

INVESTIGATING THE LIMITATIONS OF PVF FOR REALISTIC PROGRAM VULNERABILITY ASSESSMENT

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Motivation

Software development today:

- How fast does my program run?
- How much energy does it consume?
- How many software errors does it contain?

Software development in the future:

- Which functions are most vulnerable against certain hardware errors?
- Should I use implementation A or B of a certain algorithm?
- Which impact will compiler optimizations have on my program's vulnerability?



Vulnerability Analysis Today

Hardware vulnerability analysis: Architectural Vulnerability Factor (AVF)¹

- Established technique
- Requires both application and hardware model
- Hardware model out of reach in practical SW development

¹ Mukherjee et al.: A systematic methodology to compute the architectural vulnerability factors for a high-performance microprocessor, IEEE MICRO 2003

² Sridharan, Kaeli: Eliminating microarchitectural dependency from architectural vulnerability, HPCA 2009
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PVF Limitations
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Vulnerability Analysis Today

Hardware vulnerability analysis: Architectural Vulnerability Factor (AVF)¹

- Established technique
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Software-only analysis: Program Vulnerability Factor (PVF)²

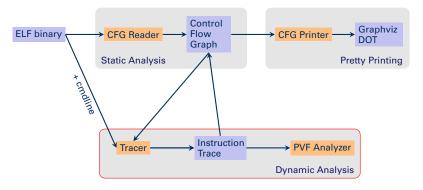
- · Replace hardware with an abstract model
- Requires application knowledge only

¹ Mukherjee et al.: A systematic methodology to compute the architectural vulnerability factors for a high-performance microprocessor. IEEE MICRO 2003

² Sridharan, Kaeli: Eliminating microarchitectural dependency from architectural vulnerability, HPCA 2009
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PVF Analysis for x86



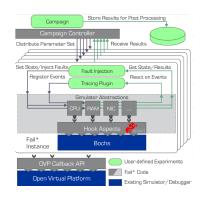
Source code: https://github.com/TUD-OS/PVFAnalyzer

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FAIL/*: Fault Injection

- Framework for performing FI campaigns³
- Interchangeable hardware model: Bochs/x86
- Benefit: real distribution of errors for a given platform



³ H. Schirmeier, M. Hoffmann, R. Kapitza, D. Lohmann, and O. Spinczyk: FAIL/*: Towards a versatile fault-injection experiment framework. ARCS 2012

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Experiment Setup

Application:

- nanojpeg JPEG decoder
- Decode a small JPEG image into bitmap
- GCC 4.4.5, fully optimized → 3.7 million instructions

Fault Injection:

- FI campaign using FAIL/*
- Flip every bit in every used register once and record outcome
- Runtime: about a week on fairly large cluster



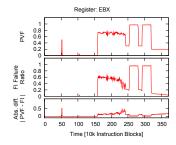
PVF Analysis:

- Analysis of a single trace
- · Same fault model as for FI
- Runtime: a couple of minutes on my laptop



Initial Results

- Program run split into blocks of 10.000 instructions
- For each block: compare computed PVF and fault ratio from FI experiments
- Question: What do we consider a fault?
 - Program crashes
 - Infinite loops
 - Silent data corruption

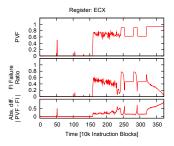


Observation #1: PVF is a fairly good prediction of the fault rate if we consider all SDC to be faults.



Understanding PVF-FI Deviations

- FI works on bit granularity, PVF analyses registers
- PVF does not analyze bit masking etc. operations
- PVF analysis stops on block boundaries – no propagation of fault state for more than 10,000 instructions





Corrupted data?









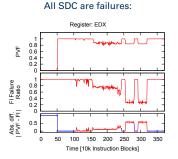
Observation #2: Not all data corruption can be considered equally faulty.

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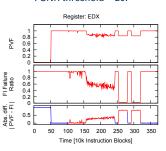


Considering Data Quality





PSNR threshold = 20:





Incorporating Quality Into PVF?

- Similar results for an analysis of H.264 video decoding: control data (high fault rate) vs. payload data (low fault rate)⁴
- Compiler extension:
 - reliable and unreliable data modifiers
 - Special protection (encoding / redundancy) for reliable data

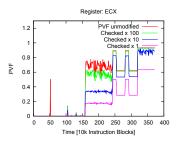
Observation #3: Using reliability annotations, the compiler can determine the reliability level of a hardware resource and propagate this information to a PVF analyzer.

⁴ A. Heinig, M. Engel, F. Schmoll, and P. Marwedel: *Improving transient memory fault resilience of an H.264 decoder*, STMEDIA 2010 Berlin, 21.01.2013 PVF Limitations slide 11 of 13



Modeling the Effect of FT methods

- PVF(app) models probability of a fault affecting execution
- Replicated execution (e.g., RMT): tolerates faults \(\frac{n}{2} 1\) independent faults
 PVF_{rep}(app) := PVF(app)\(\frac{n}{2}\)
- Encoded processing:
 - Checksummed data, periodic checking
 - Limit PVF analysis to smaller periods





Now What?

PVF seems promising, but is not there yet.

- We need to incorporate data quality information.
- Compiler support: propagate resource vulnerability information to PVF analyzer.
- Analysis of different fault models
 - Original PVF: register and ALU errors
 - How about memory bit flips, decoding errors, ...?

PVF will never fully replace full-featured FI campaigns.

- Analysis of complex applications and fault models
- Strong limitation: Abstract hardware model