TECHNISCHE UNIVERSITÄT DRESDEN

Department of Computer Science Institute of System Architecture, Operating Systems Group

SECURITY ARCHITECTURES

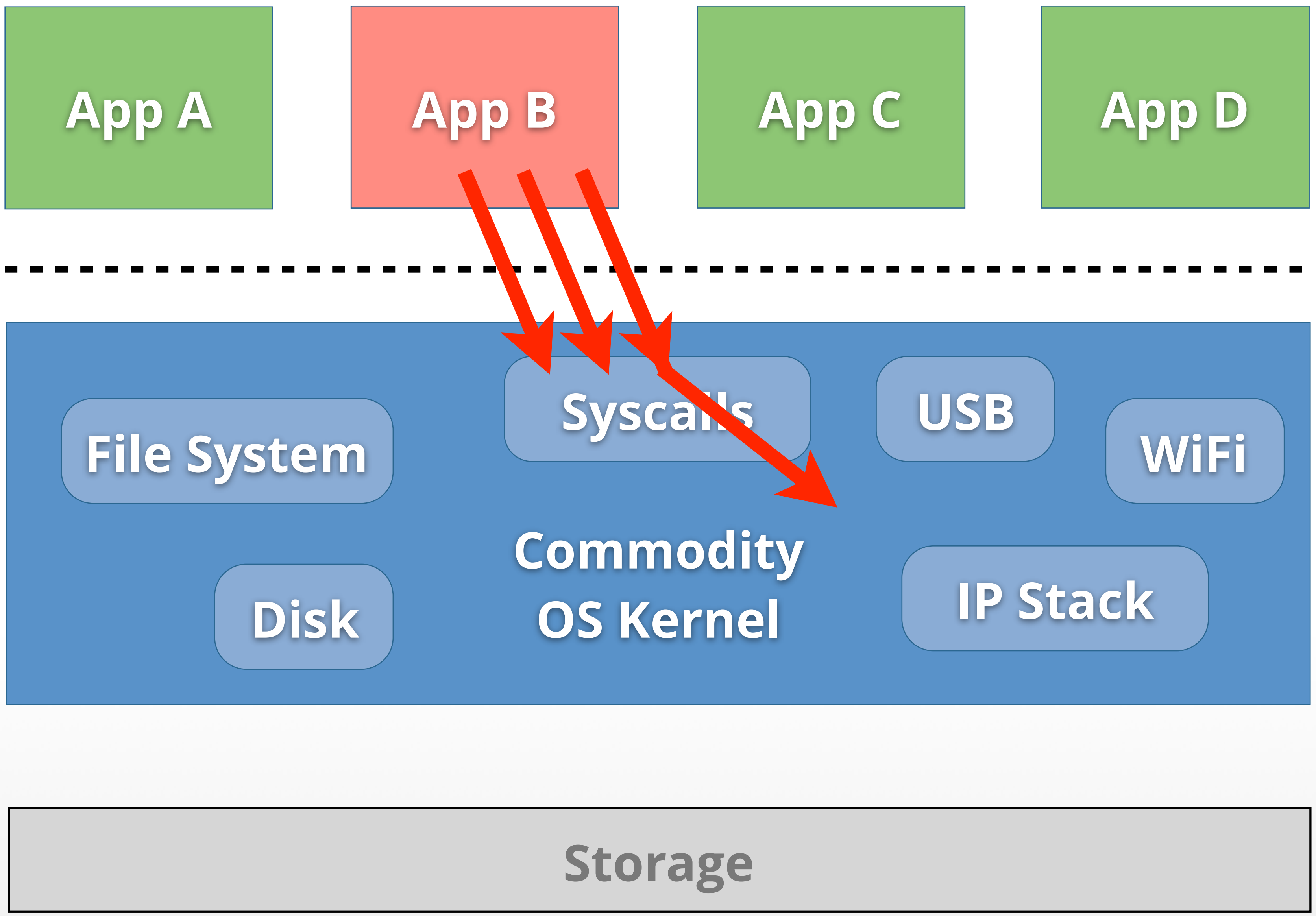
CARSTEN WEINHOLD

CLASSICAL ARCHITECTURES

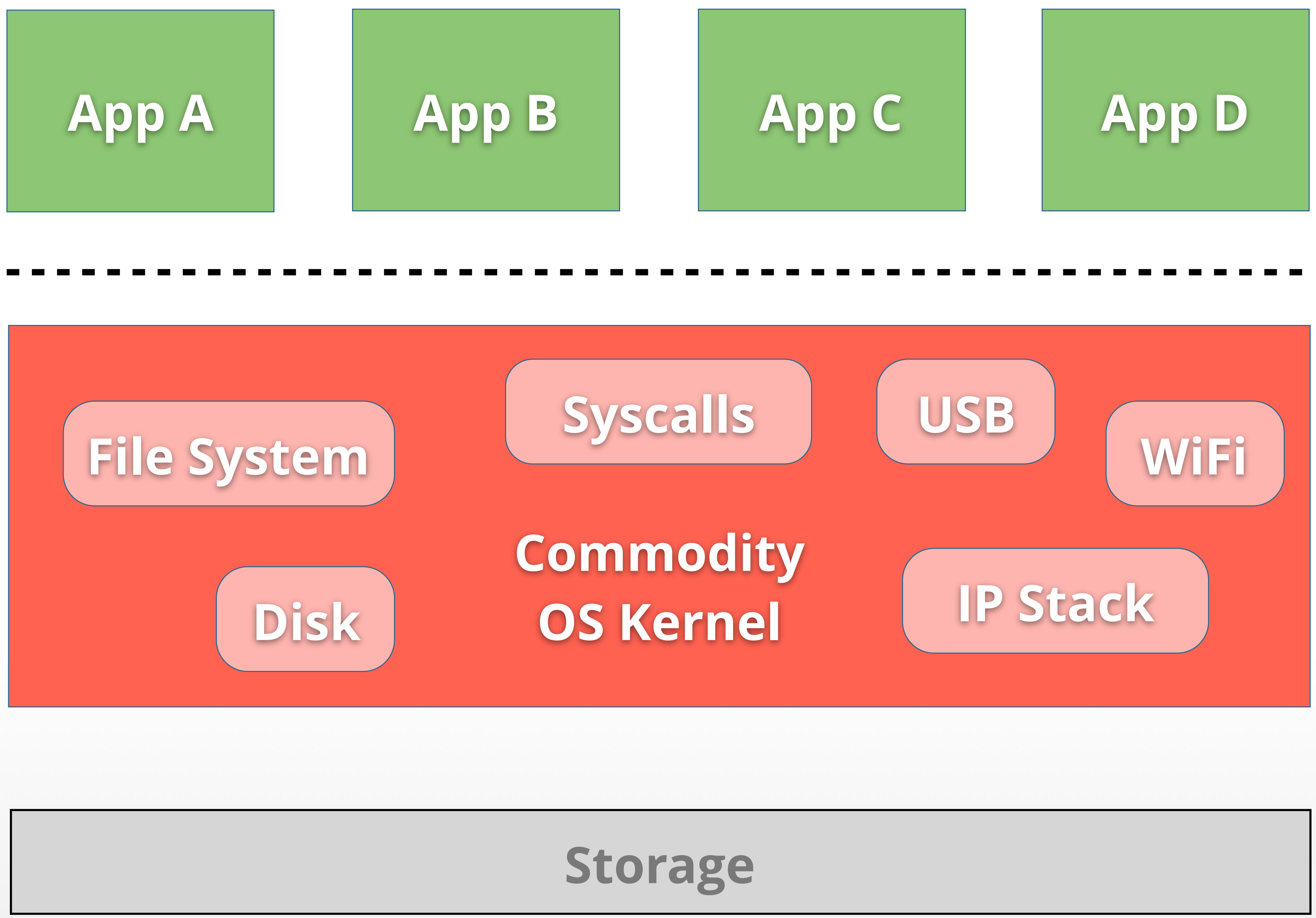
ISOLATION

- Isolation in commodity OSes:
  - Based on user accounts
  - Same privileges for all apps
  - No isolation within applications
  - Permissive interfaces (e.g., ptrace to manipulate other address spaces)

KERNEL ATTACK VECTOR



KERNEL ATTACK VECTOR



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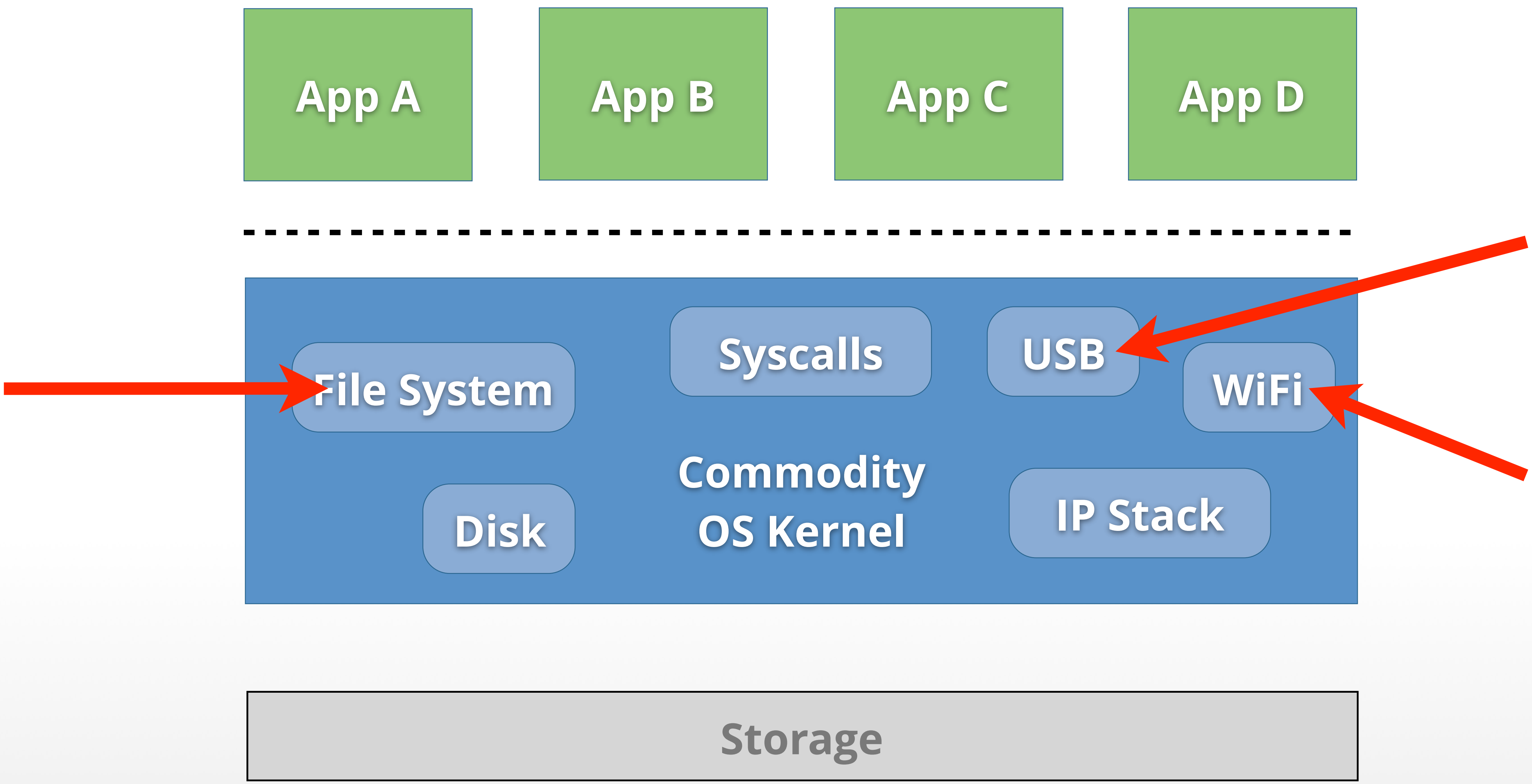

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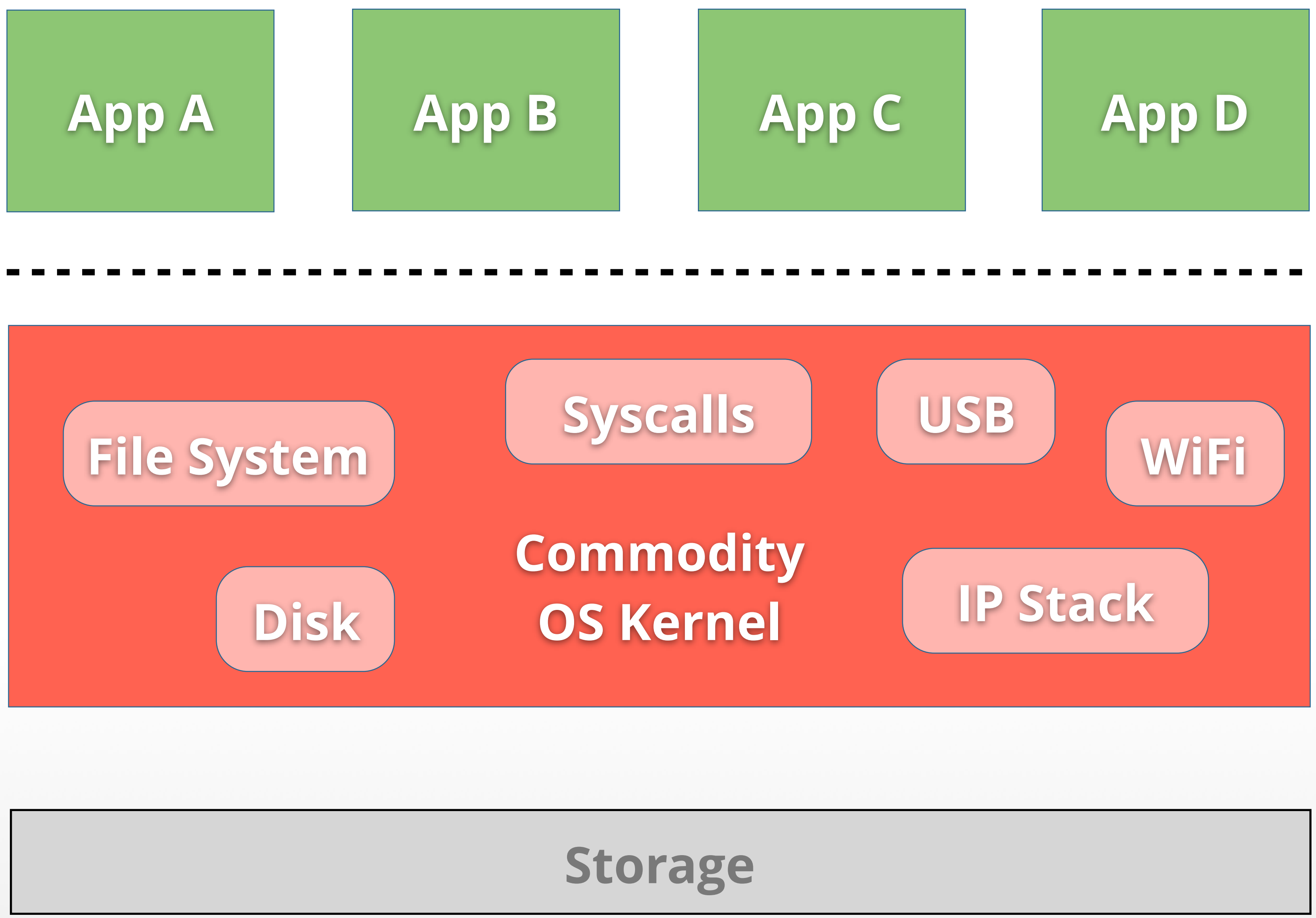
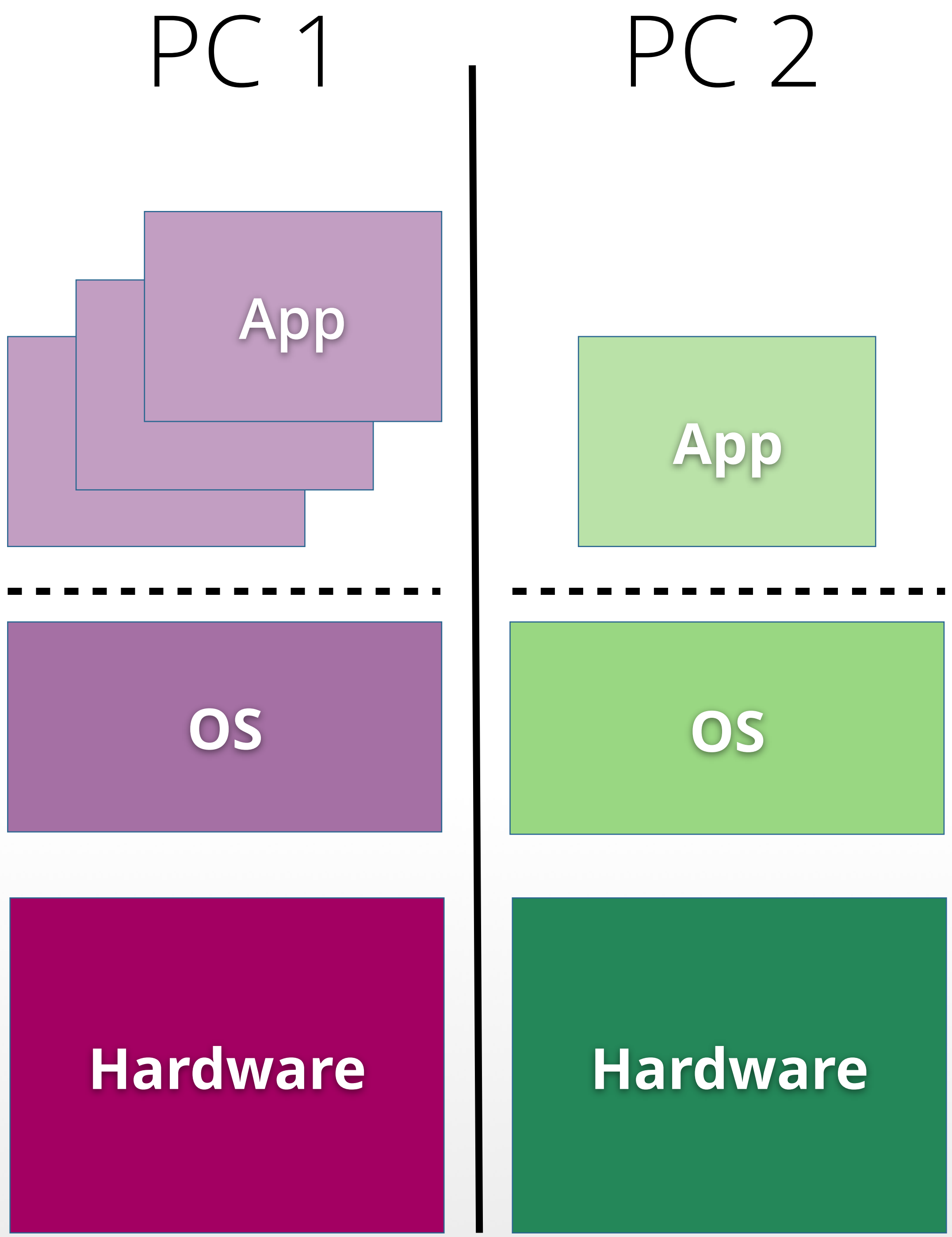


ISOLATION

- Isolation in commodity OSes:
  - Based on user accounts
  - Same privileges for all apps
  - No isolation within applications
  - Permissive interfaces (e.g., ptrace to manipulate other address spaces)
- Efforts to restrict privileges:
  - SELinux, AppArmor, Seatbelt, ...

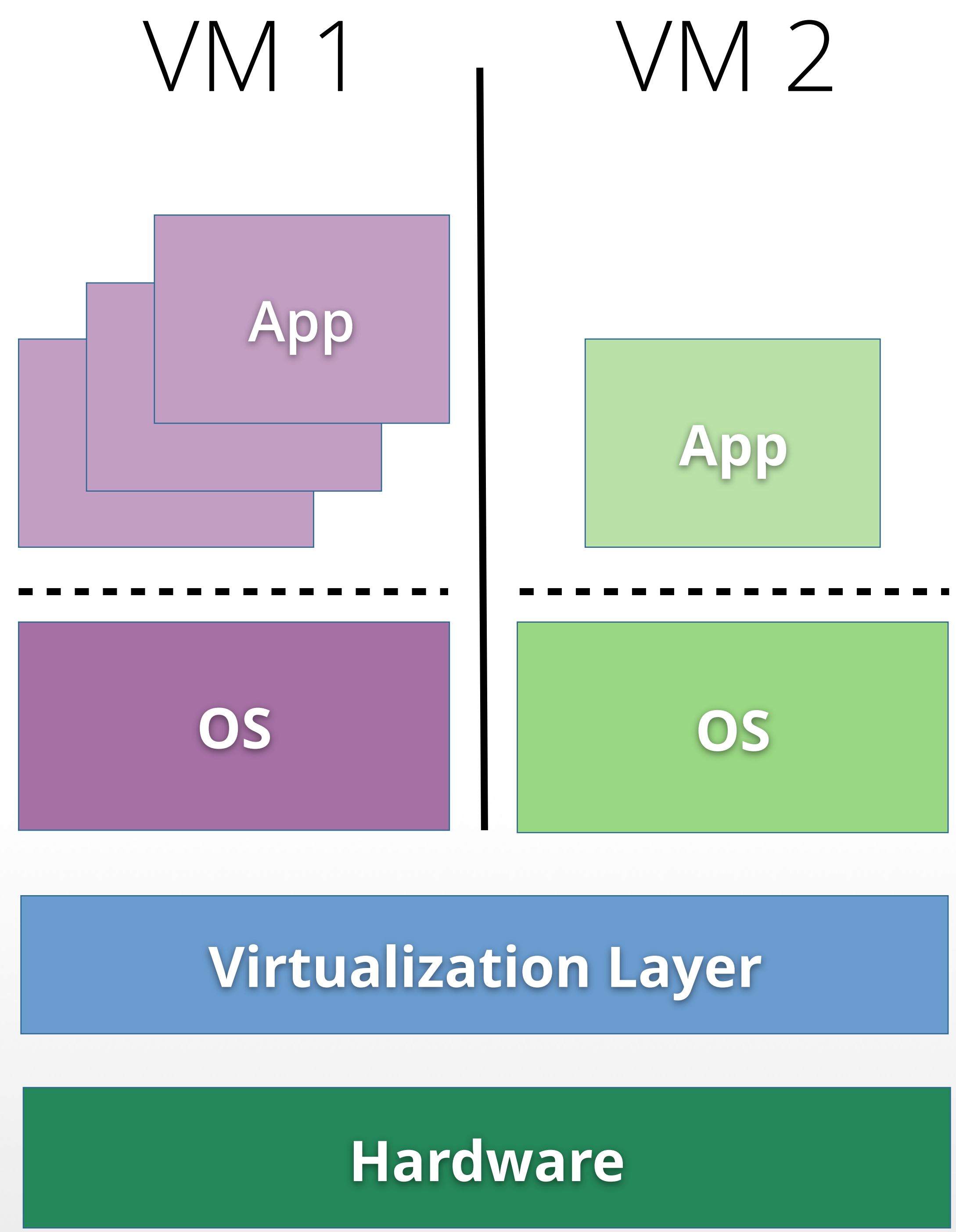
HARDWARE ISOLATION

- Separate computers
- Applications and data physically isolated
- Effective, but …
  - Higher costs
  - Needs more space
  - Inconvenient
  - Exposed to network



VM-BASED ISOLATION

- Multiple VMs, OSes
- Isolation enforced by virtualization layer
- Saves space, energy, maintenance effort
- But still …
  - Switching between VMs is inconvenient
  - Even more code



WHAT IS THE PROBLEM?

- Huge code bases remain
- Applications still the same
- Many targets to attack:
  - Applications, libraries, commodity OSes
  - Virus scanner, firewall, ...
  - Virtualization layer
- High overhead for many VMs

SECURITY ARCHITECTURES

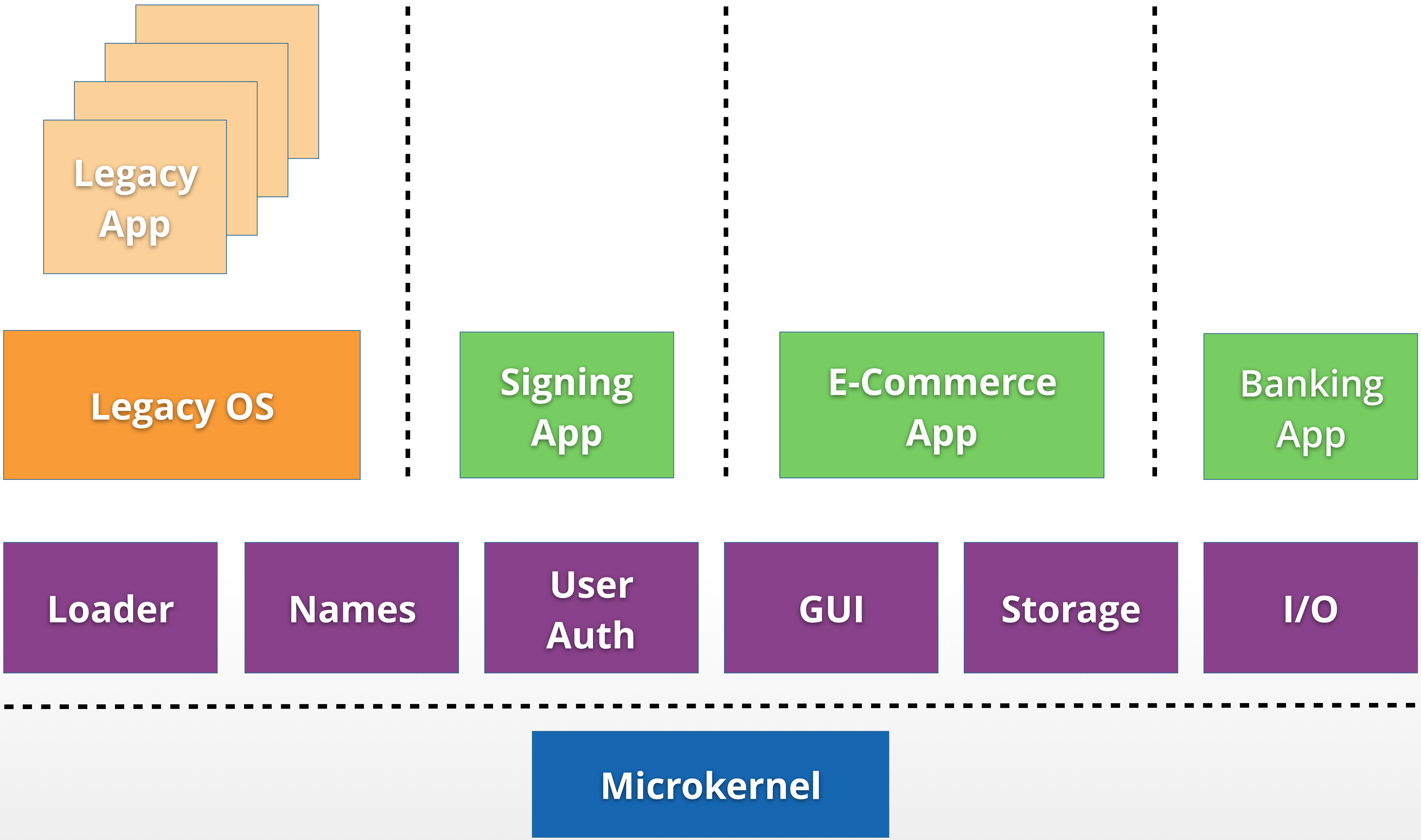
SECURITY GOALS

- Protect the user's data
- Secure applications that process data
- Acknowledge different kinds of trust, e.g.:
  - Application **A** trusted to handle its own data, but not the files of application ***B***
  - OS trusted to <u>*store*</u> data, but not to <u>*see*</u> it
- Identify and secure **TCB:** the **T**rusted **C**omputing **B**ase

APPROACH

- To improve security: Reduce size of TCB = smaller attack surface
- First (incomplete) idea:
  - Remove huge legacy OS from TCB
  - Port application to microkernel-based multi-server OS
  - Remove unneeded libc backends, etc.
  - Possible approaches discussed in lecture on „Legacy Reuse“

NIZZA ARCHITECTURE



NIZZA ARCHITECTURE

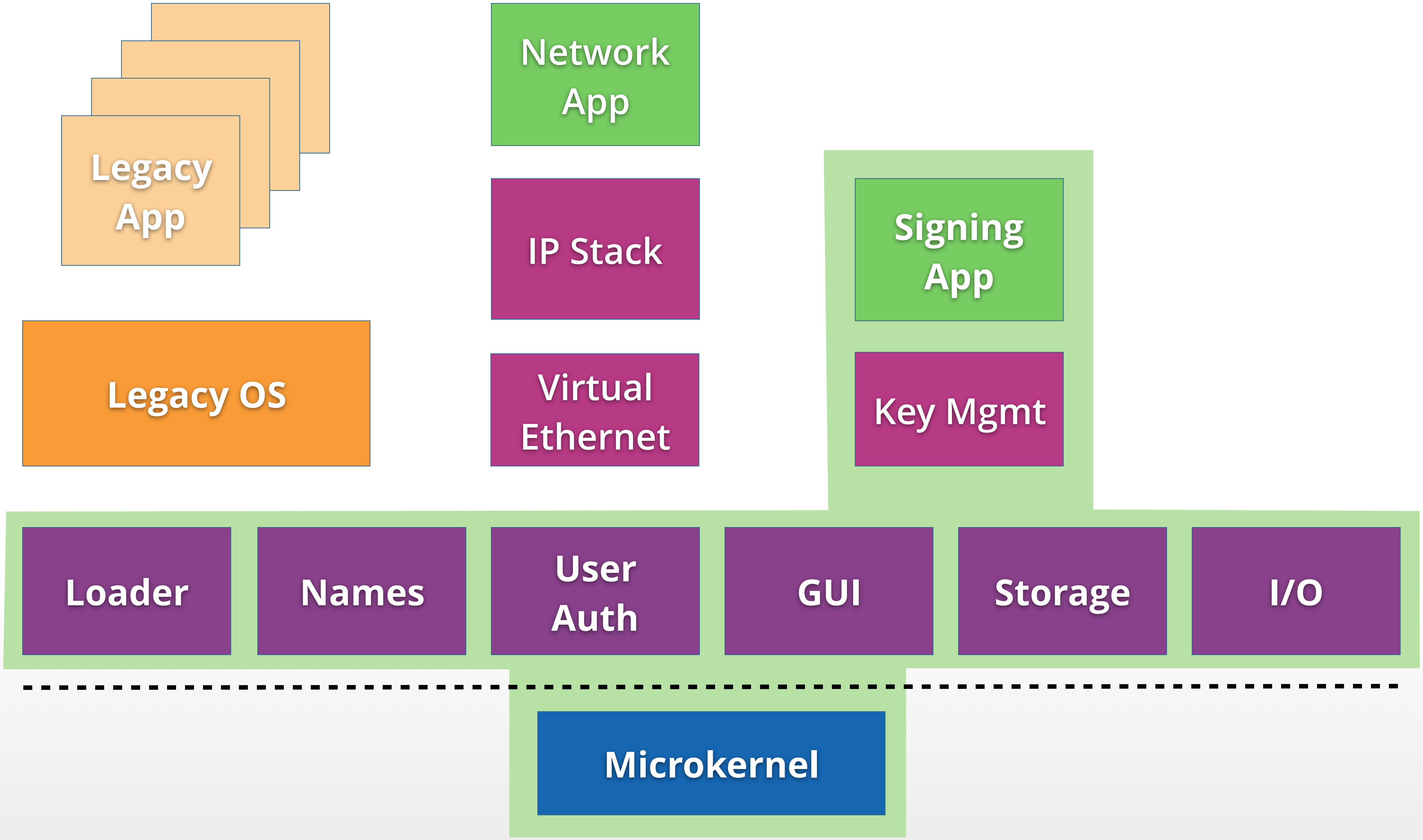
Nizza architecture: fundamental concepts:

- Strong isolation
- Application-specific TCBs
- Legacy reuse
- Trusted wrappers
- Trusted computing

APP-SPECIFIC TCB

- Reflects **Principle of Least Privilege**
- TCB of an application includes only components its security relies upon
- TCB does not include unrelated applications, services, libraries

APP-SPECIFIC TCB



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APP-SPECIFIC TCB

- Reflects **Principle of Least Privilege**
- TCB of an application includes only components its security relies upon
- TCB does not include unrelated applications, services, libraries
- Mechanisms:
  - Address spaces + IPC control for isolation
  - Well-defined interfaces

SPLITTING COMPONENTS

SPLIT APPS

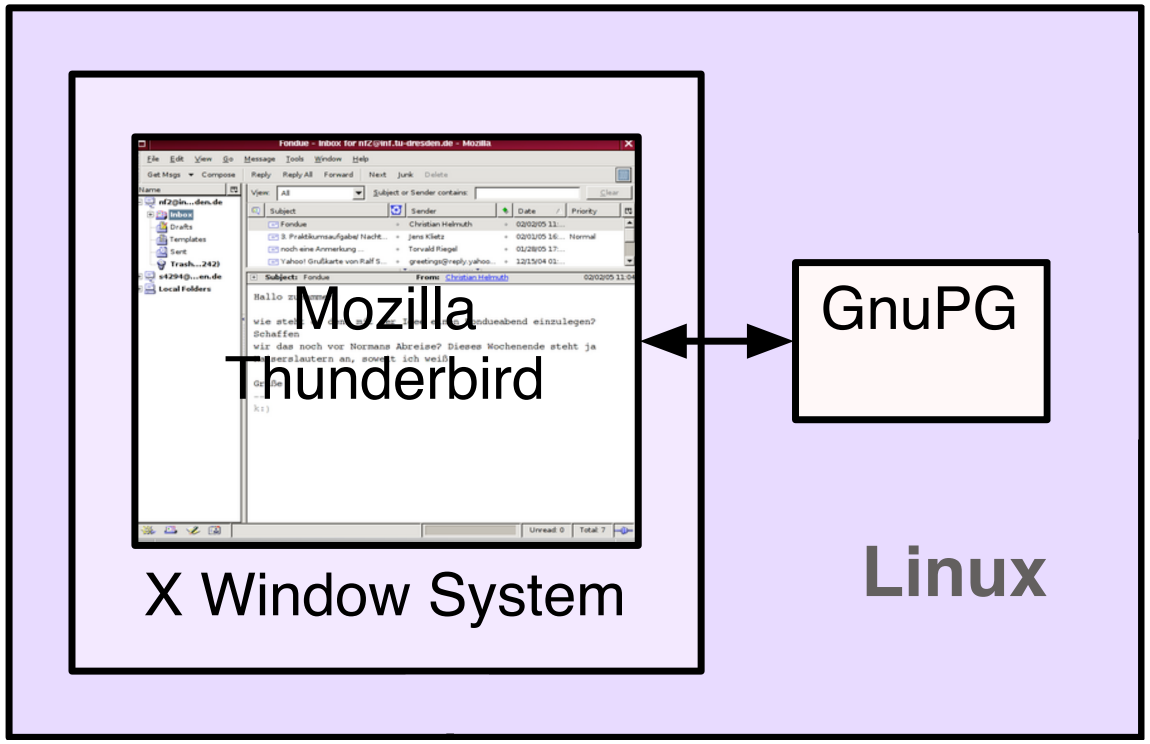
- Problems with porting applications:
  - Dependencies need to be satisfied
  - Can be complex, require lots of code
  - Stripped down applications may lack functionality / usability
- Better idea: split application
  - Make only security-critical parts run on microkernel-based OS
  - Parts of application removed from TCB

EXAMPLE 1: EMAIL

Digitally signed e-mails, what's critical?

- Handling of signature keys
- Requesting passphrase to unlock signature key
- Presenting e-mail message:
  - Before sending: „**W**hat **Y**ou **S**ee **I**s **W**hat **Y**ou **S**ign"
  - After receiving: verify signature, identify sender

STANDARD EMAIL APP



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SPLIT EMAIL APP

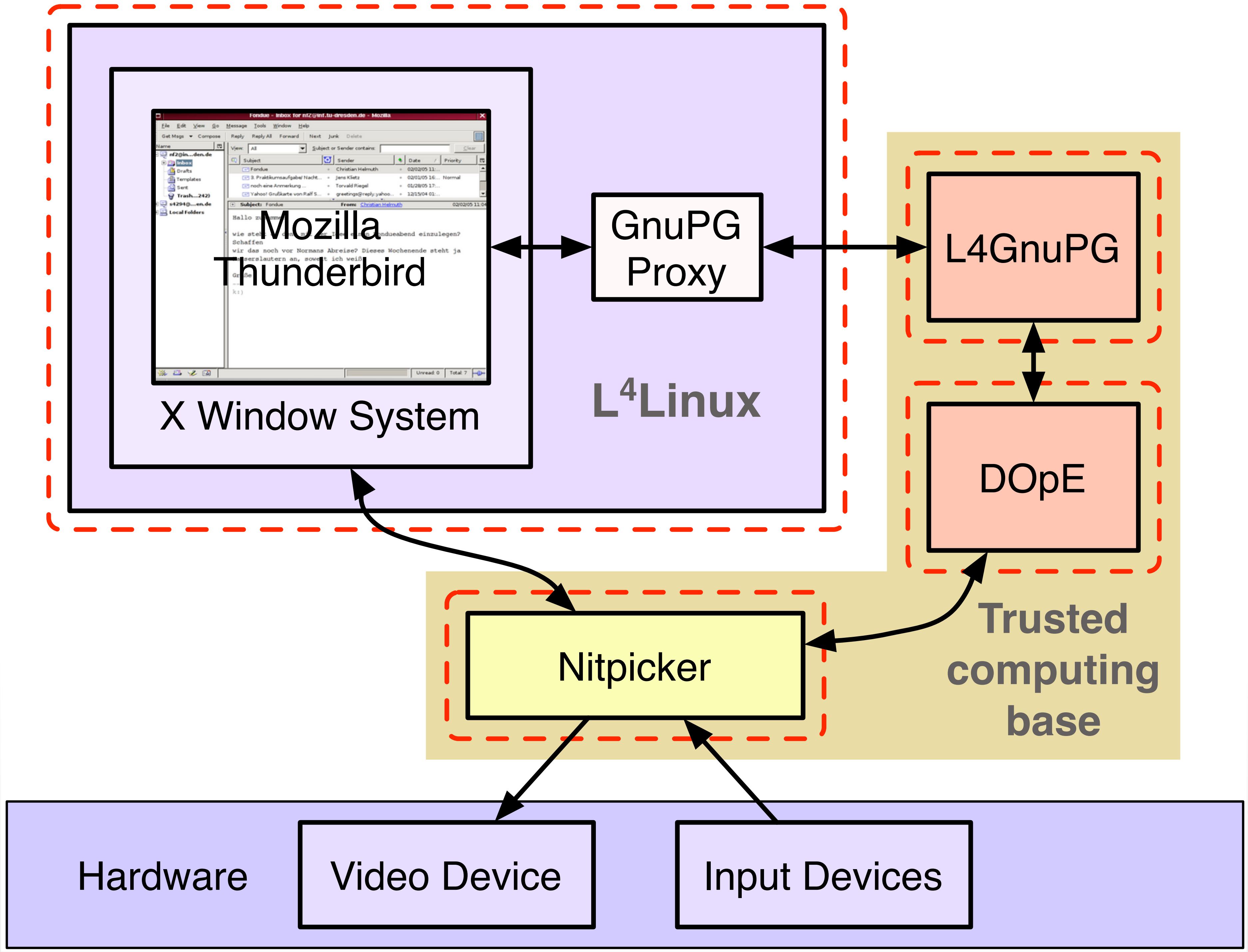


Image source: [5]

TCB REDUCTION

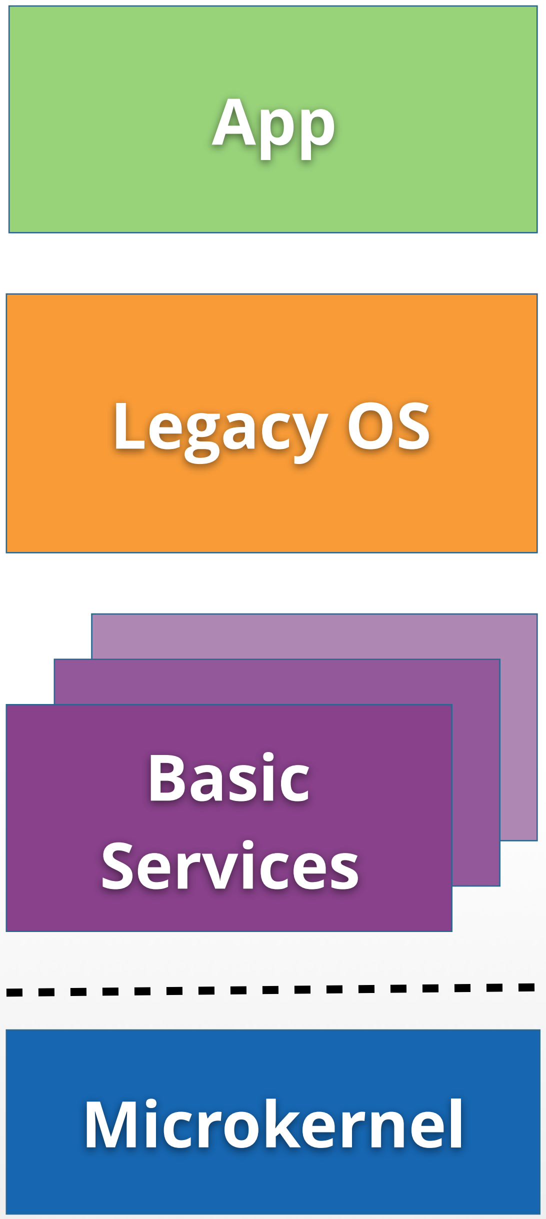
- *1,500,000*+ SLOC no longer in TCB:
  - Linux kernel, drivers, X-Server
  - C and GUI libraries, Thunderbird, ...
- TCB size reduced to ~*150,000* SLOC:
  - GNU Privacy Guard, e-mail viewer
  - Basic L4 system
- At least 10 times less code in TCB

SPLITTING THE OS

- Splitting works for applications
- What about the complex and useful infrastructure of commodity OSes?
  - Drivers (see previous lectures)
  - Protocol stacks (e.g., TCP/IP)
  - File systems
- Starting point: Virtualized commodity OS

SIMPLE REUSE

- Run legacy OS in VM
- Reuse service: net, files, …
- Legacy infrastructure isolated from applications
- But:
  - Applications still depend on legacy services … in TCB?
  - Interfaces reused, security issues as well?



COMPLEXITY + BUGS

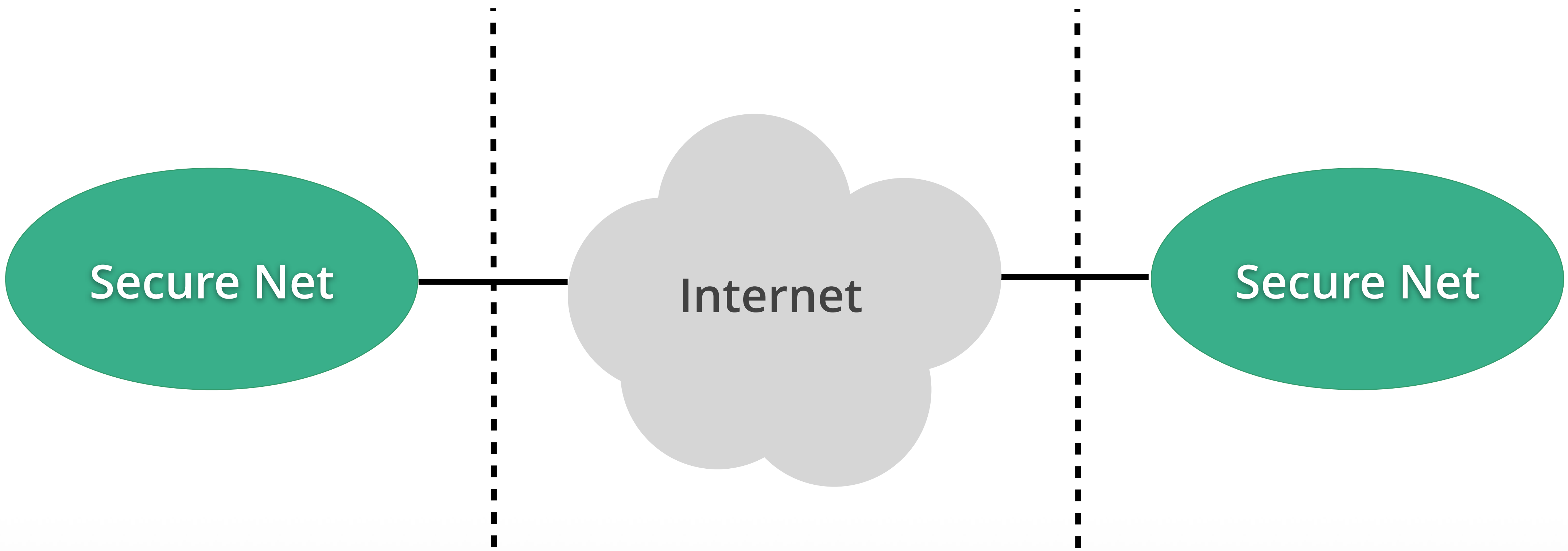
- Network and file system stacks are virtually essential subsystems
- Generally well tested
- Ready for production use
- ... but not bug free [1,2]:
  - Linux file systems (UFS, ISO 9660, Ext3, SquashFS, ...): bug hunt of just 1 month yielded 14 exploitable flaws
  - WiFi drivers: remotely exploitable bugs

REUSE + TRUST

- Complex protocol stacks should not be part of TCB (for confidentiality + integrity)
- Reuse untrusted infrastructure through **Trusted Wrapper:**
  - Add security around existing APIs
    - Cryptography
    - Additional checks (may require copy of critical data, if original data cannot be trusted)
- General idea similar to TLS, VPN

EXAMPLE 2: VPN

VPN: Confidentiality, Integrity, ~~Availability~~



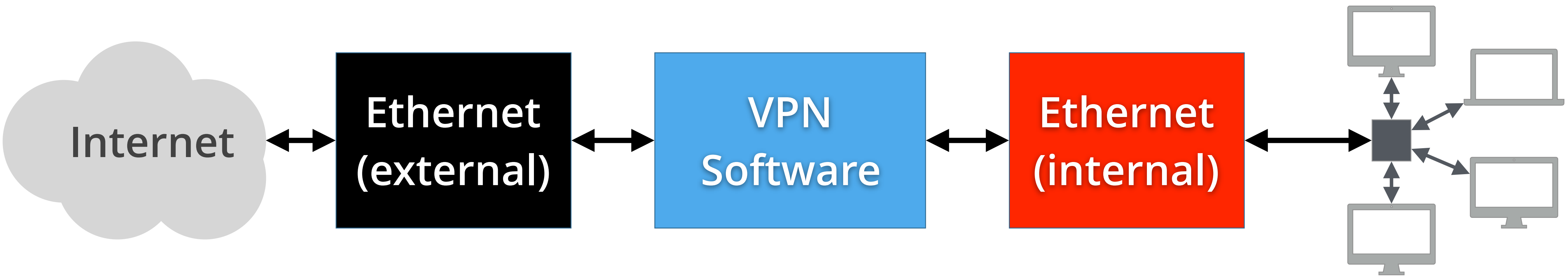
CASE STUDY: SINA BOX

- SINA box used by German „BSI“:
- VPN gateway
- Implements IPSec & PKI
- Intrusion detection & response
- Used for secure access to government networks, e.g., in German embassies

Image source: http://www.secunet.com/de/das-unternehmen/presse/bilddatenbank/

SINA BOX

- Differently trusted network interfaces:
  - **Red:** plaintext, no protection
  - **Black:** encryption + authentication codes



- VPN Software:
  - Based on minimized and hardened Linux
  - Runs only from CD-ROM or Flash

OS COMPLEXITY

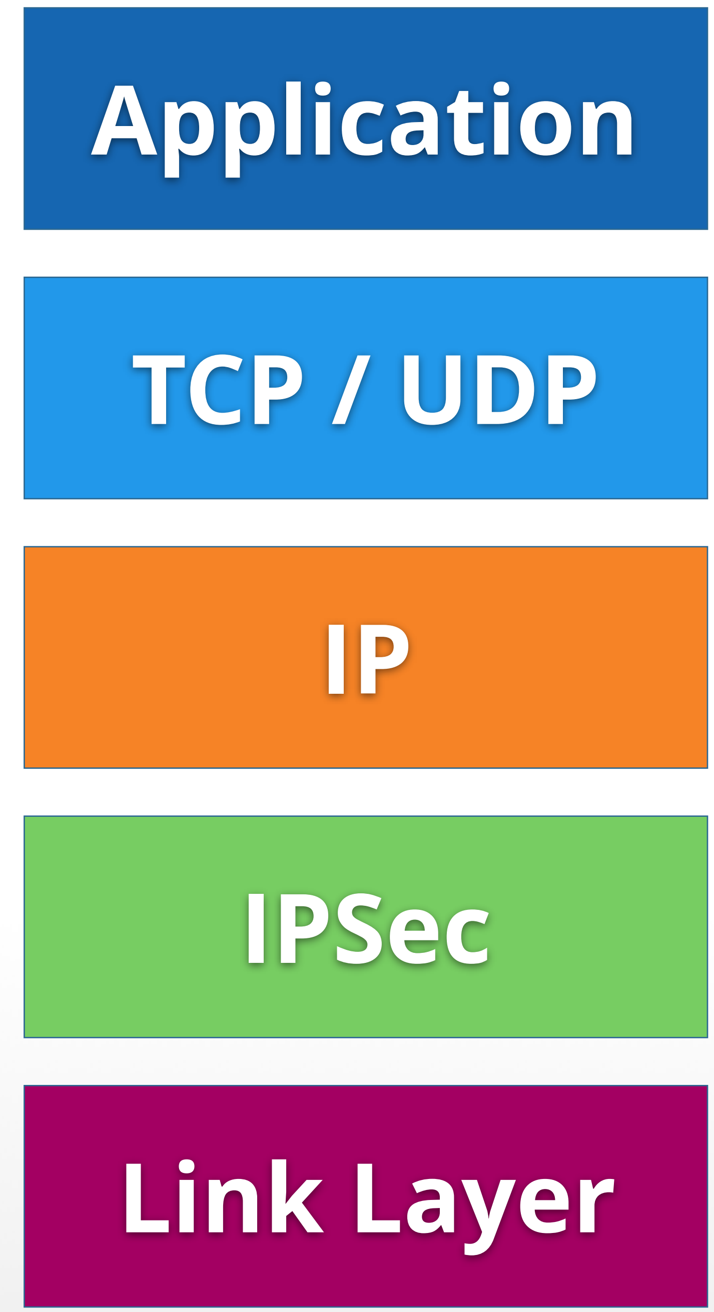
- Linux is complex!
- SLOC for Linux 2.6.18:
  - Architecture specific: 817,880
  - x86 specific: 55,463
  - Drivers: 2,365,256
  - Common: 1,800,587
- Typical config: ~ 2,000,000
- Minimized & hardened: ~ 500,000

MIKRO-SINA

- Research project „Mikro-SINA“
- Goals:
  - Reduce TCB of VPN gateway software
  - Enable high-level evaluation for high assurance scenarios
  - Ensure confidentiality and integrity of sensitive data within the VPN
  - Exploit microkernel architecture

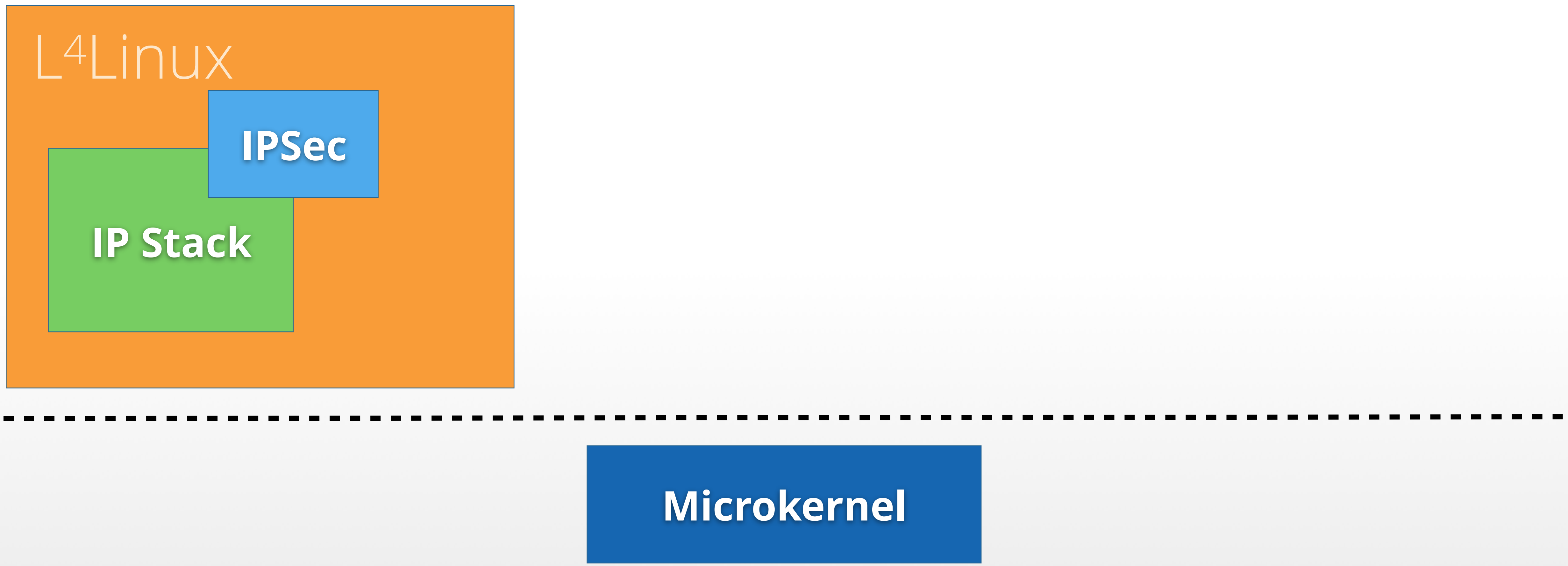
IPSEC BASICS

- Protocol suite for securing IP-based communication
- Authentication header (**AH**)
  - Integrity
  - Authentication
- Encapsulating Security Payload (**ESP**)
  - Confidentiality
- Key management / exchange



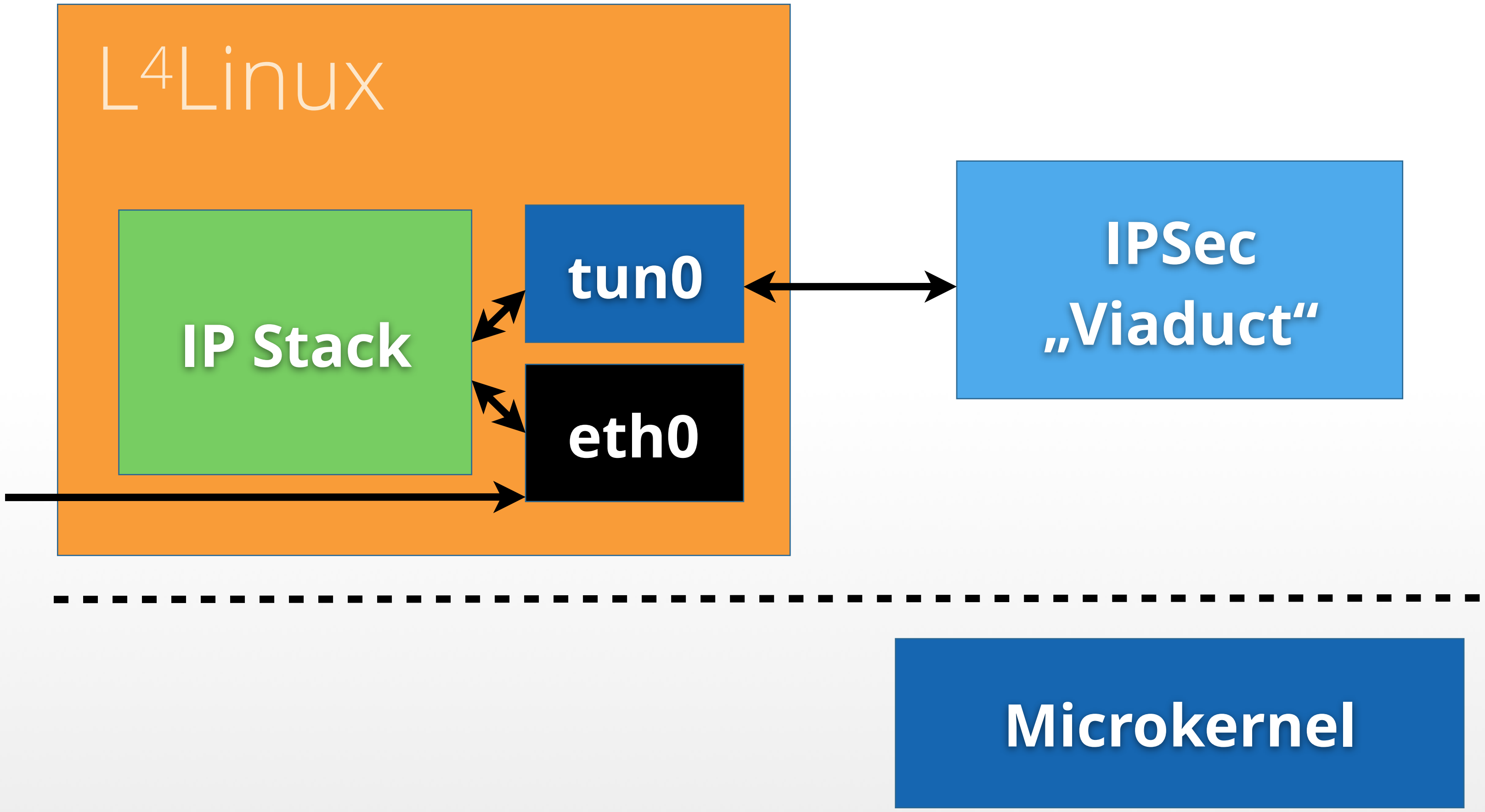
IPSEC IN L4LINUX

- IPSec is security critical component
- … but is integrated into Linux kernel



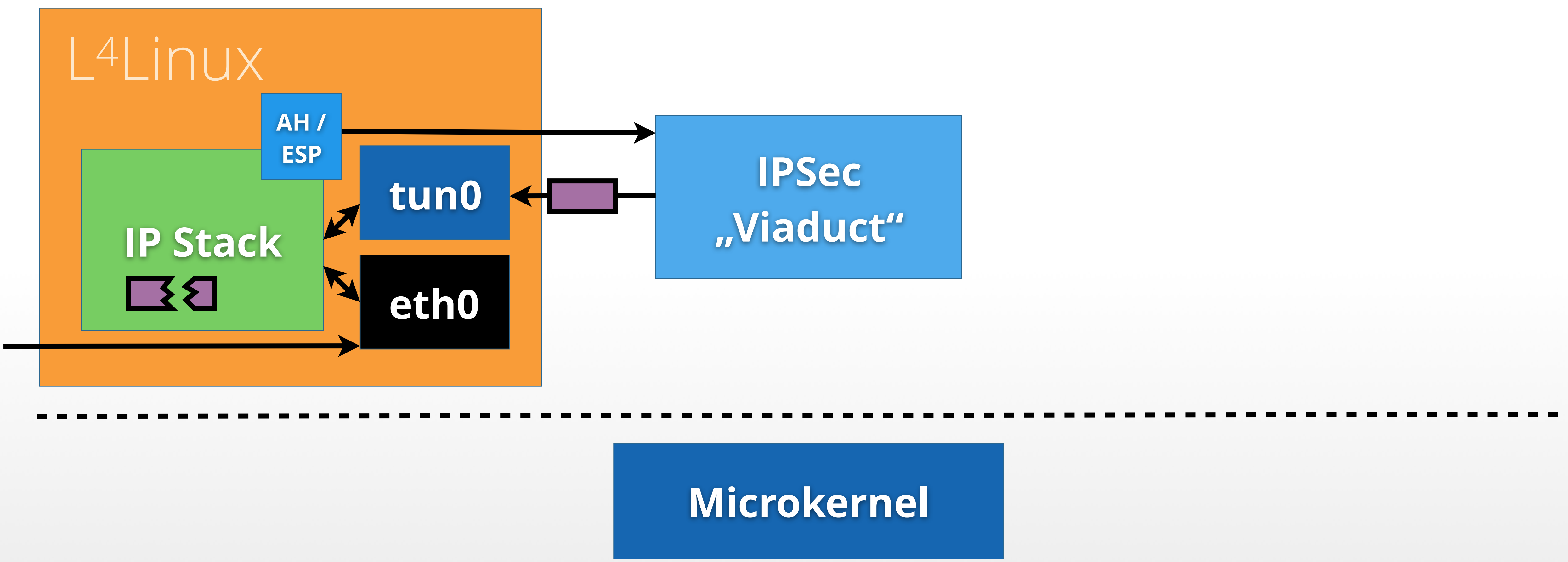
IPSEC „VIADUCT“

- **Idea:** Isolate IPSec in „Viaduct“
- IPSec packets sent/received through TUN/TAP device



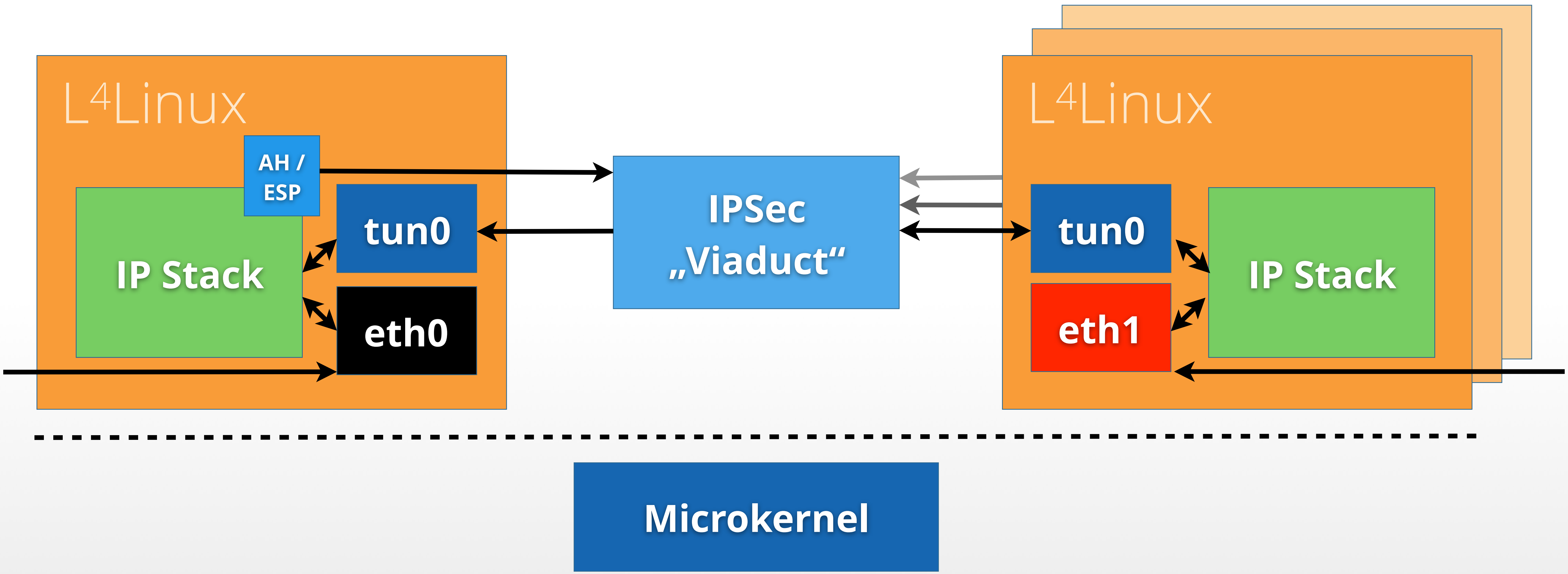
FRAGMENTATION

- Problem: Routers can fragment IPSec packets on the way
- Let L4Linux reassemble them



CONFIDENTIALITY

- Untrusted L4Linux must not see both plaintext and encrypted data
- Dedicated L4Linux for black/red networks

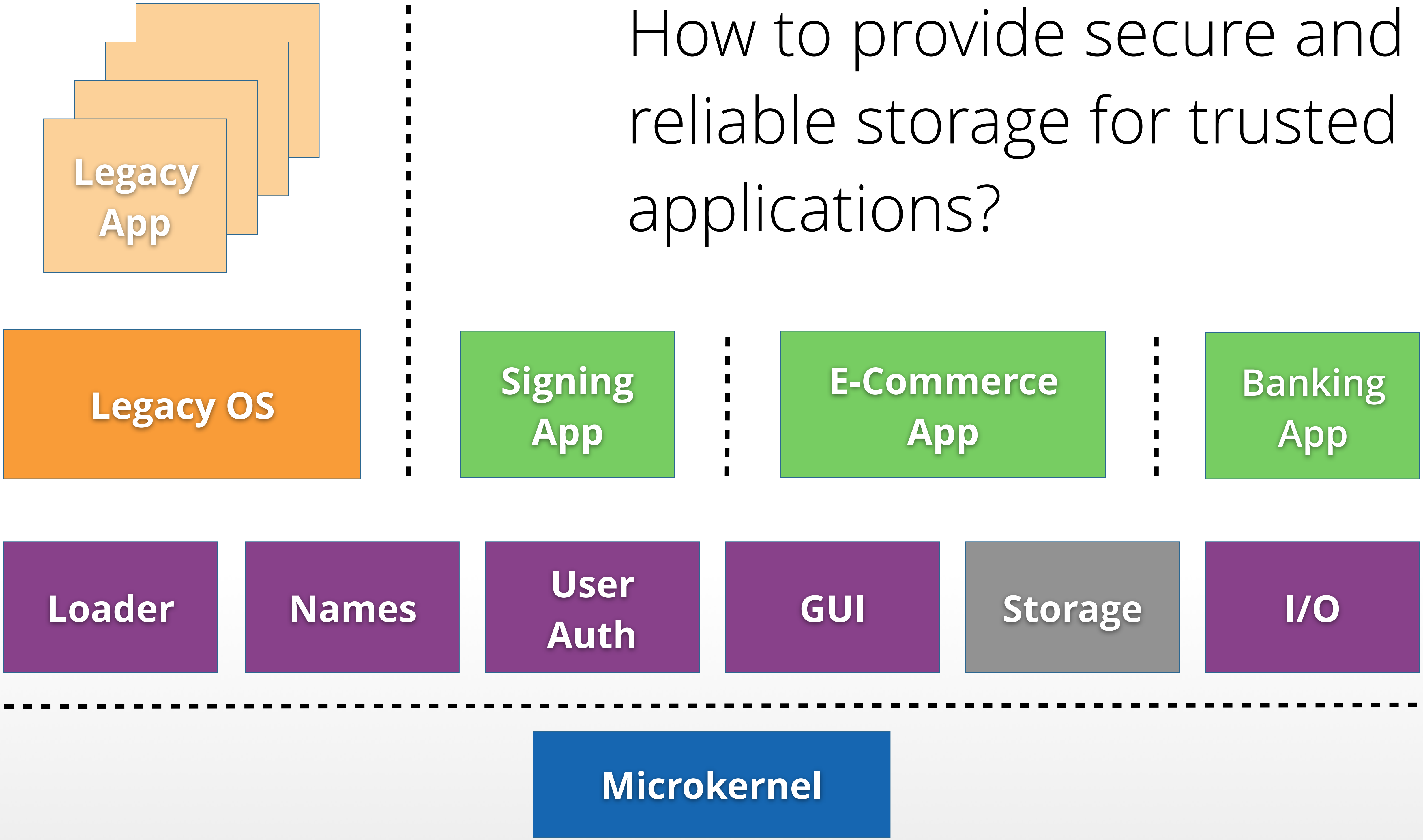


MIKRO-SINA

- Result: Trusted Wrapper for VPN
- Small TCB (see [6] for details):
  - 5,000 SLOC for „Viaduct“
  - Fine grain isolation
  - Principle of least privilege
- Extensive reuse of legacy code:
  - Drivers
  - IP stack

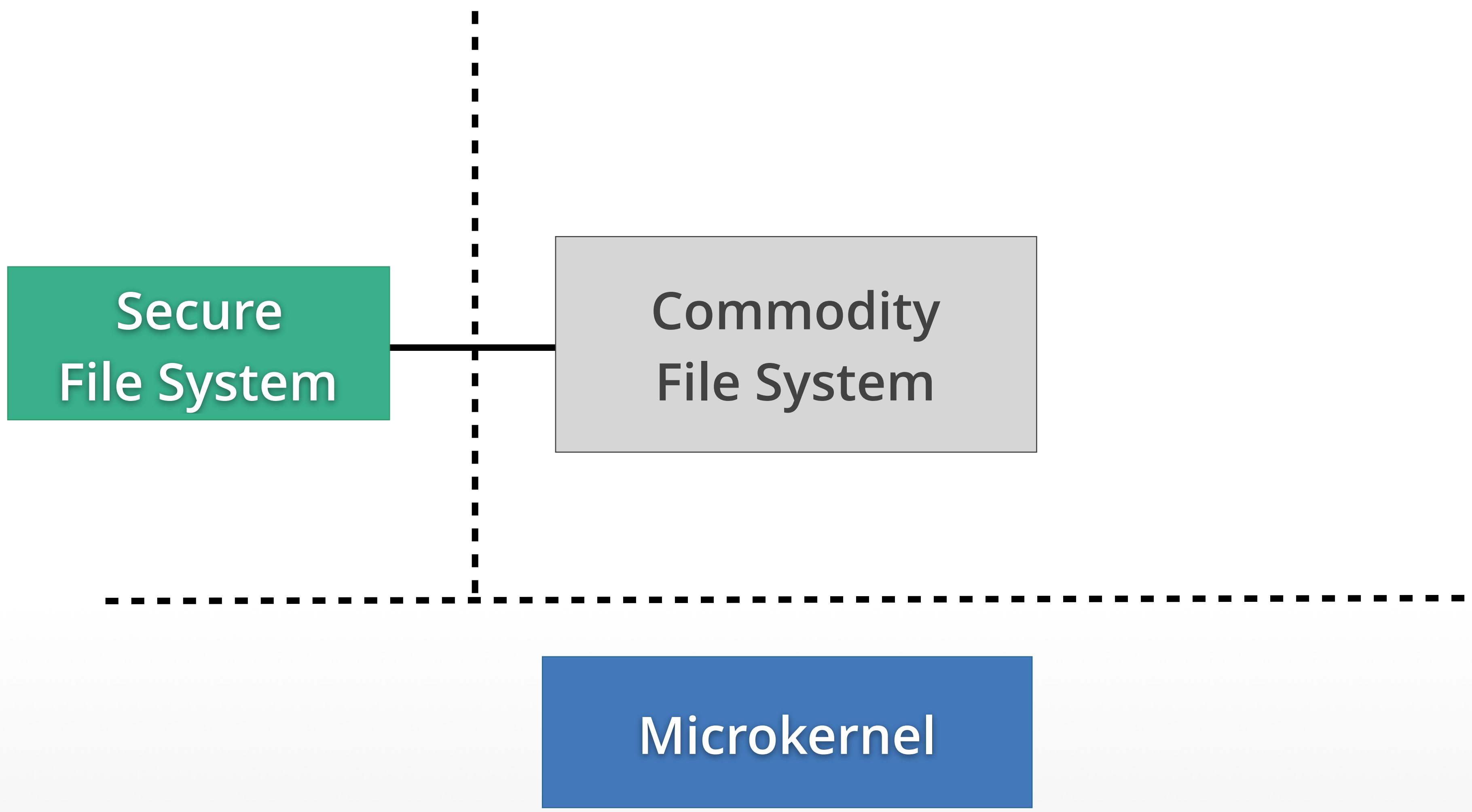
EXAMPLE 3: STORAGE

How to provide secure and reliable storage for trusted applications?



VIRTUAL PRIVATE…

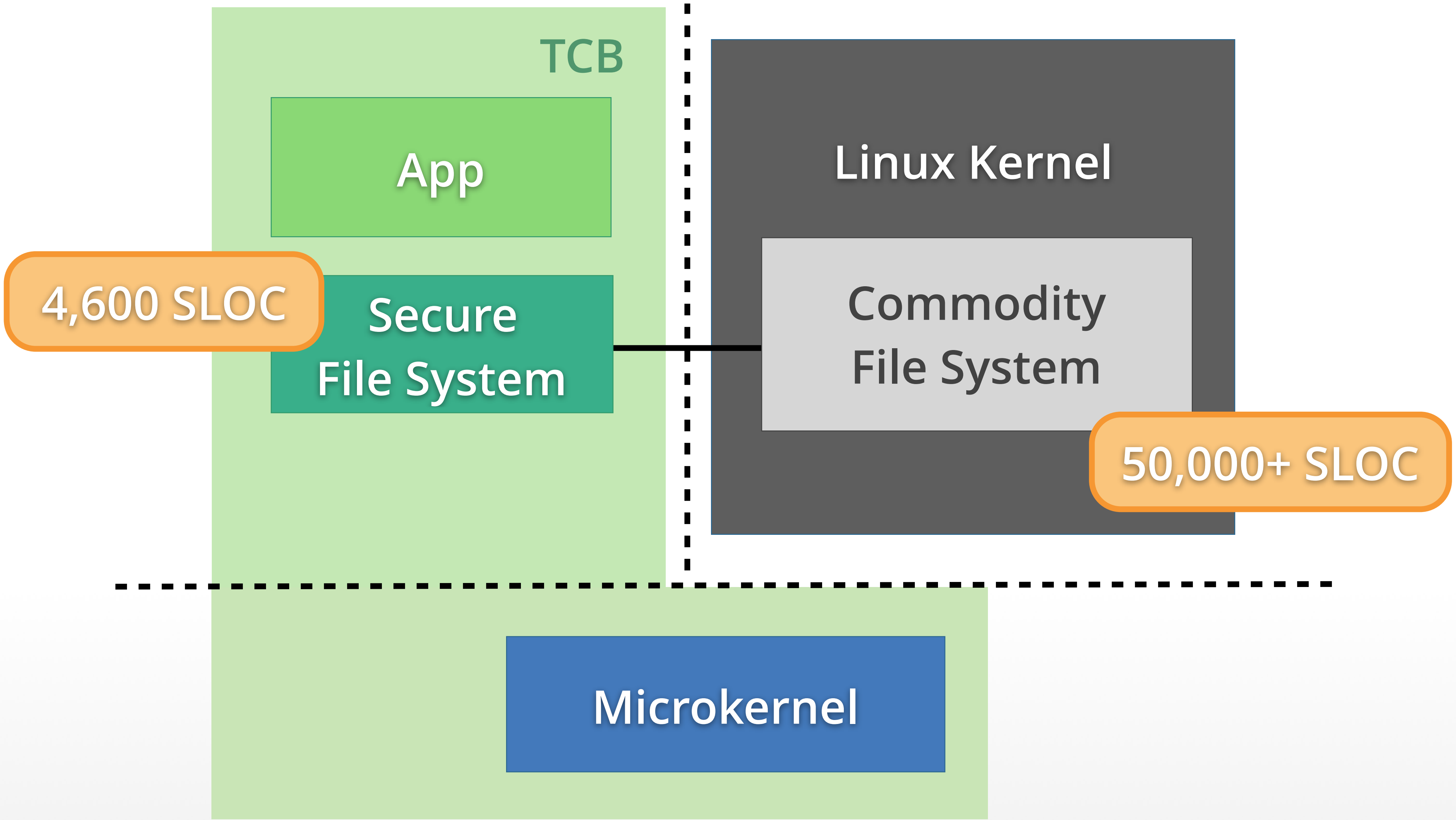
VPFS: Confidentiality, Integrity, ~~Availability~~



See [3] for details

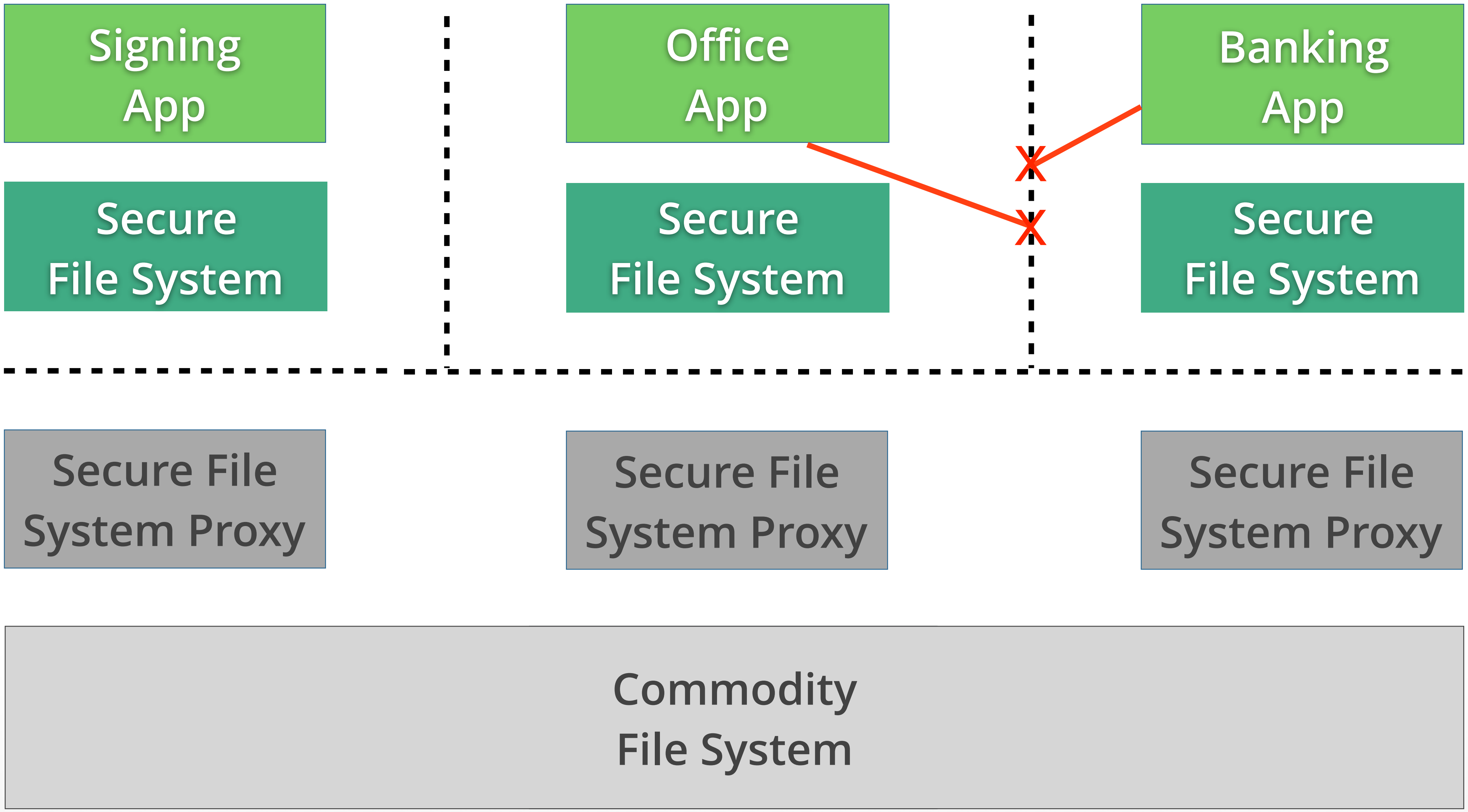
VIRTUAL PRIVATE...

VPFS: Confidentiality, Integrity, ~~Availability~~



See [3] for details

VPFS STACK

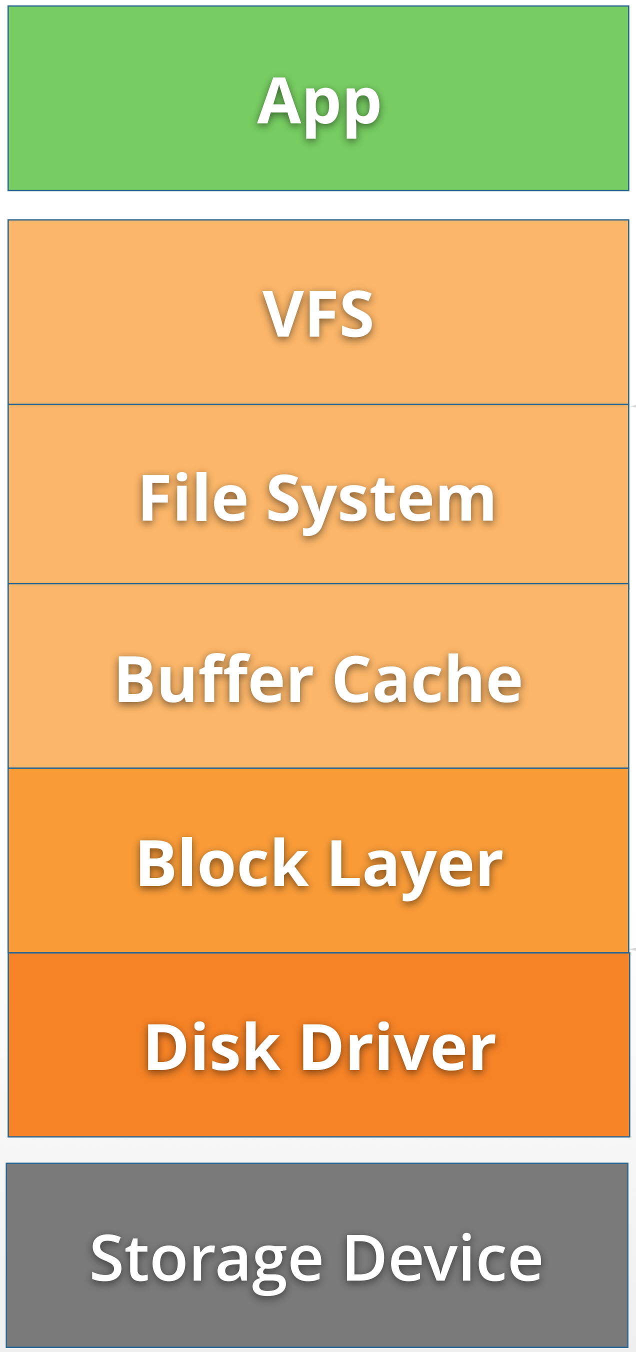


Isolate applications and their private storage: configure communication capabilities such that each application can access its private instance of the secure file system exclusively

SECURITY GOALS

- **Confidentiality:** only authorized applications can access file system, all untrusted software cannot get any useful information
- **Integrity:** all data and meta data is correct, complete, and up to date; otherwise report integrity error
- **Recoverability:** damaged data in untrusted file system can be recovered

POPULAR SOLUTIONS



File-level protection

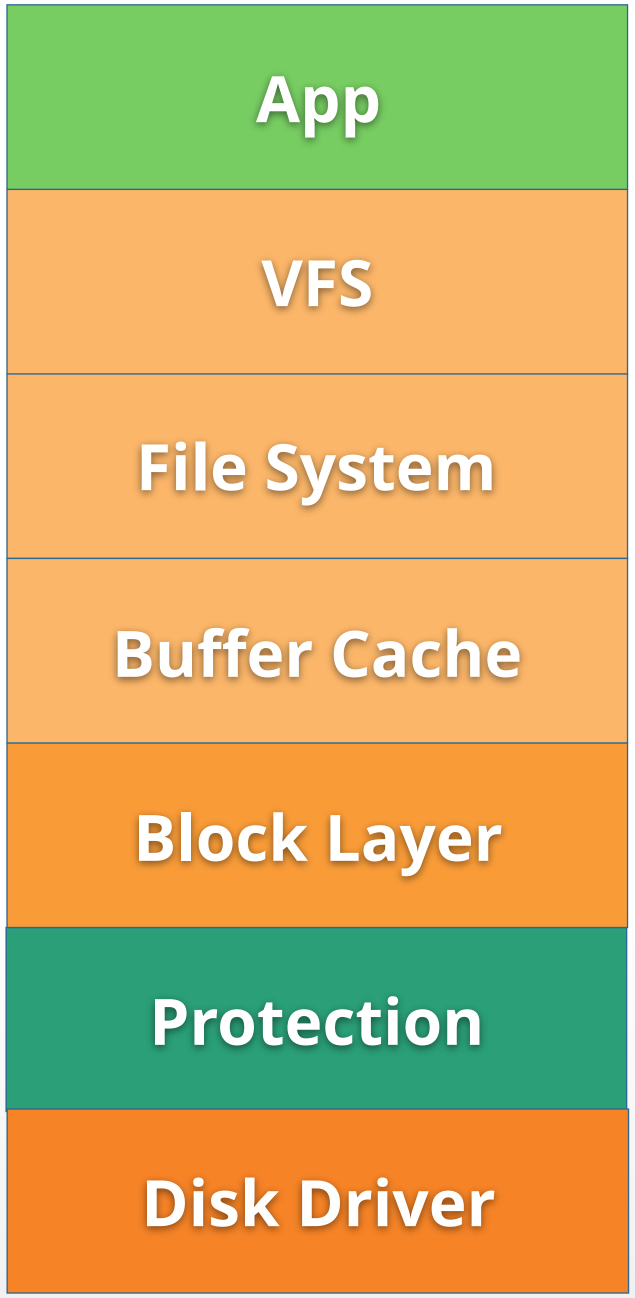
CFS	Cryptographic File System for UNIX
EFS	Microsoft Encrypting File System
ecryptfs	Linux kernel support + tools
EncFS	Based on FUSE

Volume-level protection

TrueCrypt, **Filevault 2**
dm_crypt
Bitlocker
Encrypted volumes in smartphones, etc.

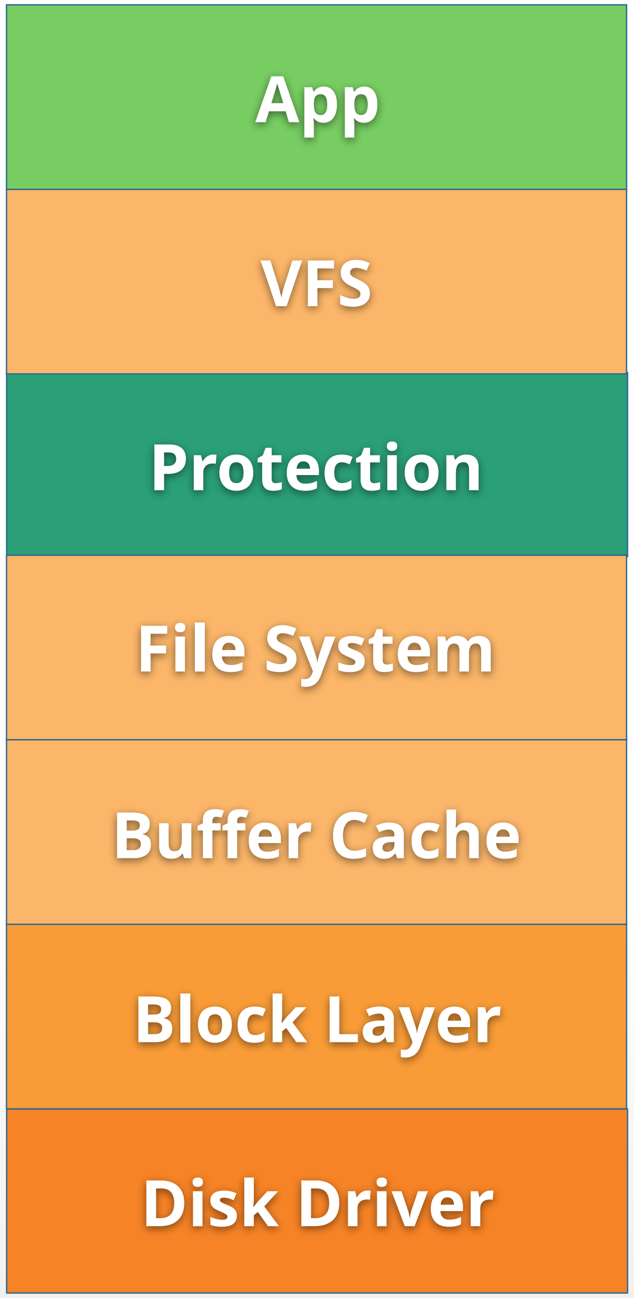
DESIGN OPTIONS

- First end of design space: Protect at block layer
  - Transparent encryption of all data and metadata
  - Block-level integrity ???
  - Most parts of file system stack are part of TCB
  - Attack surface still big

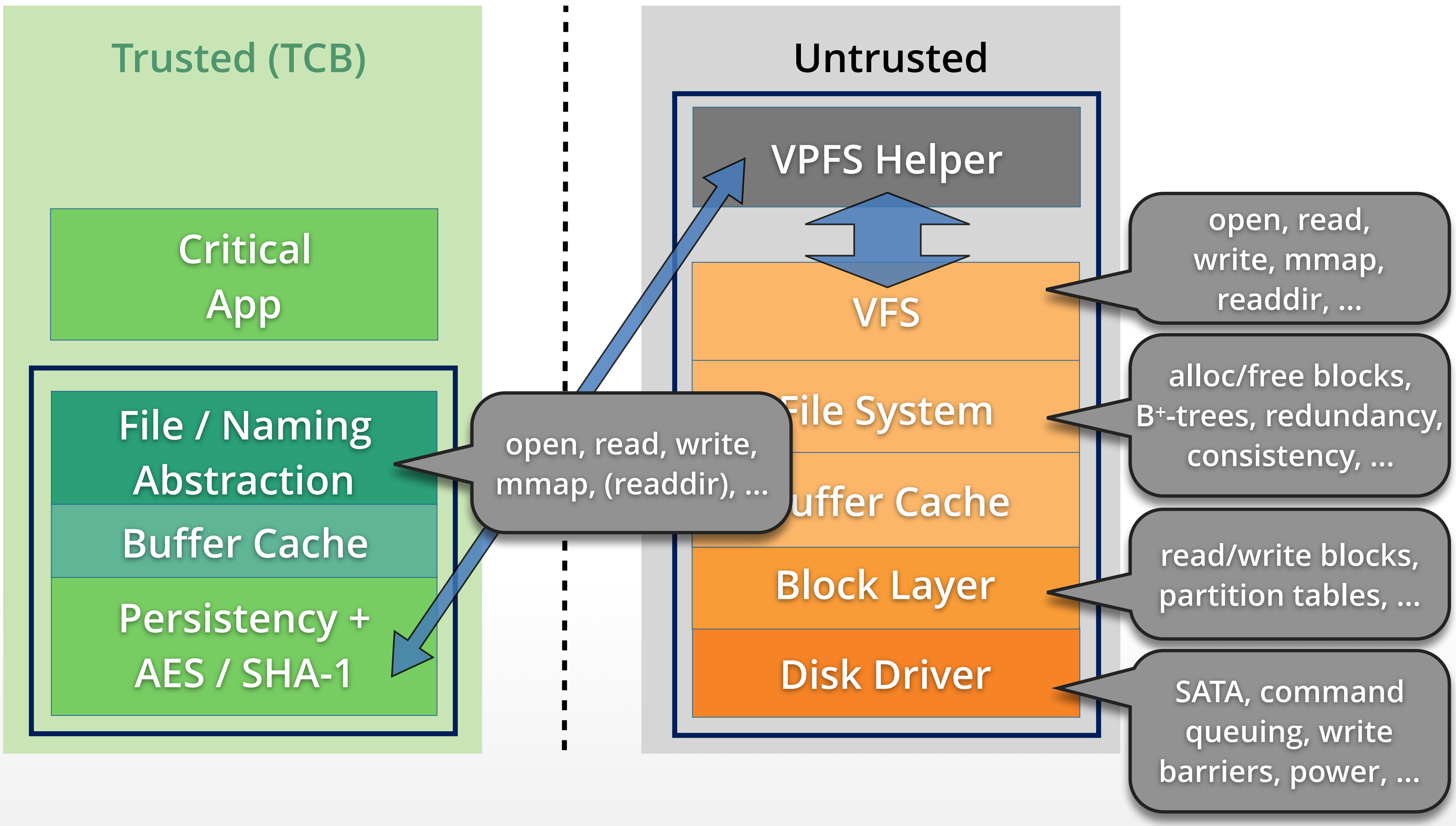


DESIGN OPTIONS

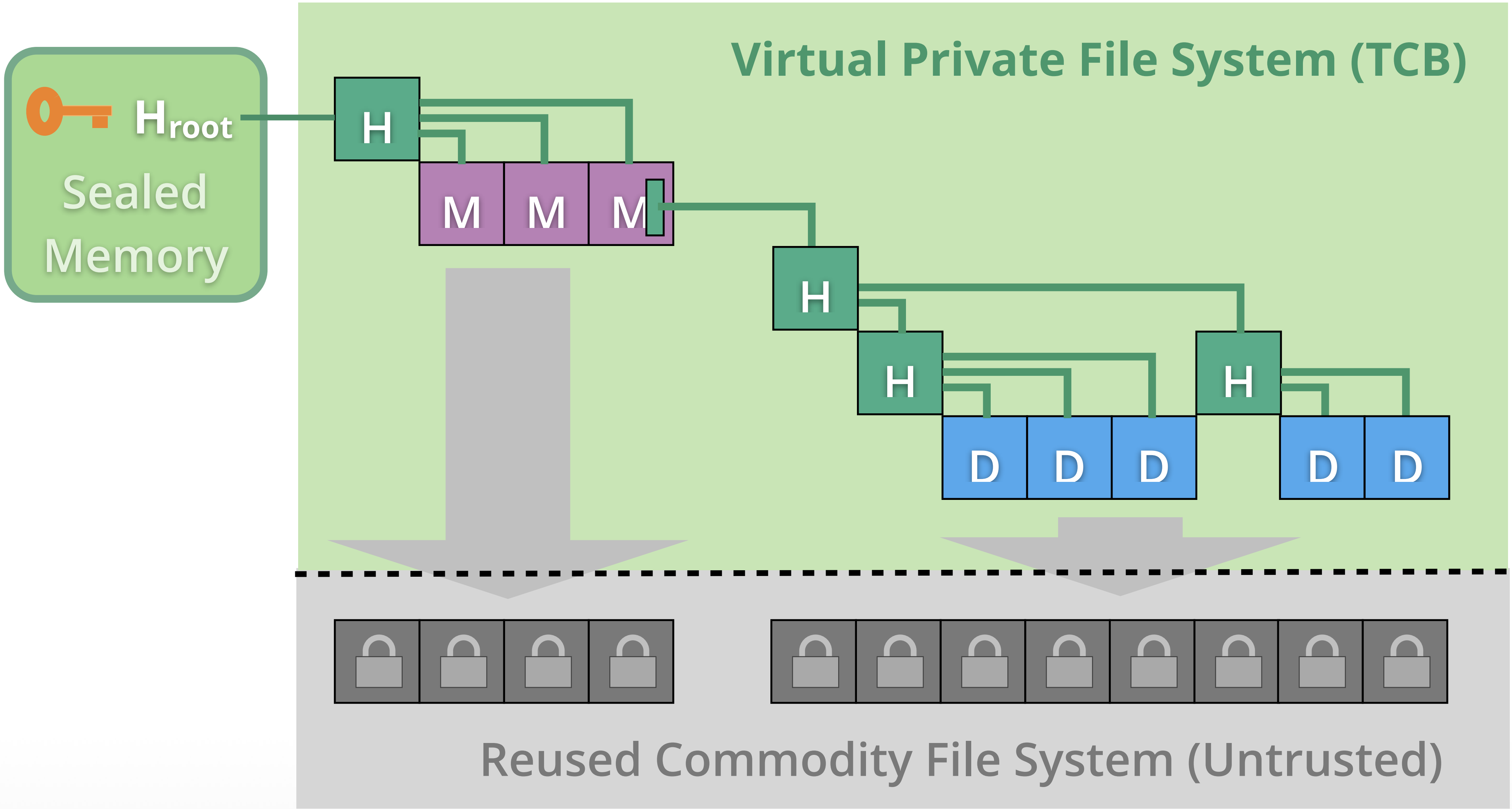
- Second end of design space: Protect individual files
  - Stacked file system
  - Encrypt all data and some metadata (directories, …)
  - More flexibility for integrity
  - Most parts of file system stack <u>not</u> part of TCB
  - Ideal for trusted wrapper



TRUSTED WRAPPER



VPFS APPROACH

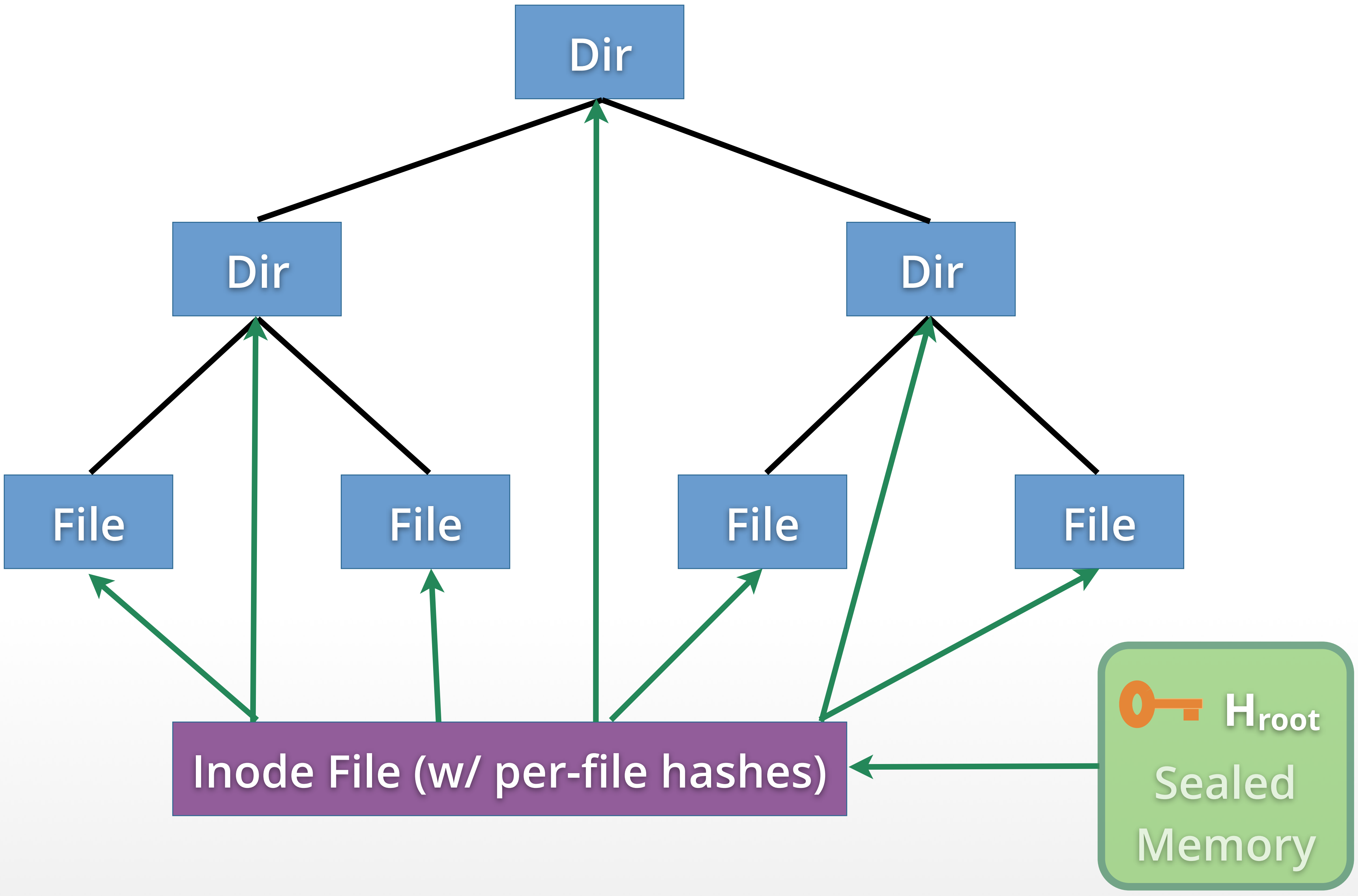


- Encrypted files in commodity file system
- Merkle **hash tree** to detect tampering

VPFS APPROACH

- Trusted part of VPFS enforces security:
  - Encryption / decryption on the fly
  - Plaintext only in trusted buffer cache
  - Files in untrusted commodity file system store encrypted blocks
  - Hash tree protects integrity of complete file system
  - Single hash of root node stored securely

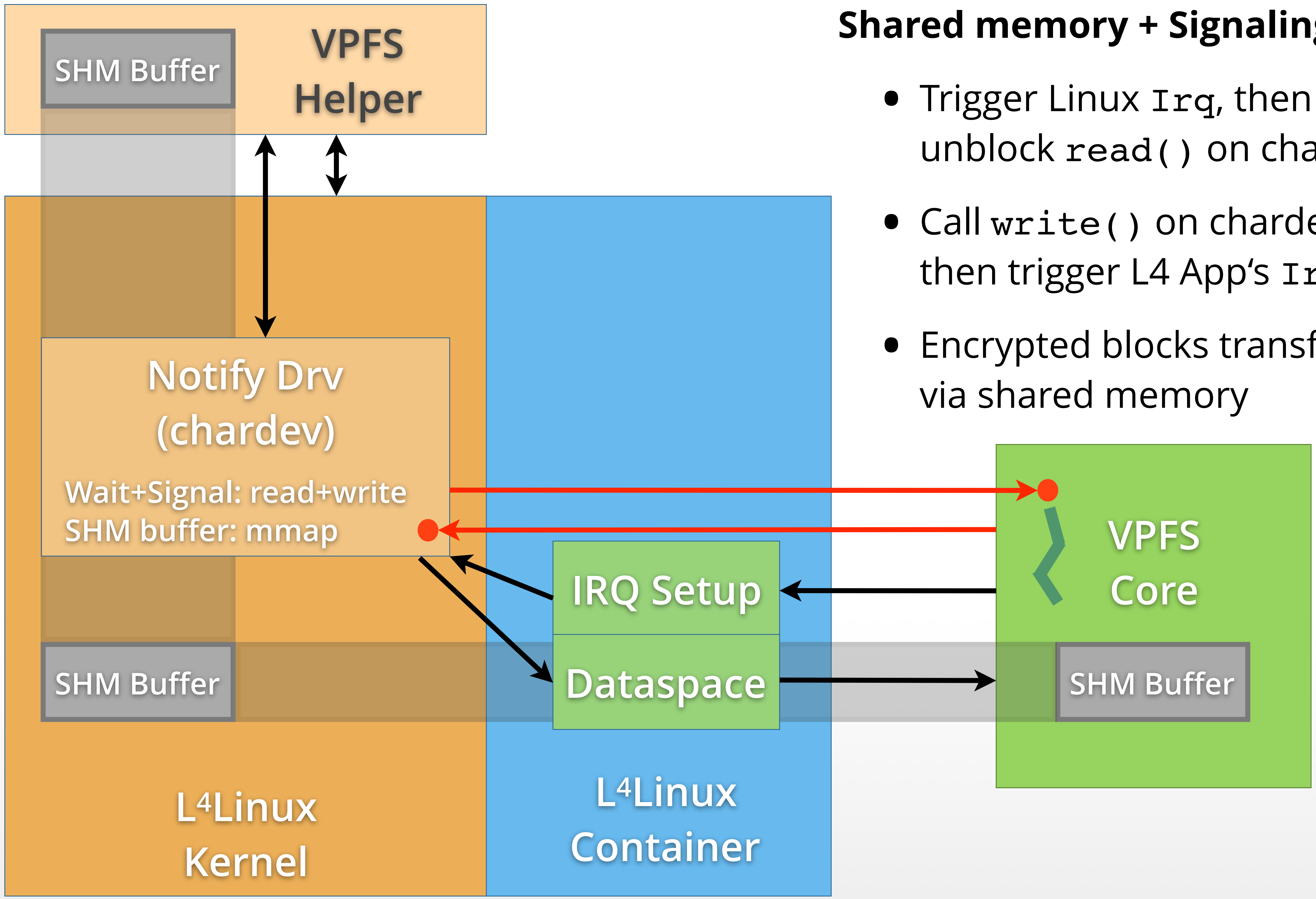
MULTIPLE FILES



EXTENSIVE REUSE

- VPFS reuses Linux file system stack:
  - Drivers, block device layer
  - Optimizations (buffer cache, read ahead, write batching, ...)
  - Allocate / free disk storage for files
- Cooperation: proxy driver in L4Linux

VPFS PROXY DRIVER



Shared memory + Signaling:

- Trigger Linux `Irq`, then unblock `read()` on chardev
- Call `write()` on chardev, then trigger L4 App's `Irq`
- Encrypted blocks transferred via shared memory

VPFS SUMMARY

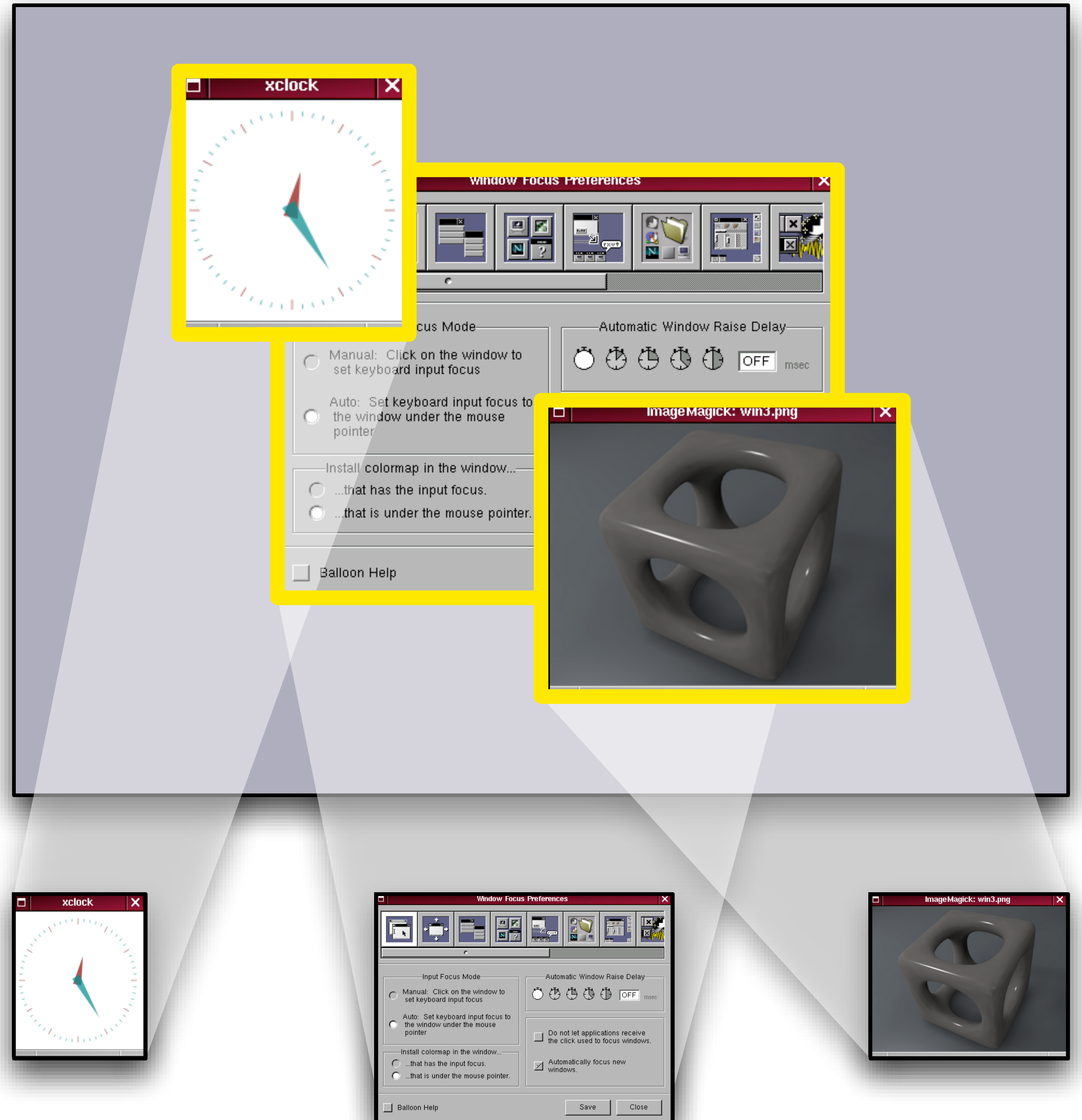
- Trusted wrappers for file systems work!
- VPFS is general purpose file system
- Significant reduction in code size:
  - Untrusted Linux file system stack comprises **50,000+** SLOC
  - VPFS adds **4,000** to **4,600** SLOC to application TCB [3]
  - jVPFS adds another **350** SLOC for secure journaling to protect against crashes [4]

USER INTERFACES

SCREEN SHARING

- Isolated applications run in different domains of trust, but separate screens are inconvenient
- The Nitpicker solution [5]:
  - Let all windows share the same screen ...
  - ... but securely:
    - Make windows & applications identifiable
    - Prevent them from spying on each other: route input securely, no screenshots

CONCEPTS



Views

Buffers

Image source: [5]

HOW IT WORKS

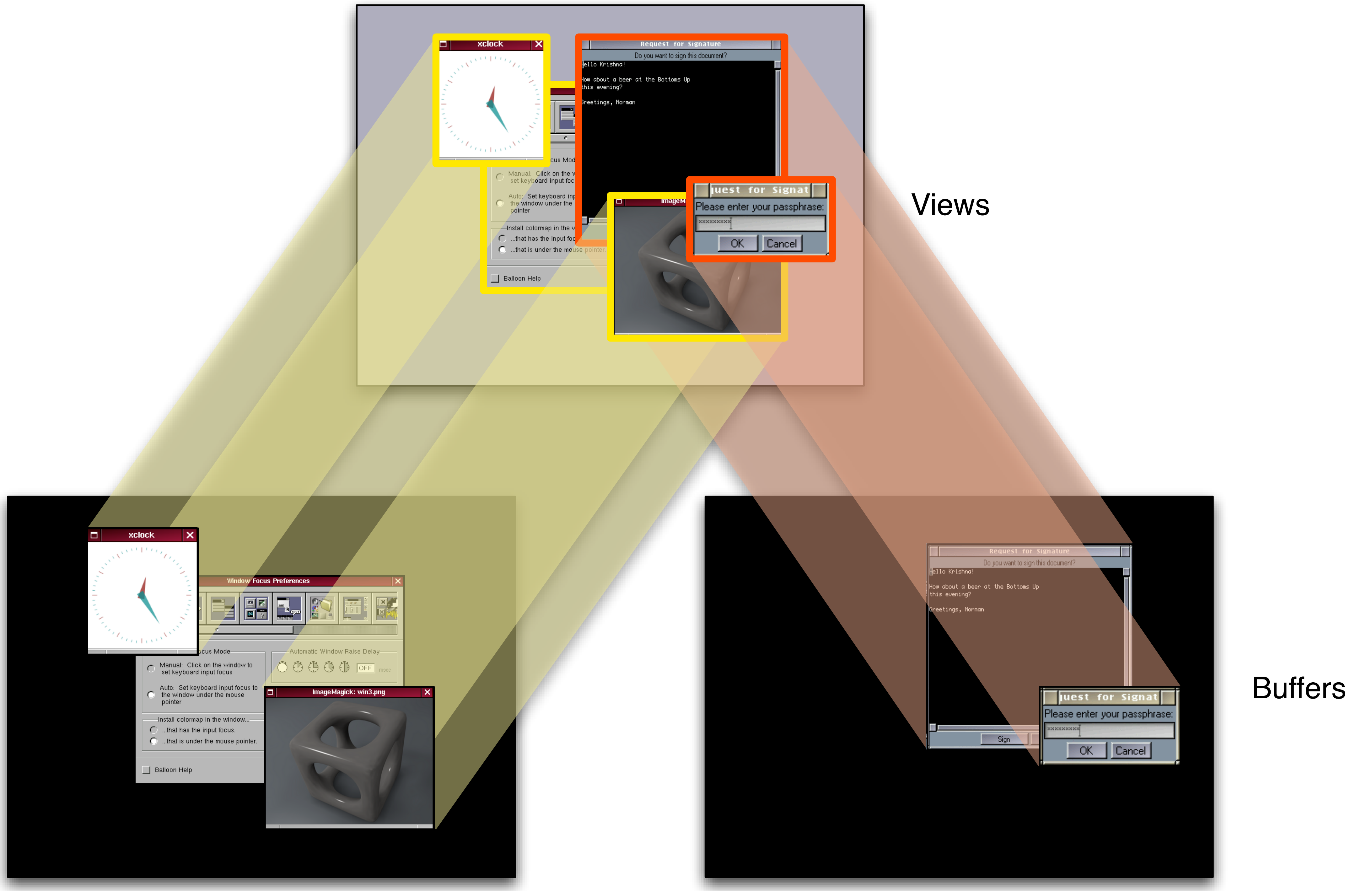


Image source: [5]

DEMO

NITPICKER IN ACTION

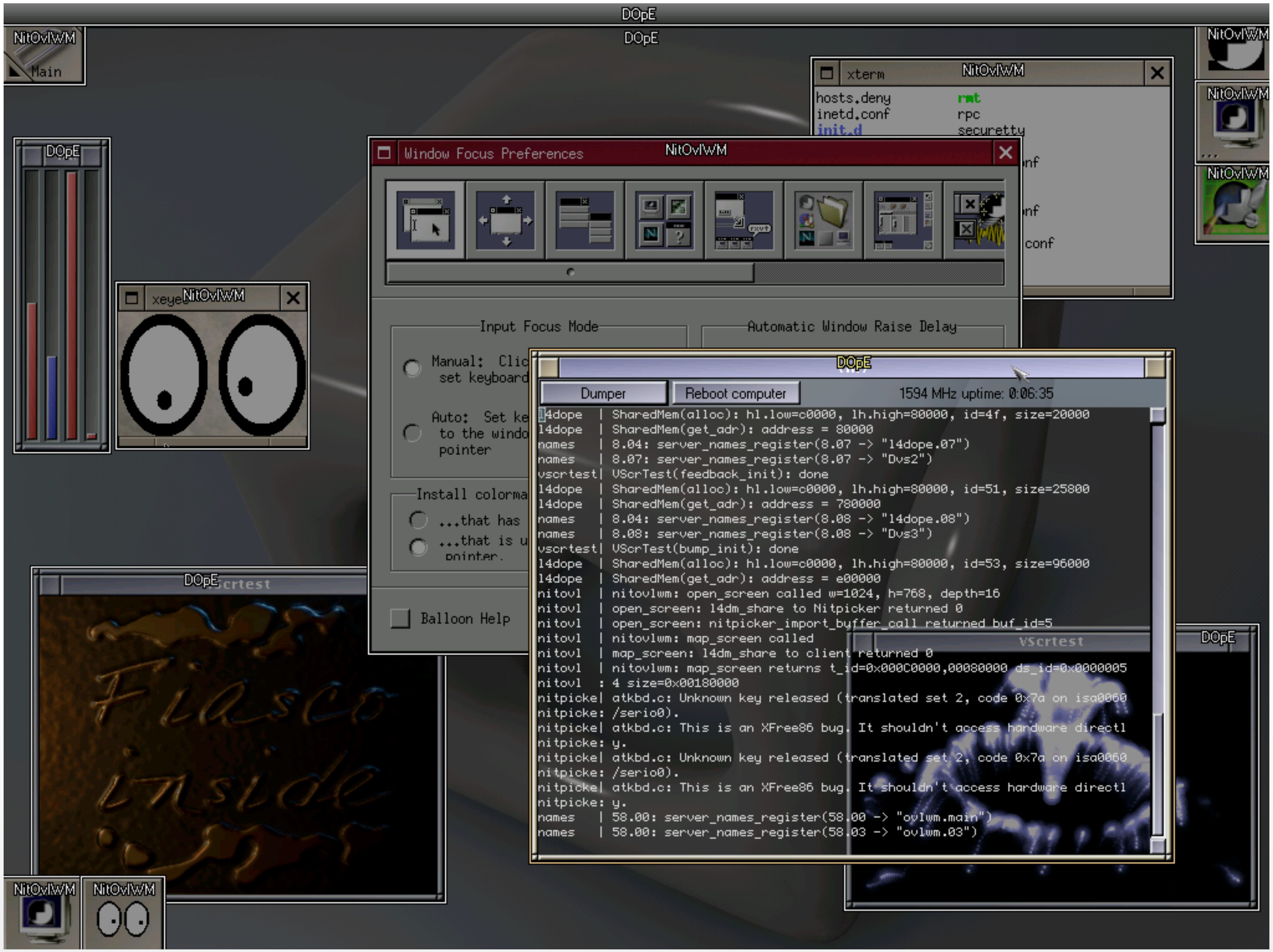


Image source: [5]

SUMMARY

- Secure reuse of untrusted legacy infrastructure
- Split apps + OS services for smaller TCB
- Nizza secure system architecture:
  - Strong isolation
  - Application-specific TCBs
  - Legacy Reuse
  - Trusted Wrapper

COMING UP

- Next week, January 14:
  - Lecture on “Trusted Computing”
    - Where does VPFS store its secrets?
    - How to prevent tampering with stored data?
    - How to trust in what Nitpicker shows on screen?
  - *No exercise or complex lab!*

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REFERENCES

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- [4] Carsten Weinhold and Hermann Härtig, „jVPFS: Adding Robustness to a Secure Stacked File System with Untrusted Local Storage Components“, Proceedings of the 2011 USENIX Annual Technical Conference, Portland, OR, USA, June 2011
- [5] Norman Feske and Christian Helmuth, „A Nitpicker's guide to a minimal-complexity secure GUI“, ACSAC '05: Proceedings of the 21st Annual Computer Security Applications Conference, 2005, Washington, DC, USA
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