

DISTRIBUTED OPERATING SYSTEMS

- Name no more precise →
 Interesting/advanced Topics in Operating Systems
 - scalability
 - systems security
 - modeling
- Some overlap with "Distributed Systems" (Prof Schill) and some classes by Prof Fetzer
- In some cases no written material (except slides)



DISTRIBUTED OPERATING SYSTEMS

- Several lectures presented by research-group members.
- Strongly suggested: register for mailing list!!!
 - for questions and discussions
 - only way to inform about short-term issues
 - use: "tu-dresden" mail-adresses





- oral
- about 1 exam date per month until September (except August)
- Diplom/Master INF study programmes:
 can be combined with other classes in complex modules
- for application:
 - email to angela.spehr@tu-dresden.de for exam date/time
 - provide paperwork (forms) at least 2 weeks before exam
 (after that no more cancellation except for sicknesss)



Faculty of Computer Science Institute of Systems Architecture, Operating Systems Group

SCALABILITY IN COMPUTER SYSTEMS EXAMPLE: DNS/BIND

HERMANN HÄRTIG, DISTRIBUTED OPERATING SYSTEMS, SS2019



GOAL OF LECTURES ON SCALABILITY

Outline:

- scalability: terminology, problems
- basic approaches
- case studies

Goal:

 understand some of the important principles how to build scalable systems



OUTLINE AND GOAL OF TODAY'S LECTURE

Outline:

- scalability ...
- names in Distributed Systems:
 purposes of naming, terminology
- application of scalability approaches on name resolution

Goal:

 understand some of the important principles how to build scalable systems ... using DNS as example



MORE CASE STUDIES IN THE CLASS

- memory consistency
- locks and advanced synchronization approaches
- file systems
- load balancing (MosiX) and HPC



GENERAL DEFINITION: SCALABILITY

Scalability:

the ease with which a system or component can be modified to fit the problem area http://www.sei.cmu.edu/str/indexes/glossary/

Dimensions of Scalability:

- resources: CPUs, memory
- software (versions, better libs, etc.)
- heterogeneity (different hardware / SW = portability)



SPECIFIC: SCALABILITY IN SYSTEMS

 A system is described as scalable if it remains effective when there is a significant increase in the number of resources and the number of users.
 (Coulouris, Dollimore, Kindberg: Distributed Systems)

 Scalability [in telecommunication and software engineering] indicates the capability of a system to increase performance under an increased load when resources (typically hardware) are added (Wikipedia)



A SW ENGINEERING ASPECT OF SCALABILITY

Prepare for change in functionality

- software engineering
- choose sufficiently large logical resources
- provide hooks for extension

Not subject of the course



- performance bottlenecks / Amdahl's Law
- failures / abuse
- administration



AMDAHL'S LAW

Speedup: original execution time enhanced execution time

- P: section that can be parallelized
- 1-P: serial section
- N: number of CPUs

Speedup(P,N) =
$$\frac{1}{1-P+\frac{P}{N}}$$

■ if N becomes VERY large, speedup approaches: 1/(1-P)



THE "RPC" PRINCIPLES

partitioning

 split systems into parts that can operate independently to a large extent

replication

- provide several copies of components
 - that are kept consistent eventually
 - that can be used in case of failure of copies

locality (caching)

 maintain a copy of information that is nearer, cheaper/faster to access than the original

MORE PRINCIPLES FOR SCALABILITY

identify and address bottlenecks (!!!)

- specialize functionality/interfaces
- right level of consistency
 caches, replicates, ... need not always be fully consistent
- lazy information dissemination
- balance load



MORE ON LOAD BALANCE

- balance load
 - keep load under reasonable threshold
 - at the processing components
 - in the communication subsystems
- load balancing can be static or dynamic.

 Will study a detailed example for dynamic load balancing later(MosiX).



SOME MORE ISSUES

minimize the delay induced by "RPC"

prepare for change

information dissemination



DOMAIN NAME SYSTEM AS AN EXAMPLE

- names and name resolution etc in general
- a bit of history of internet names
- DNS general properties
- RPC in DNS



NAMES, IDENTIFIERS, ADDRESSES

names

- symbolic
- have a meaning for people
- identifiers
 - identifies a component (uniquely)
 - are used by programs
- addresses
 - locates a component & can change
 - can change



NAME RESOLUTION

- name resolution:
 - map symbolic names to objects
 - in detail: to a set of attributes such as:
 identifiers, addresses, other names, security properties

- Principle interface:
 - Register (Name, attributes, ...)
 - Lookup (Name) -> attributes



- compilers
 - statically map names to addresses
- dynamic libraries
 - dynamically remap addresses
- port mapper (SUN RPC)
 - map service to port
- Name resolution is a form of dynamic mapping of pathnames to attributes.



OBSERVATION

- Many services, tools, ... provide their own name resolution
 - file systems (UNIX: path names to I-Nodes)
 - login
 - RPC (remote procedure call) systems (portmapper)



PURPOSE OF DIRECTORY SERVICES

- integration of name services
- generic name service
- world-wide use of names
- pervasively used:
 - email/web
 - computer attributes (IP addresses)
 - people attributes (certificates, ...)
 - **...**



A BIT OF HISTORY

- UUCP/MMDF (cum grano salis):
 - ira!gmdzi!oldenburg!heinrich!user (path to destination)
 - user@ira!heinrich%gmdzi (mixing identifiers and path information)

A BIT OF HISTORY

- ARPA-Net at the beginning:
 - a single file: hosts.txt
 - maintained at Network Information Center of SRI (Stanford)
 - accessed via ftp
 - TCP/IP in BSD Unix => chaos name collisions, consistency, load
- DNS: Paul Mockapetries (84) ...

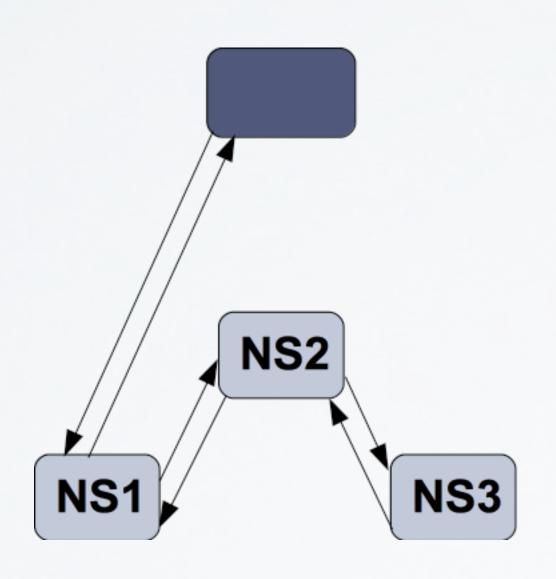


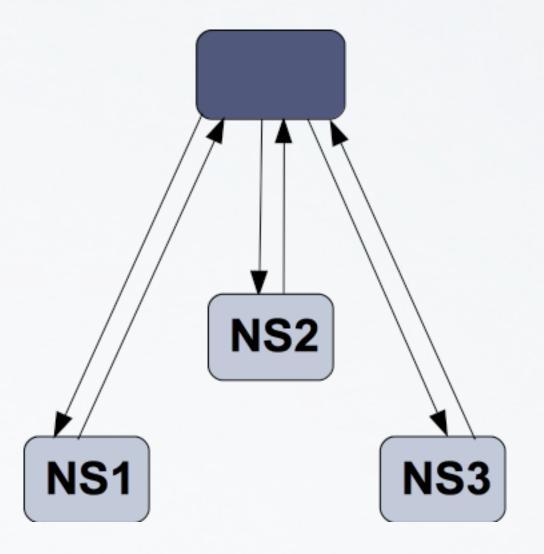
SOME DNS TERMINOLOGY

- naming domain
 subtree in the hierarchy of DNS contexts
- zone (aka Zone of authority) Subset of a domain over which an authority has complete control. Subzones (starting at apices of a zone) can be delegated to other authorities.
- navigation
 querying in a set of cooperating name spaces



RECURSIVE./. ITERATIVE





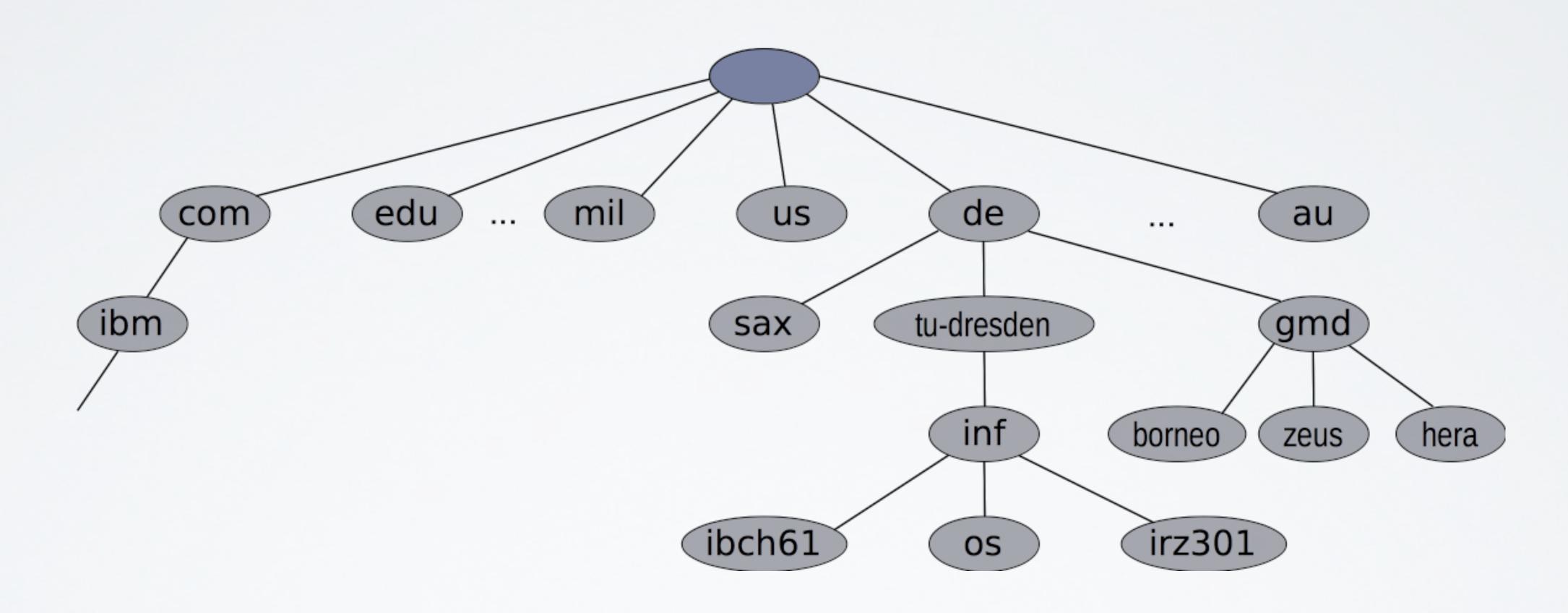


REQUIREMENTS / PROPERTIES

- arbitrarily large numbers
- arbitrary units of administration
- long living names, the higher in the hierarchy the longer
- high robustness
- restructuring of name spaces
- consistency
- efficiency



DNS NAME SPACE (ORIGINAL)



TODAY: hundreds of "top level domains"



inf.tu-dresden.de domain

os.inf.tu-dresden.de computer

heidelberg.ibm.com domain

ftp ftp.inf.tu-dresden.de

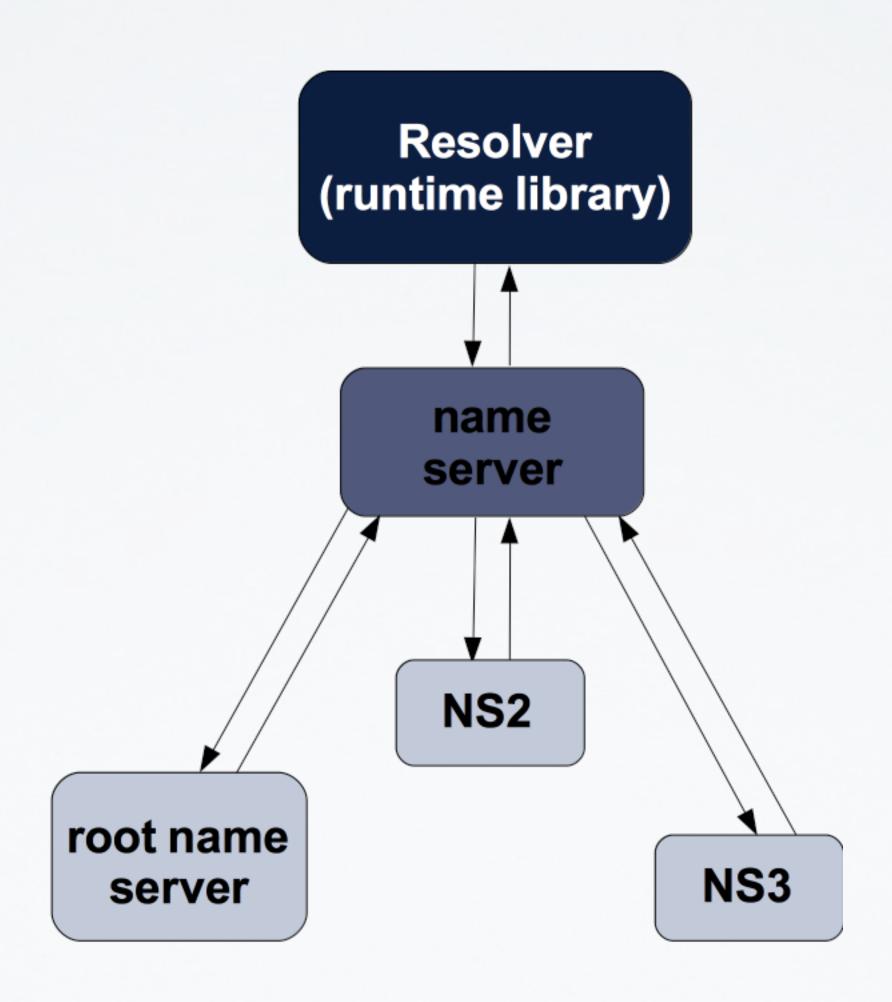
■ DNS: → IP address: 141.76.2.3

■ ftp daemon: IP address, port 21

properties:
 location independent / not very deeply nested



IMPLEMENTATION STRUCTURE (BIND)

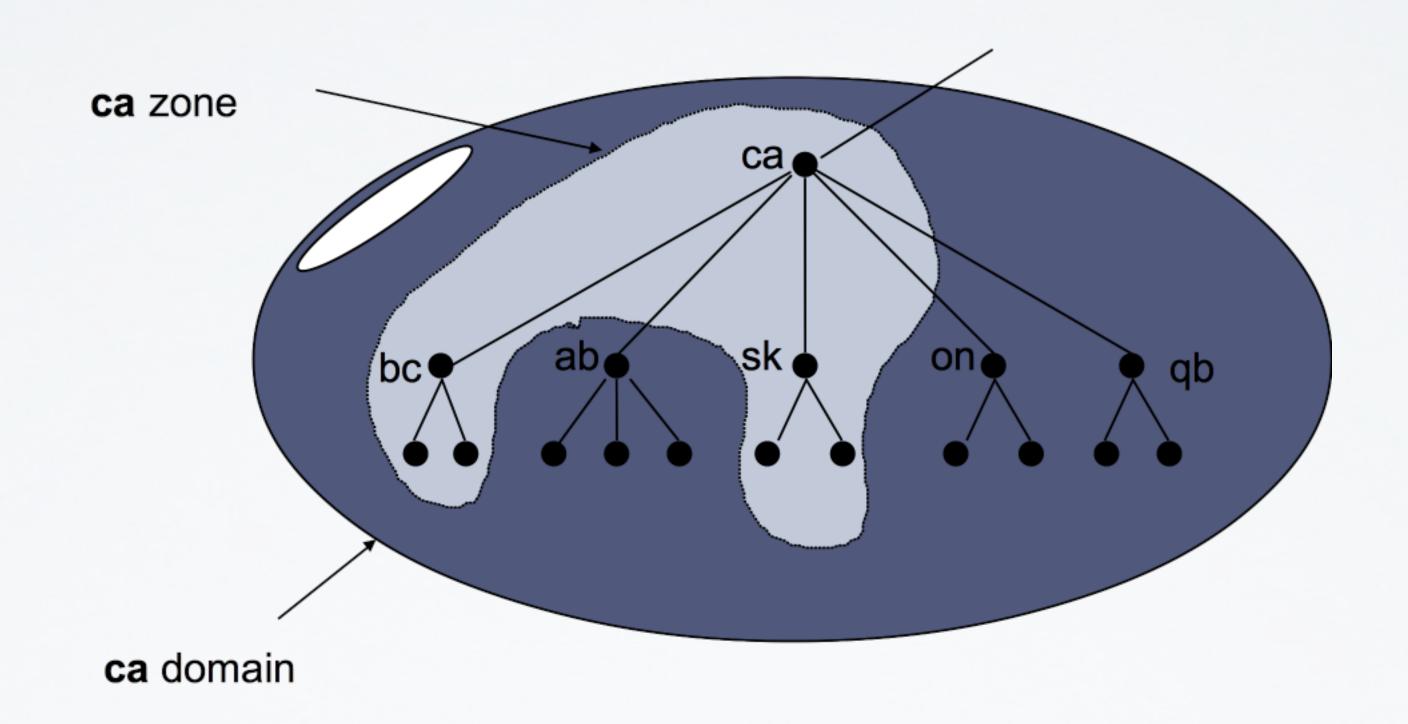






- Zones:
 - administrative unit
- Name Server:
 - maps to names and addresses of name servers responsible for sub zones
 - maintains management data
 - process doing the name resolution for one zone
- key interface: Resource records (RR)





example taken from Coulouris et al, Distributed Systems



RPC: REPLICATION

- 2 ways of replication:
 - several IPs/names
 - "any cast" (send packet to one of many servers with same IP)
- 13 root name server IPs, several hundreds of any cast
- each zone has at least
 one primary and one secondary IP



- each name server caches resource records
- time to live attribute
- authoritative versus non-authoritative answers



RESOURCE RECORDS

Record type	Interpretation	Content
Α	address	IPv4 address
AAAA	address	IPv6 address
NS	Name server	DNS name
CNAME	Symbolic link	DNS name of canonicial name
SOA	Start of authority	Zone-specific properties
PTR	IP reverse pointer	DNS name
HINFO	Host info	Text description of host OS
•••		



- Paul Albitz & Cricket Liu
 DNS and BIND
 O'Reilly & Associates, Inc.
- Mark Hill, Michael Marty
 Amdahl's Law in the Multicore Era IEEE
- Couluris, Tollimore, Kindberg
 Distributed systems