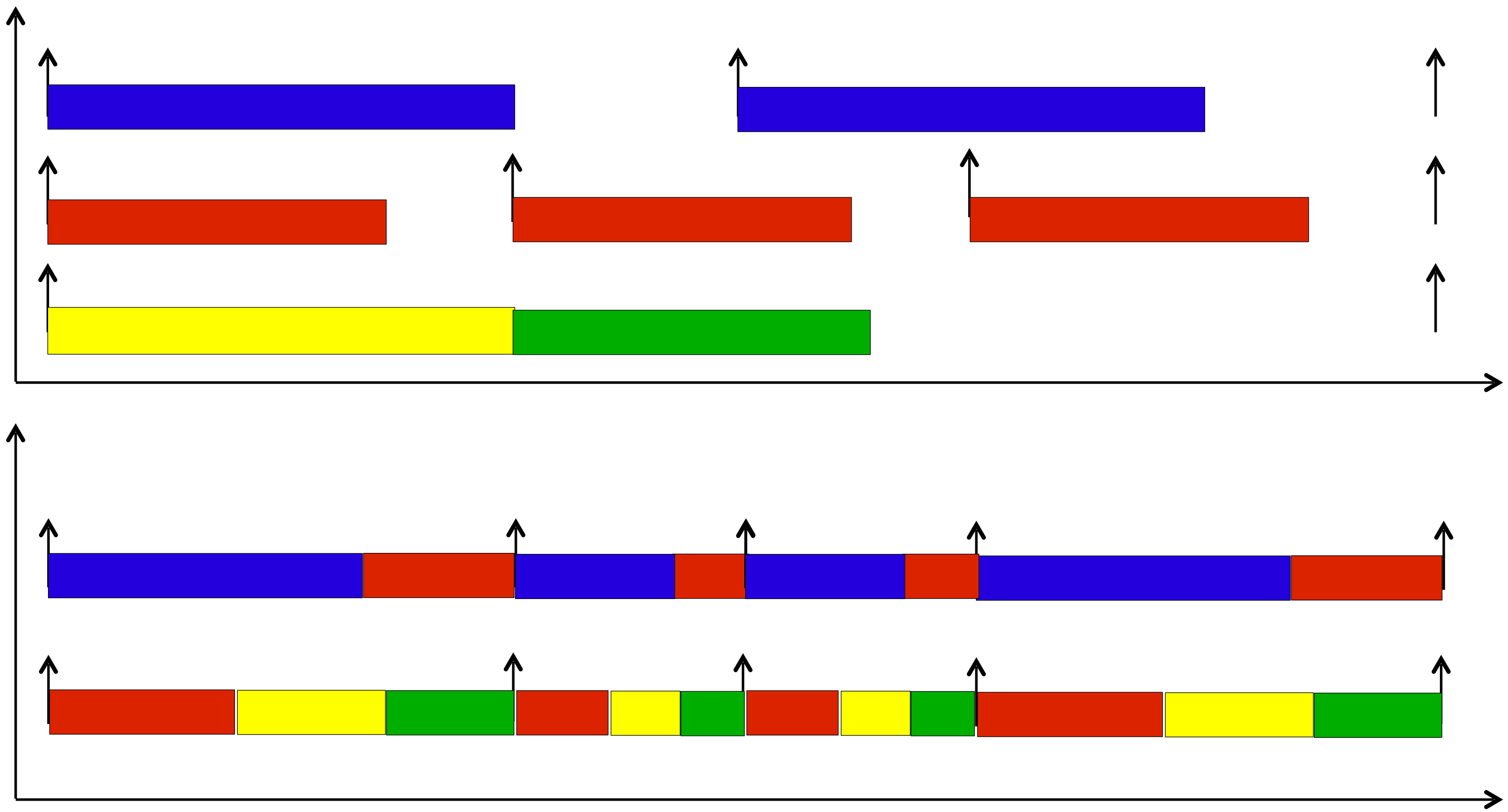
DP-Wrap – a DP-Fair Scheduler



wrap around to obtain schedule for 2nd CPU, ...

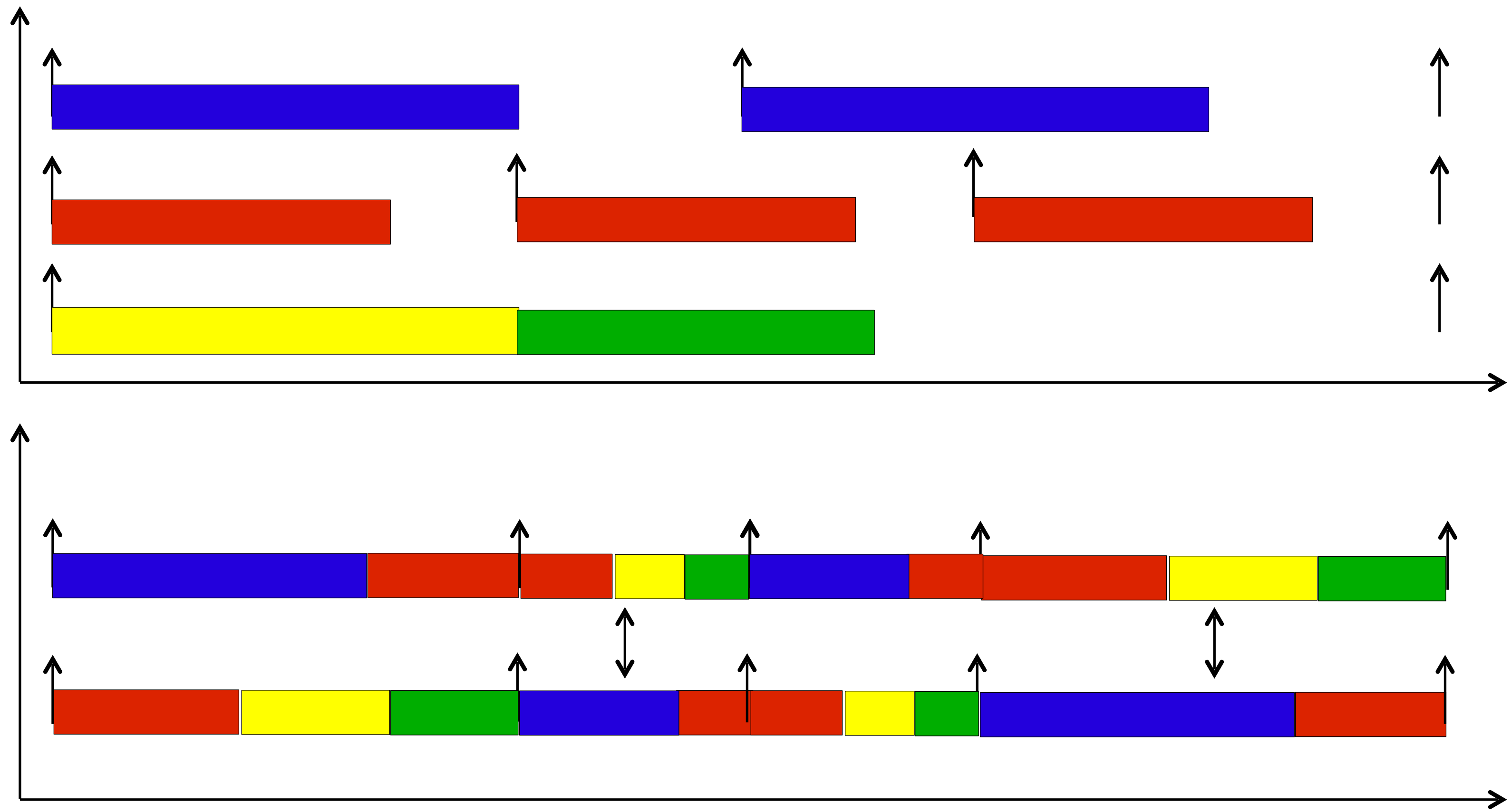
m – 1	migrations per chunk
n – 1	context switches per chunk

DP-Wrap – a DP-Fair Scheduler



Unnecessary migration of red task at chunk boundaries
=> mirror processor assignment of every second chunk

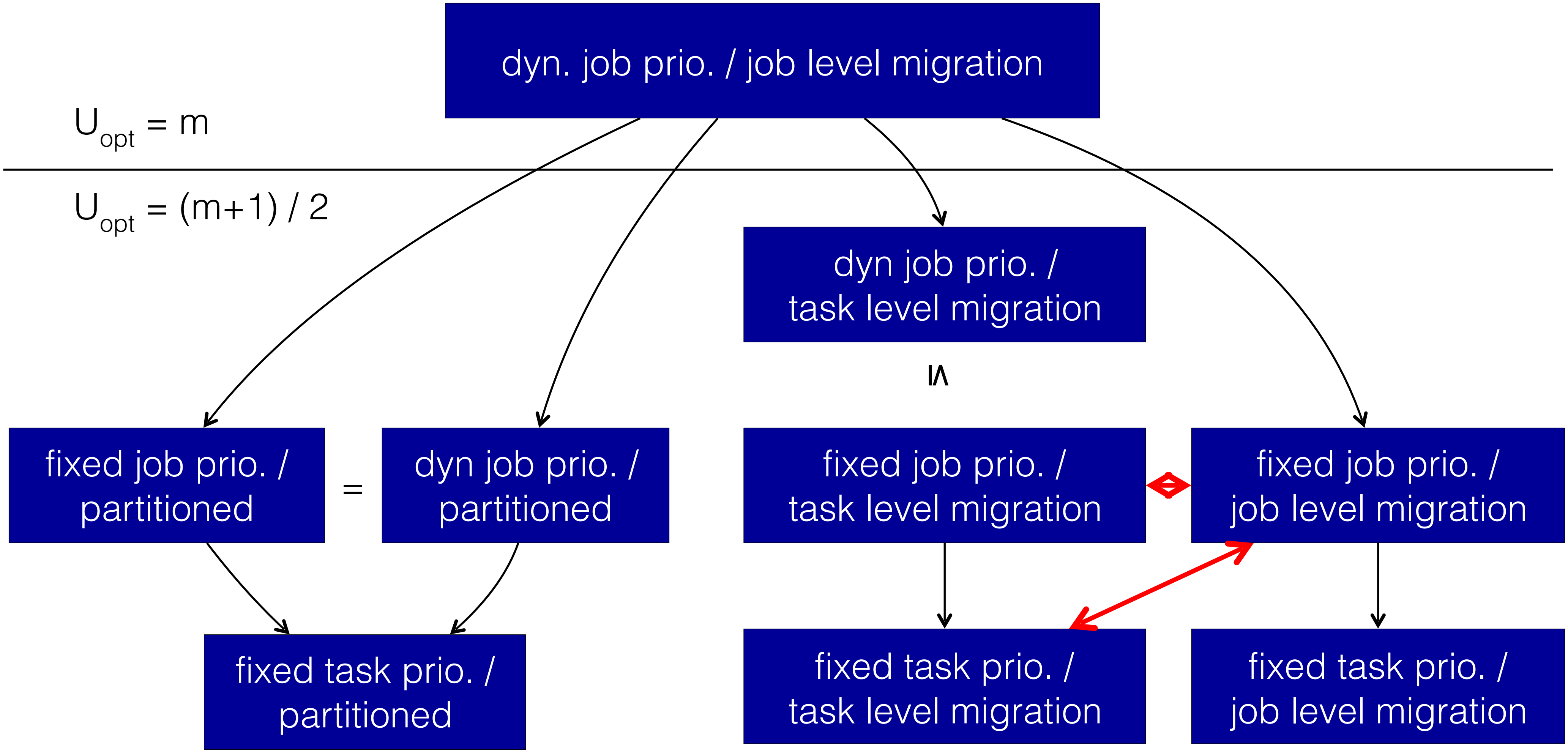
DP-Wrap (mirrored)



Design Space of MP Scheduling

dyn job prio. / partitioned	dyn job prio. / task level migration	dyn. job prio. / job level migration
fixed job prio. / partitioned	fixed job prio. / task level migration	fixed job prio. / job level migration
fixed task prio. / partitioned	fixed task prio. / task level migration	fixed task prio. / job level migration

Design Space of MP Scheduling



A → B => A can schedule any taskset that B can schedule and more

A ↔ B => dominance is not yet known

Outline

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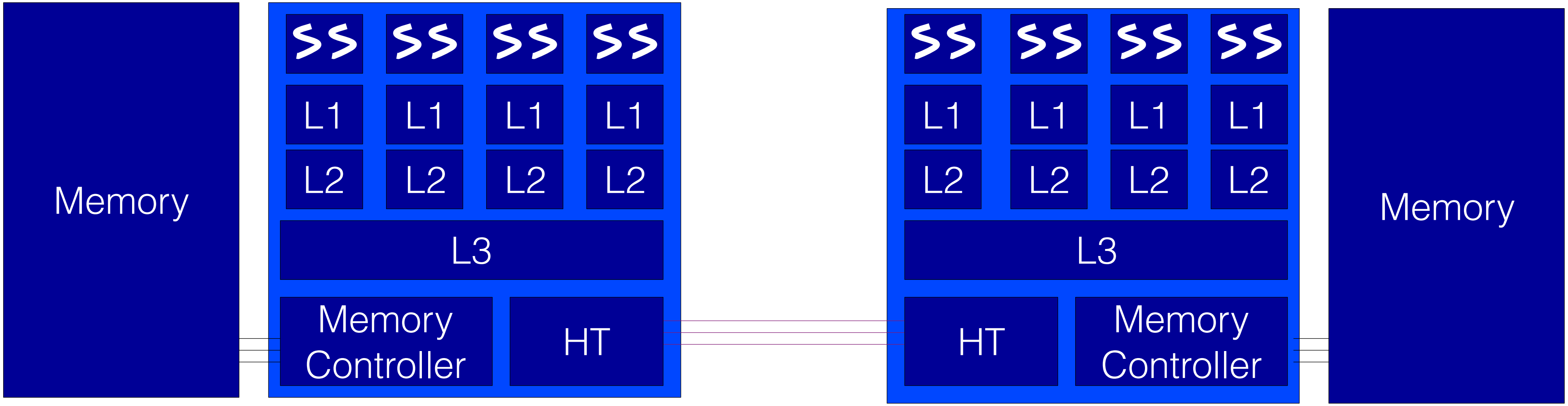
- Terminology, Notation and Assumptions
- Anomalies + Impossibility Results
- Partitioned Scheduling
- Global (Task-Lvl migration) Scheduling
    - G-FTP (e.g., G-RMS)
    - G-EDF
- Optimal MP Scheduling
- MP – Resource Access Protocols
- Open Research Issues

MP Resource Access Protocols

- UP:
  - Basic Priority Ceiling Protocol BPCP
  - Stack Resource Protocol SRP (Ceiling Priority Protocol CPP)
    - bounded priority inversion: | CS |
    - BPCP does not influence unrelated threads
- General Idea:
  - run UP protocol on every CPU of MP system
- Ceiling Priority i of Resource R_i: $\hat{R}_i$ = max prio(t_j)
  - here, priorities have a global meaning
- System Ceiling $\hat{S}$ = max $\hat{R}_i$ of held resources
- Synchronization Processor: CPU on which R_i is executed

Locking for Clustered Scheduling

- [Brandenburg '11]:
  - clustered scheduling: global within the cluster; partitioned in between



- Idea:
  - Every task helps out resource holders for a bounded time
  - Only the n-highest prioritized threads may acquire resources

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Open Issues [Burns '09]

- Limited processor Utilization
  - minimally dynamic algorithms
  - novel partitioning approaches
    - increase the guaranteed processing capability; overheads
- Ineffective Schedulability Tests (in particular, sporadic workloads)
  - large gap between feasibility / infeasibility tests
  - identify finite collection of worst-case job arrival sequences
- Overheads
  - migration costs; run queue manipulations; context switching
  - algorithms that permit intra-cluster migration; task-level migr.
- **Task Models**
  - intra-task parallelism, runtime integration
  - heterogeneous resources, Turbo Boost, GPUs

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