

# **Tradeoffs for Virtualization Mechanisms**

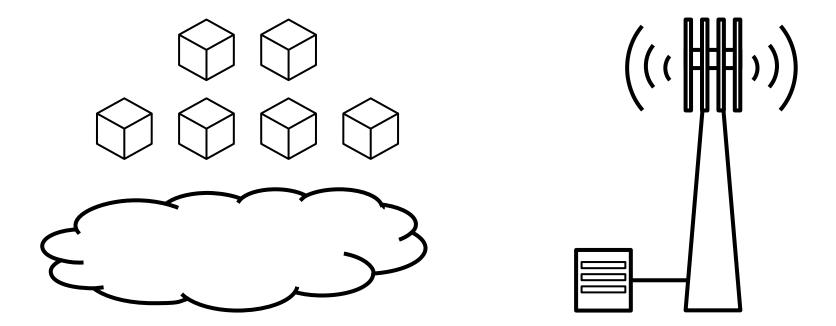
**Lectures on Distributed Operating Systems (SS'25)** 

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### **Recap: Public Clouds**



Large-scale systems that run applications on behalf of untrusted users



Outsourcing to cloud providers often due to (alleged) economic benefits

### **Requirements in Cloud Computing**



- Users want their applications to run reliably.
  - Data protection (confidentiality / integrity), availability
  - Do not depend on third parties in any form
- Users want performance close to that of native deployment.
  - Get resources that you pay for

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- Cloud providers want to maximize resource utilization.
  - If possible, customers should share physical hardware
  - Function as a service (FaaS) / serverless instead of static instances
  - Requires highly dynamic systems

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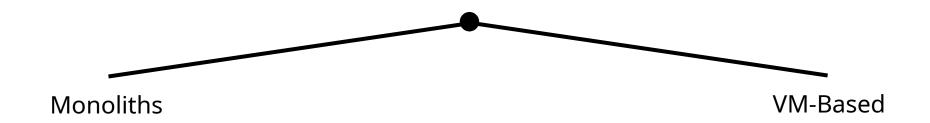




# **Isolating Cloud Applications – A Naive Approach**

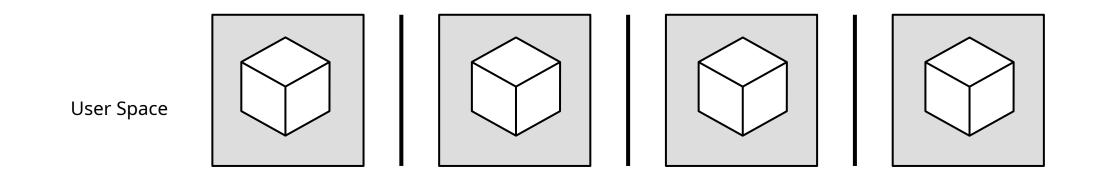
# **Hardening Operating Systems for the Cloud**







Dominated by Linux, a monolithic OS design



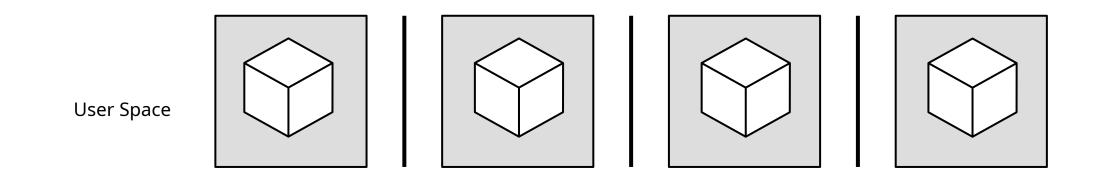
Kernel Space

FS Net Sched Memory IPC

Drivers



Monolithic OS: large API, large trusted computing base (TCB)



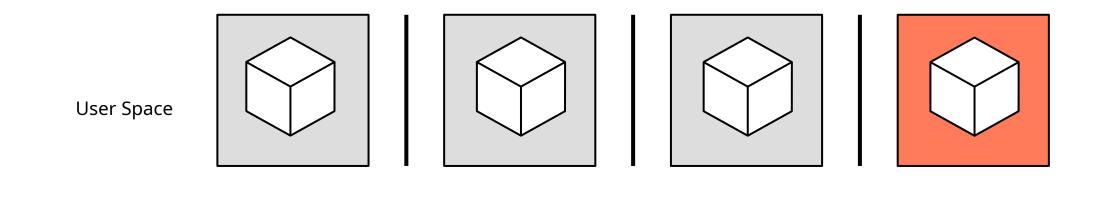
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Monolithic architectures are prone to privilege escalation



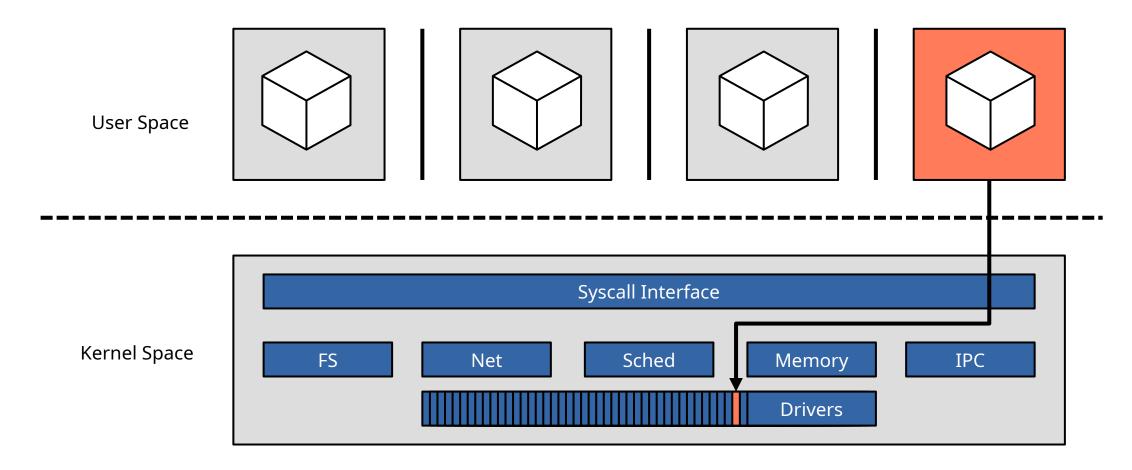
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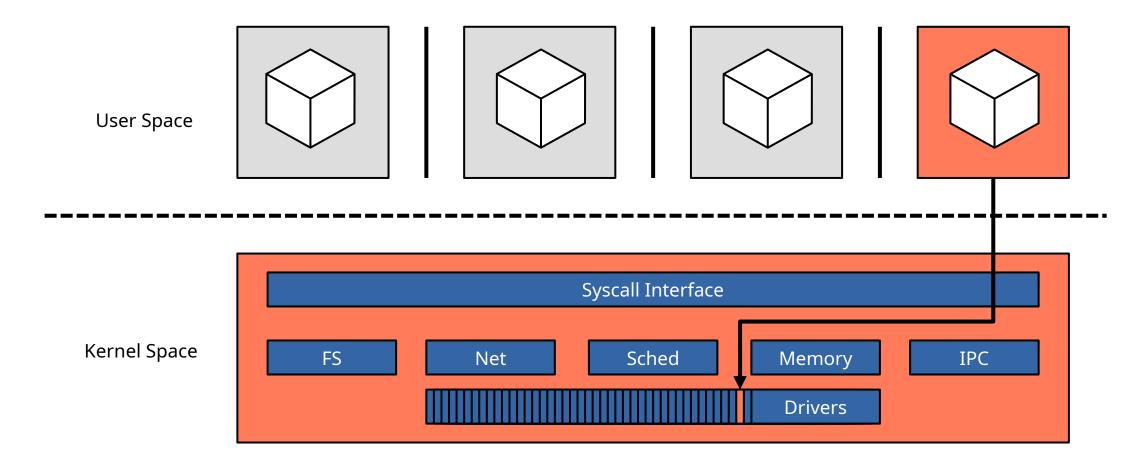


Monolithic architectures are prone to privilege escalation



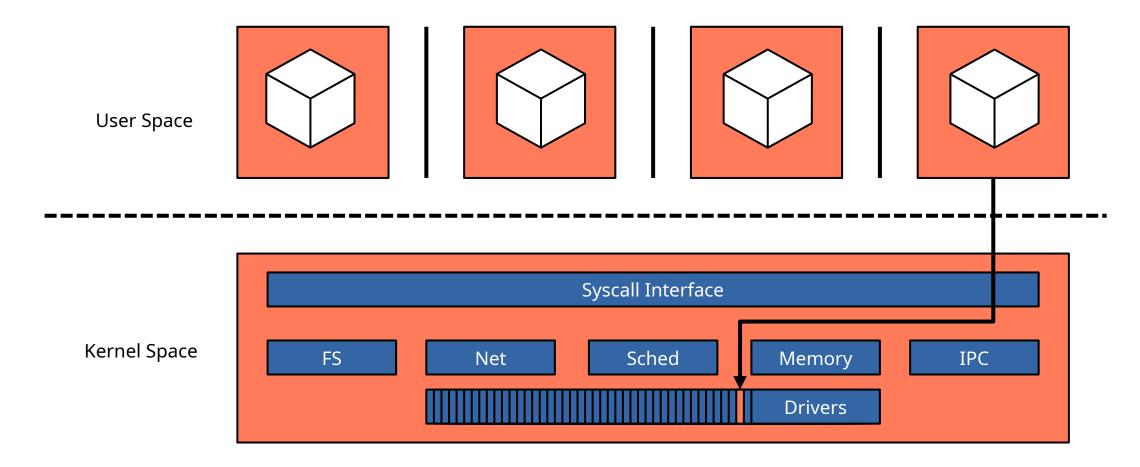


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Monolithic architectures are prone to privilege escalation



### Why Is Process-Based Isolation on a Monolith Not Optimal?



- Architectural weaknesses
  - No restriction of accessible kernel API (ambient authority)
- Insufficient means for enforcing resource limits
  - If applicable, rlimit too coarse-grained
- Processes share several global "namespaces"
  - E.g., process IDs, network devices, view on file system, ...

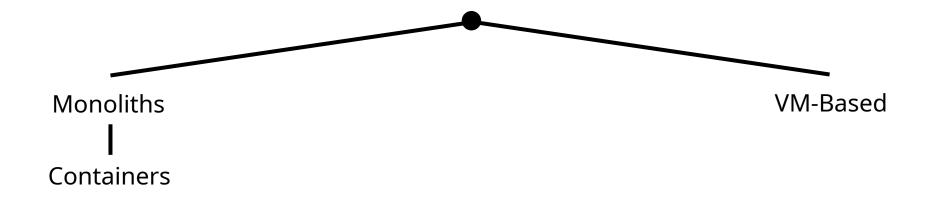
But: Excellent performance on standard hardware, compatible, easy to use



# **Isolating Cloud Workloads in Containers**

# **Hardening Operating Systems for The Cloud**

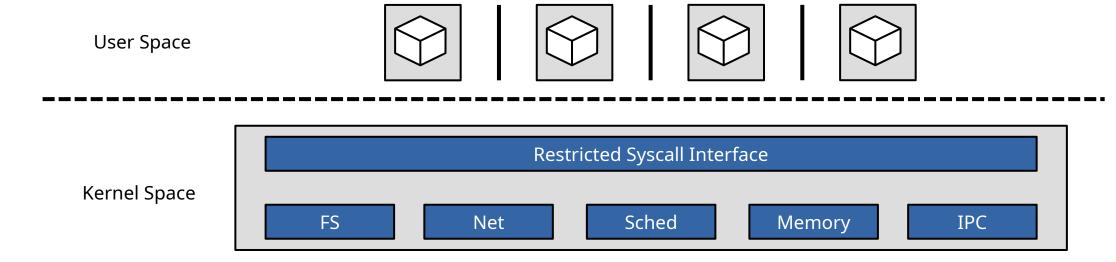




#### **Containers**



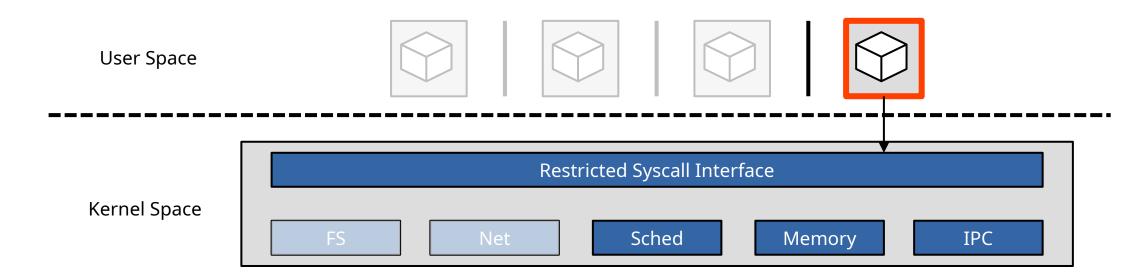
- Lightweight virtualization (BSD Jails [KW00], Solaris Zones [TC04], Linux [FF+15])
- Increase isolation level of groups of processes using in-kernel facilities
  - Comes with an entire ecosystem for deployment
  - Intuition: Provide separated userland instances on top of a shared kernel
  - No application modification and / or awareness needed, run images as-is



### **Linux Containers - Namespaces**



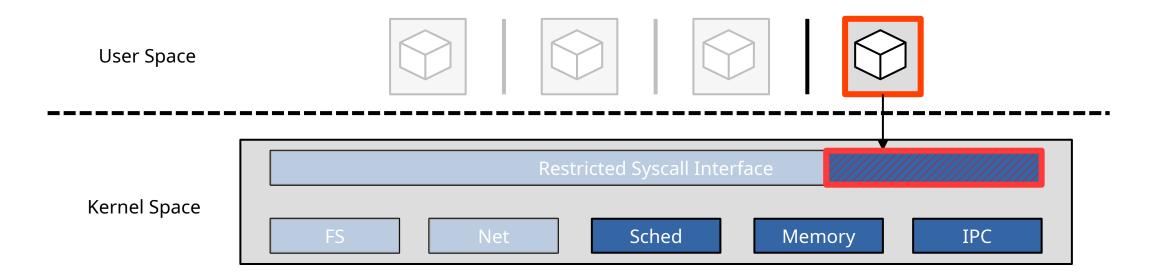
- Restrict visibility of system components inside a container
- Used to virtualize a variety of kernel abstractions shared by stock processes
  - PIDs, mount points, network interfaces, IPC channels, ...
- Often additional restrictions for visibility of file system subtrees (chroot, ...)



### **Linux Containers – Seccomp-BPF**



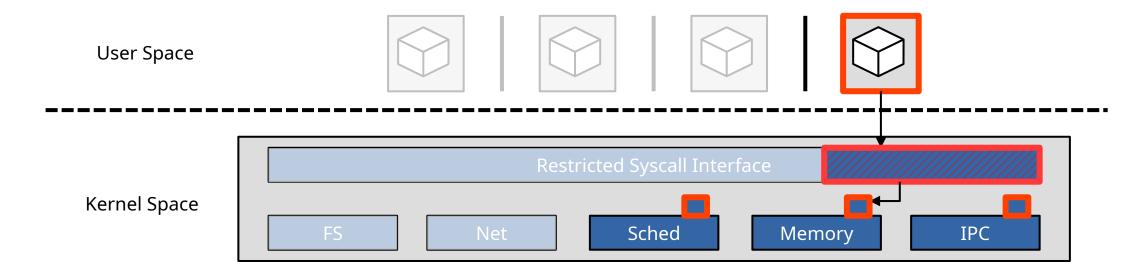
- Enables administrator to limit execution of system calls
  - Inject BPF programs interpreted at certain probe points in the kernel
  - Possible to load / unload BPF programs at runtime
  - Also enables restriction of valid parameter range for system calls



### **Linux Containers – Cgroups**



- Unified framework for resource restriction and accounting
  - CPU time / share, memory budgets, network bandwidth, ...
  - Hierarchical organization possible → Pass fraction of own budget to other process
  - Offers priorities (e.g., differentiate between different customers)



#### **Containers – Additional Benefits**



- Execution inside virtual machines for performance reasons [SS+19]
  - Deploy customized system services per containerized application
- Implementation of Linux features in userspace container runtime [YZ+19]
  - Idea: Reduction of API used by the container
  - gVisor widely applied by Google
- Adding and removing container parts on demand [TB+18]
  - Reduce TCB inside containers



# **Hardening Containers**

### **Everything Secure Now?**



- Various approaches for further hardening container mechanisms
  - Automatic generation of system call filters [CW+21]
  - Transparent change of kernel substrate the containers run on [NS+23]

- ...

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- ...

- The complexity of container mechanisms remains.
  - Various CVEs (e.g., CVE-2018-18955 for namespaces, CVE-2022-0492 for cgroups)
  - Executing user-provided code in kernel dangerous [HG+23]
  - Huge TCB for confidentiality, integrity, and availability



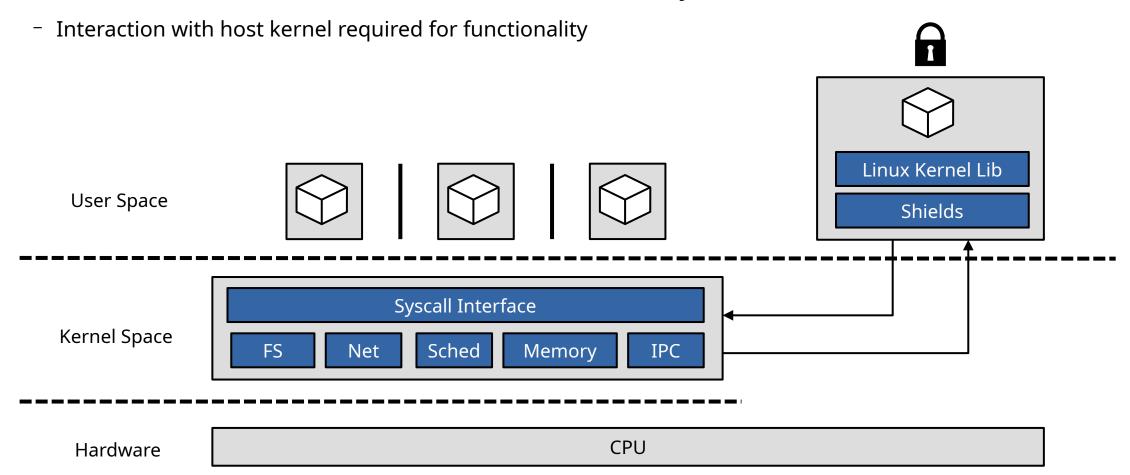
25

- SCONE: Trusted execution environment (TEE) to shield container [AT+16]
  - Based on Intel SGX
  - Idea: Let hardware provide secure compartments for executing applications
- SGX provides extensive security guarantees (→ see lecture on trusted execution)
  - Enclaves provide encrypted execution of applications in ring 3 (user space)
  - Enclave set up by OS, inaccessible from any software after initialization (including SMM and DMA)
  - OS, BIOS, firmware are *not* a part of a container's TCB
  - Only need to trust CPU vendor (for attestation)
- TEEs like SGX can incur significant performance overhead



26

Container in TEE runs isolated from the rest of the system



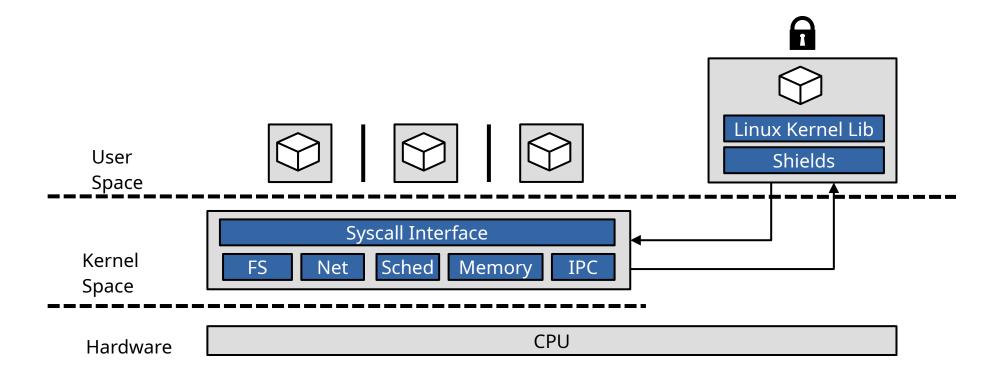


27

- SCONE applies optimizations to reduce SGX overhead
- Async communication interfaces - Implementation of kernel functions inside enclave SCONE only incurs moderate performance penalty (0.6 – 1.2x throughput) Linux Kernel Lib **User Space** Shields Syscall Interface Kernel Space Net Sched Memory IPC **CPU** Hardware

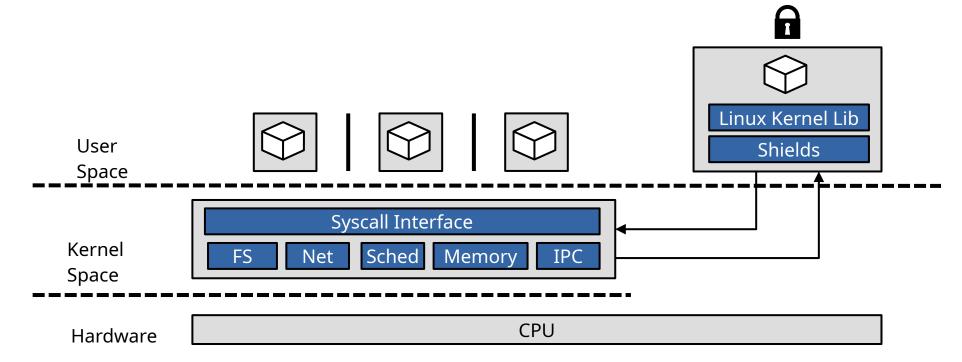


• What are the security implications of TEEs?





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  - Reliant on hardware vendor
  - Complexity moved from soft- to hardware
  - Effective TCB?





30

- What are the security implications of TEEs?
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# FORESHADOW: Extracting the Keys to the Intel SGX Kingdom with Transient Out-of-Order Execution

Jo Van Bulck, imec-DistriNet, KU Leuven; Marina Minkin, Technion; Ofir Weisse,
Daniel Genkin, and Baris Kasikci, University of Michigan; Frank Piessens, imec-DistriNet,
KU Leuven; Mark Silberstein, Technion; Thomas F. Wenisch, University of Michigan;
Yuval Yarom, University of Adelaide and Data61; Raoul Strackx, imec-DistriNet, KU Leuven

https://www.usenix.org/conference/usenixsecurity18/presentation/bulck

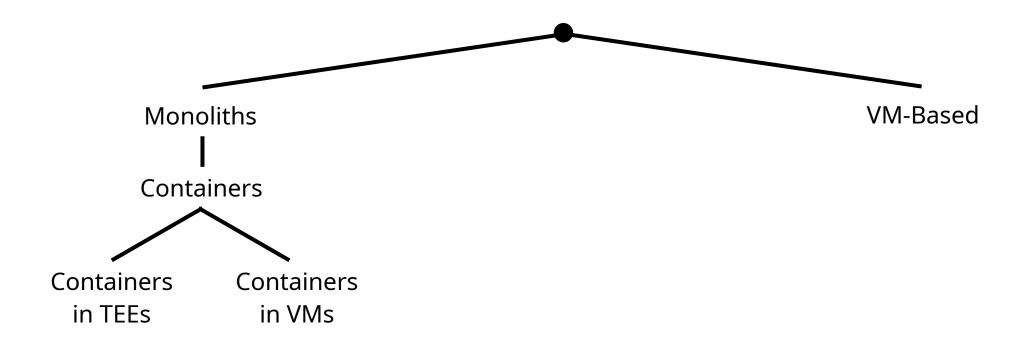
This paper is included in the Proceedings of the 27th USENIX Security Symposium.

August 15-17, 2018 • Baltimore, MD, USA

ISBN 978-1-939133-04-5

### **Hardening Operating Systems For The Cloud**





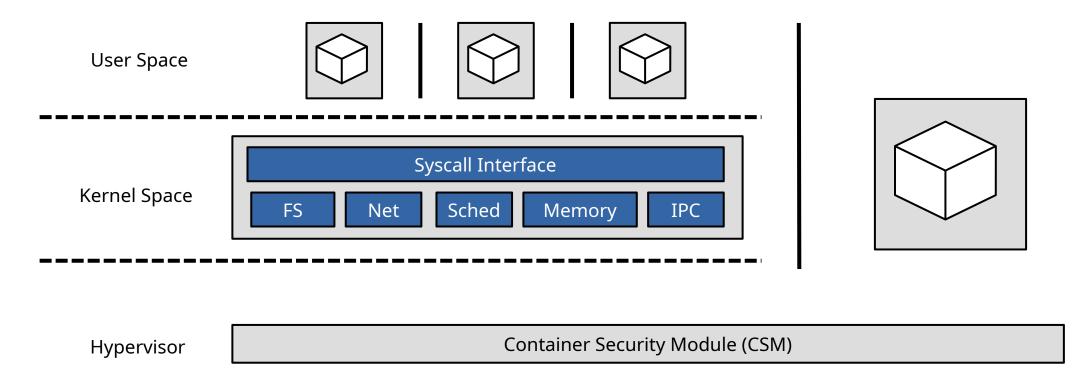


- Simpler hardware, yet strong isolation possible (BlackBox [VN22])
  - Runs on standard virtualization hardware

- Run containers on top of a container security module (CSM)
  - Linux VM with drivers etc. running beside application containers
  - CSM shields application containers
  - Applications should not have to trust function donor VM

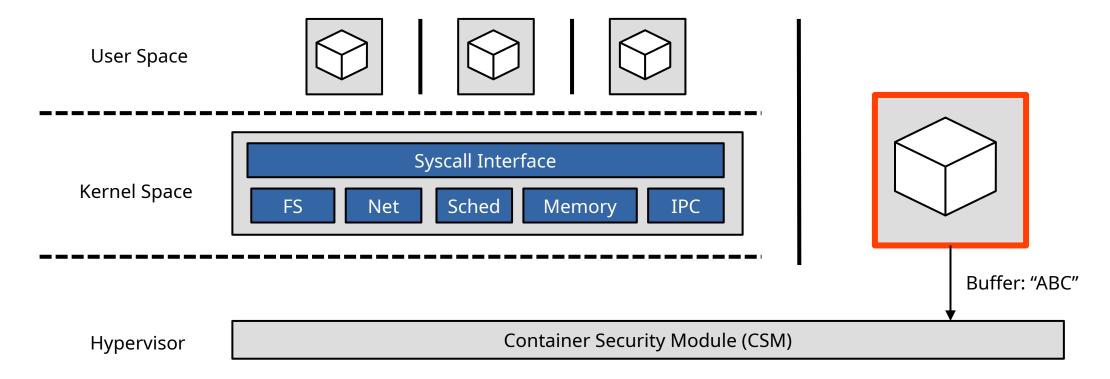


- Simple hardware, yet strong isolation possible (BlackBox [VN22])
  - Only needs to trust CSM for security, integrity (~10k LoC)
  - Moderate performance overhead compared to native containers



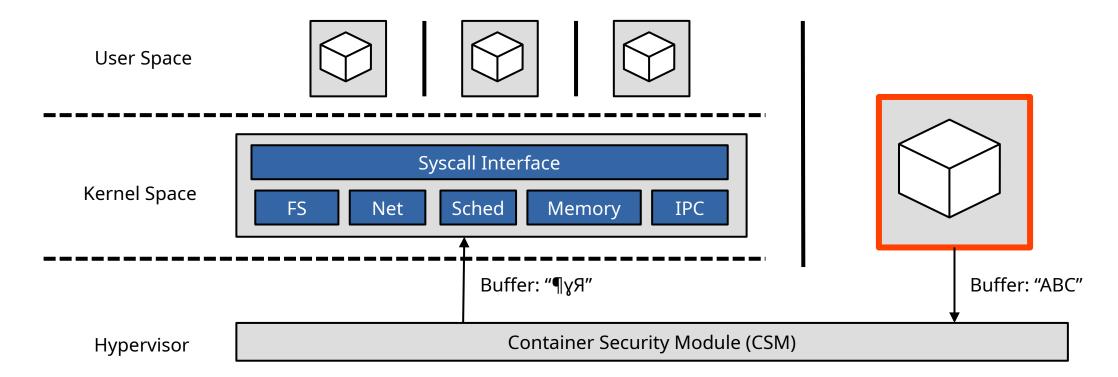


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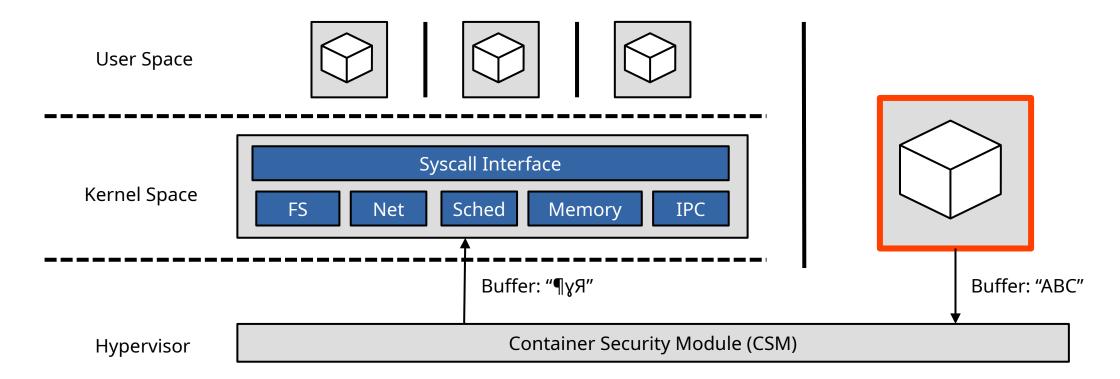
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36

- Simple hardware, yet strong isolation possible (BlackBox [VN22])
  - What about the TCB size regarding availability?





# **Lightweight Virtual Machines**

### **Virtual Machines in Dynamic Cloud Settings**



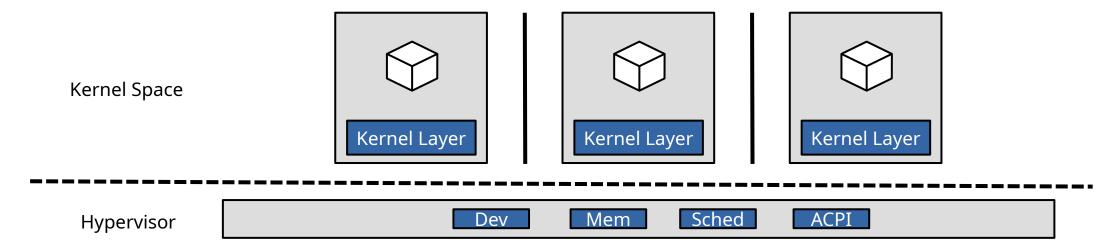
- Strong isolation between VMs
  - Hypervisor has a smaller interface to provide than an OS kernel
  - Hardware acceleration is wide-spread, also for many devices (SR-IOV [YY+08])
- Performance overhead cannot be avoided.
  - Nested paging
  - Device emulation

- Load complete kernel, startup considered slow [ML+17]
  - Partially caused by device initialization
  - Cost for being able to run arbitrary operating systems

## **Lightweight Virtual Machines**



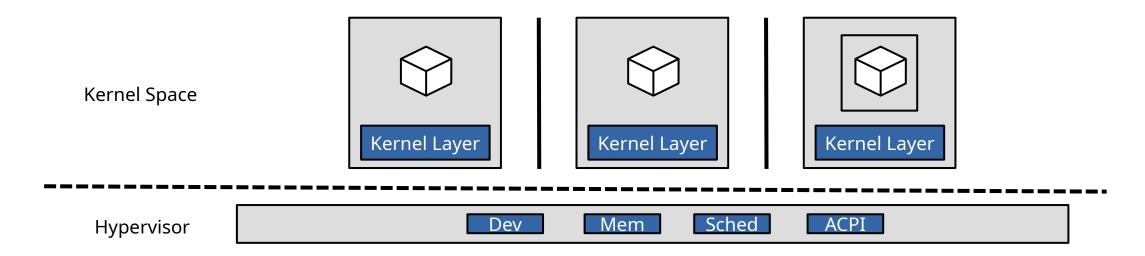
- Reduce VM startup overhead, but retain isolation (e.g. LightVM [ML+17])
  - Remove unused devices etc.
  - Boot time as low as a couple of milliseconds
- Dedicated hypervisors offering tight feature set
  - E.g., Firecracker [AB+20] (used by AWS)



## **Lightweight Virtual Machines**



- Flavor I: Unikernel guests (e.g., Unikraft [KB+21])
  - Performance and TCB advantages possible
- Flavor II: Containers inside lightweight VM (e.g., Kata containers [KC24])
  - Second layer of defense, used with Alibaba Cloud [LC+22]

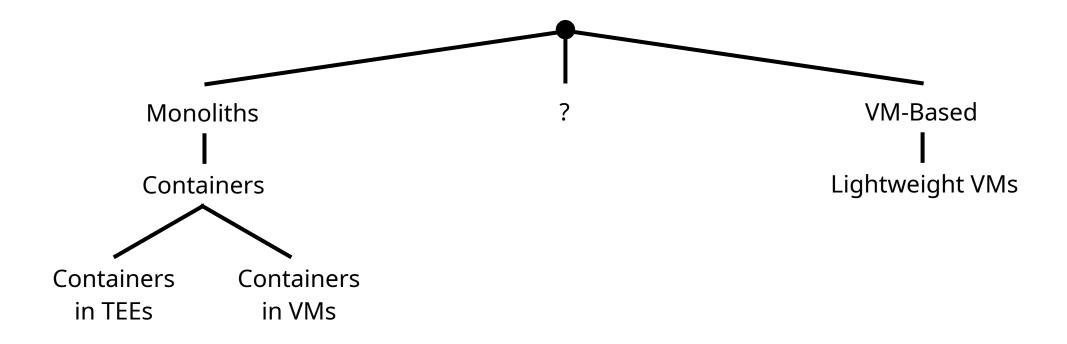




### **Current Research: Microkernels as Data Center Substrate**

## **Hardening Operating Systems for the Cloud**







43

- General idea of microkernels: Minimize code running in kernel mode
  - Drivers, file systems, network etc. implemented in user space
  - Only requires basic CPU architecture (standard virtual memory and threads)
- Modern microkernels use capabilities for access control
  - Processes can only access resources for that they have been granted permission explicitly
  - No ambient authority
- Inter process communication (IPC) important for performance
  - Microkernels by design have performance disadvantages compared to monoliths

**Till Miemietz**, Viktor Reusch, Matthias Hille, Max Kurze, Adam Lackorzynski, Michael Roitzsch, Hermann Härtig: A Perfect Fit? - Towards Containers on Microkernels. WOC@Middleware 2024



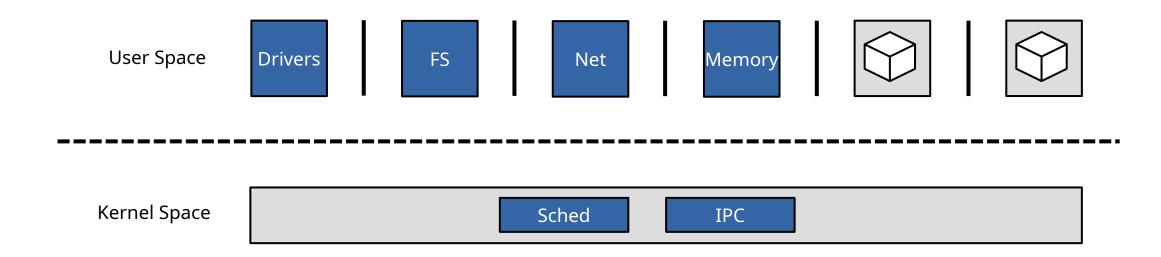
- What if we used an OS with proper process isolation to begin with?
  - Capability-based access control
  - Small TCB for confidentiality, integrity, and availability

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45

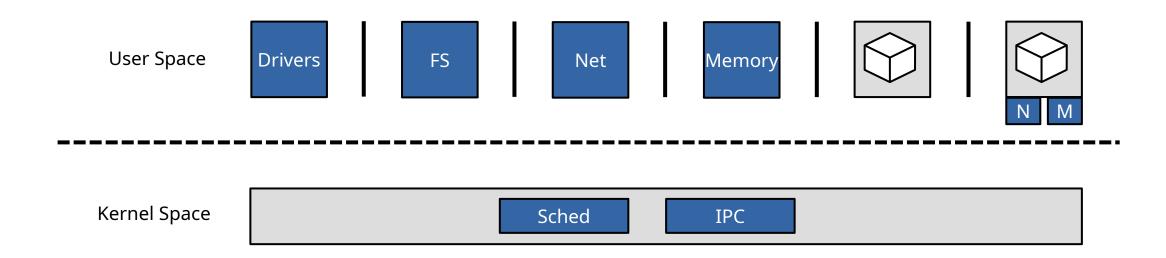
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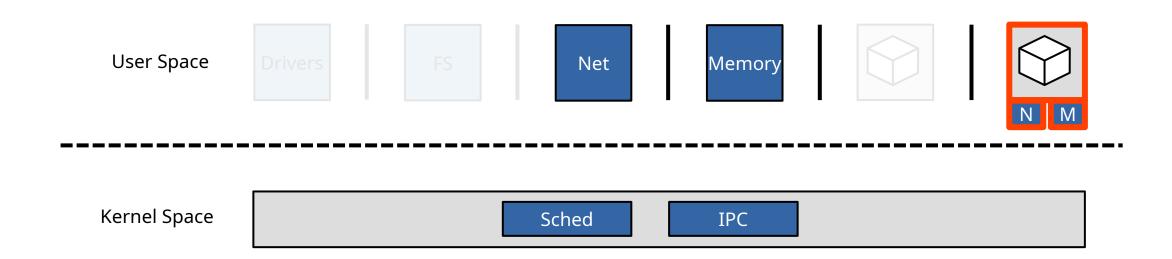


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47

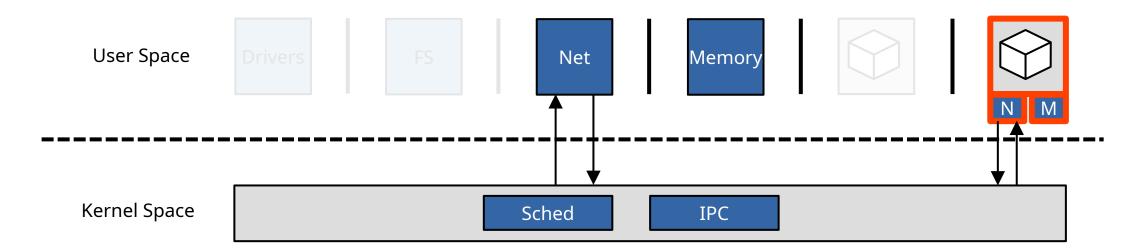
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- What if we used an OS with proper process isolation to begin with?
  - Capability-based access control
  - Small TCB for confidentiality, integrity, **and** availability
  - Interaction with system services more expensive



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- Visibility restrictions
  - Integrated in capability system
  - Namespace abstraction also exists on L4Re
- Interface restrictions: Containers can only access a minimal interface by design
  - There is no ambient authority.
- Resource restrictions
  - Tied to capability
  - To be implemented by individual service

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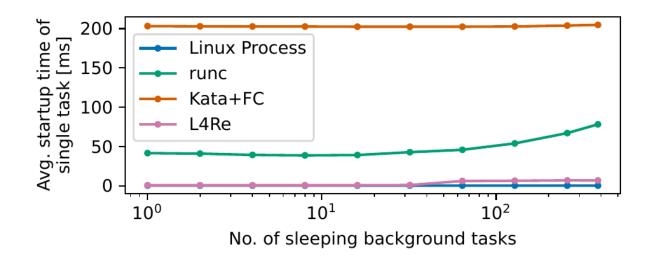
- Simplicity reduces attack surface by orders of magnitude
- What about performance?

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51

- Simplicity reduces attack surface by orders of magnitude
- Starting a single container

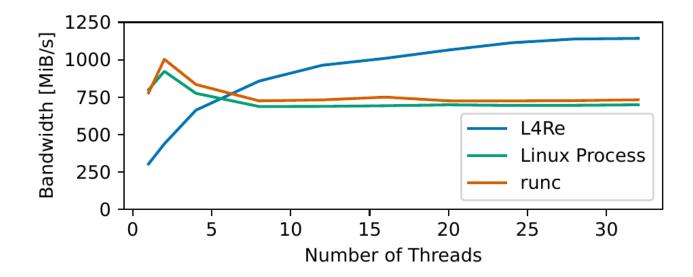


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52

- Simplicity reduces attack surface by orders of magnitude
- I/O performance on a 10G NIC

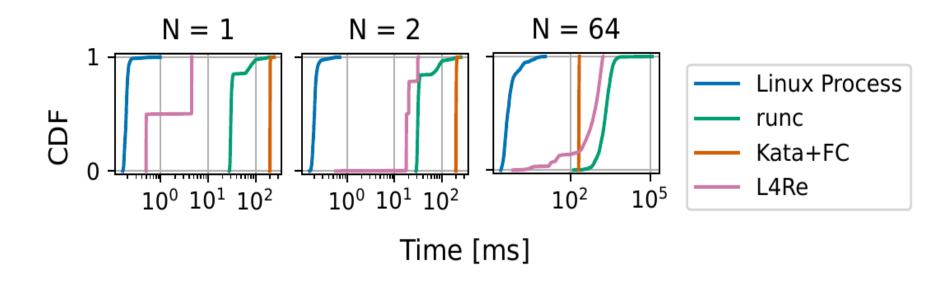


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53

- Simplicity reduces attack surface by orders of magnitude
- Parallel container startup



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Approach	Isolation	Small TCB	Low Overhead	Scalability	Compatibility
Monolith	X	X	✓	✓	✓



Approach	Isolation	Small TCB	Low Overhead	Scalability	Compatibility
Monolith	X	X	✓	✓	
Container	0	X	✓	✓	✓



Approach	Isolation	Small TCB	Low Overhead	Scalability	Compatibility
Monolith	X	X		<b>✓</b>	
Container	0	X	✓	✓	✓
TEE	<b>✓</b>	X	X	?	0



Approach	Isolation	Small TCB	Low Overhead	Scalability	Compatibility
Monolith	X	X		<b>✓</b>	✓
Container	0	X		✓	✓
TEE	<b>✓</b>	X	X	?	0
Light VM	✓	?	0	✓	0



Approach	Isolation	Small TCB	Low Overhead	Scalability	Compatibility
Monolith	X	X		✓	<b>✓</b>
Container	0	X	✓	✓	✓
TEE	<b>✓</b>	X	X	?	0
Light VM	✓	?	0	✓	0
Microkernel	<b>✓</b>		?	?	?



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TEE	✓	X	X	?	0
Light VM	✓	?	0	✓	0
Microkernel	<b>✓</b>	<b>✓</b>	?	?	?

### **Summary**



- Current isolation mechanisms in dynamic cloud settings complex
  - Requirements: Strong isolation between clients, fast startup, good performance
  - Large attack surface hard to avoid with standard system design
- Capability-based microkernels with strong built-in isolation
  - Container isolation is conceptually simple on a microkernel
- Simplicity is key for building secure systems
  - Holds for both software and hardware
  - Modularity to have the choice of including complex, high-performance implementations in TCB

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