



**Paper Reading Group:
The Transfer of Information and Authority
in a Protection System – Matt Bishop**

Marcus Völp

- `71 Butler Lampson:
 - Definition of Access Control Matrix
- `76 Harrison, Ruzzo, Ullman:
 - In the most general abstract case, security (leakage) is undecidable.
- `76 Jones, Lipton, Snyder:
 - Specific system in which leakage is decidable in linear time:
 - Take-Grant Protection Model
- **`79 Bishop, Snyder:**
 - De Facto leakage in the take grant model
- `84 Bobert, Karger: Unmodified capability systems cannot enforce * property (no write down)
- `88 Karger: Unmodified capability system cannot enforce confinement

Idea: Implement Turing Machine with general ACM; Reduction to Halting Problem

Turing Machine:

K – States (p, q, ...)

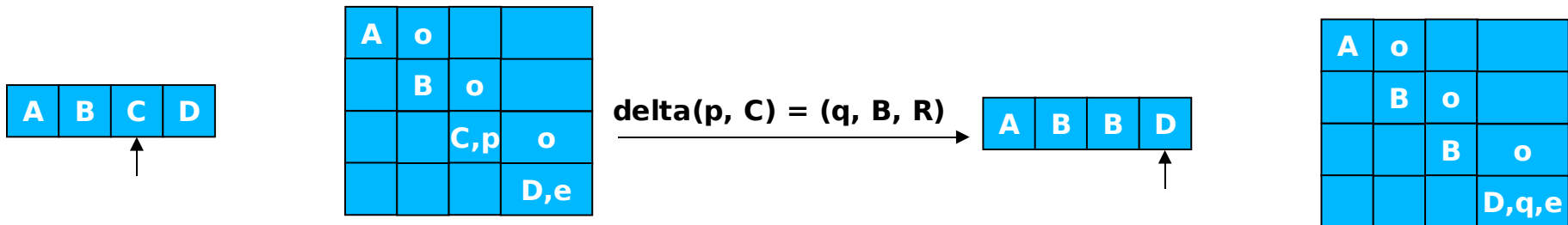
M – Type Symbols (A, B, ...)

delta – Transition Function: $K \times M \rightarrow K \times M \times \{L,R\}$

Subjects s_i = Cells of Tape

Access Rights:

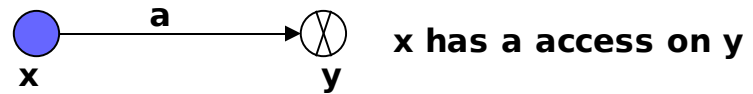
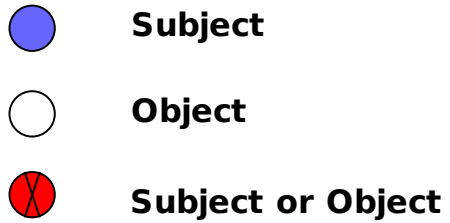
- Type Symbols as Access Rights in $s_i \times s_i$
- own in $s_{i+1} \times s_i$
- States as Access Rights in $s_i \times s_i$ iff head is at cell i



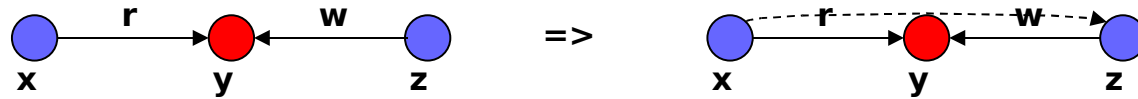
if own in $a[s_i, s_{i+1}]$ and p in $a[s_i, s_i]$ and C in $a[s_i, s_i]$
delete p from $a[s_i, s_i]$;
delete C from $a[s_i, s_i]$;
enter B into $a[s_i, s_i]$;
enter q into $a[s_{i+1}, s_{i+1}]$

- De Jure:
 - Obtain permissions to read an object
- De Facto:
 - Effectively read the content of an object
 - De jure => de facto
 - De jure gives right to obtain up-to-date information
 - De facto relies on transmitting agents
(de jure in general also but only to obtain the right)

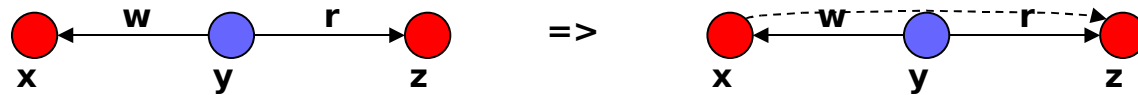
De facto rules



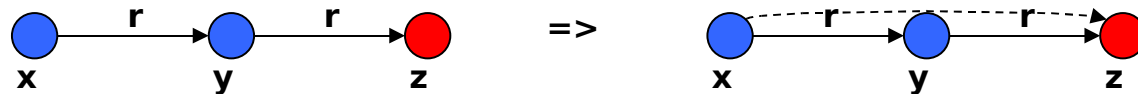
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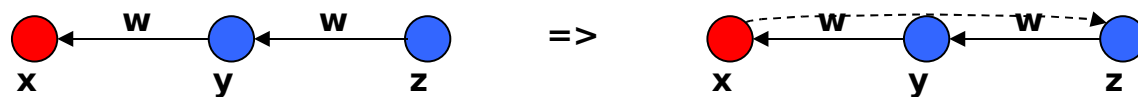
Pass



Spy

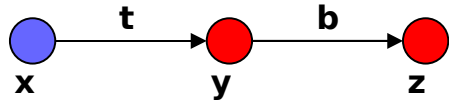


Find

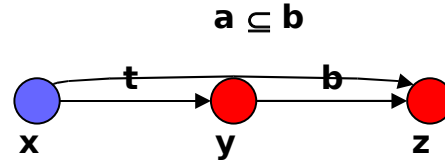


De jure rules

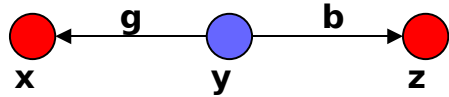
Take



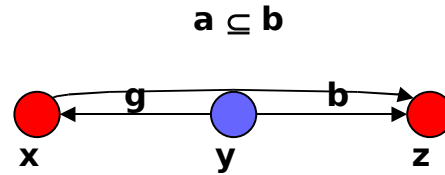
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Grant



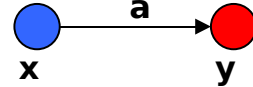
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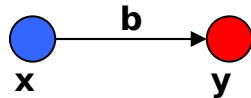
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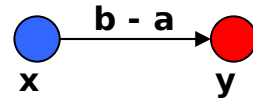
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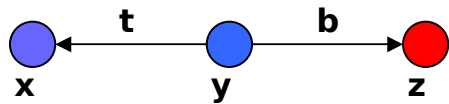
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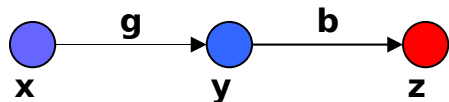
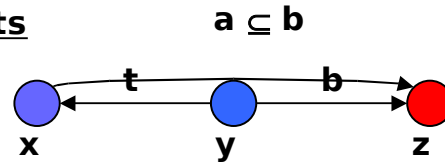
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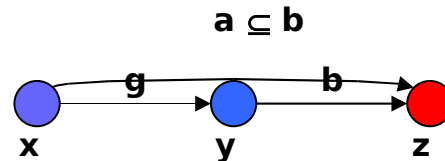
Sharing of Rights



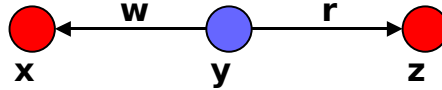
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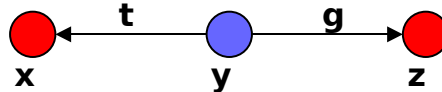
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- **RW-Path** – Sequence of vertices $v_0 \dots v_k$; $v_i \rightarrow v_{i+1}$; r or w in $\text{Label}(e_i)$



- **TG-Path:**



- De Facto Predicate:

- p **can know** q in $G_0 \Leftrightarrow$

- Exists Sequence of Graphs $G_1 \dots G_n$ with
DF
 $G_i \vdash G_{i+1}$

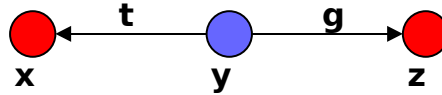
- De Facto Edge $p \rightarrow q$ in G_n

- **admissible RW Path**

- associated word $(r \vec{u} w)^*$ and
- $a_i = \vec{r} \Rightarrow v_{i-1}$ is subject; $a_i = w$ then v_i is subject

Can know \Leftrightarrow Exists admissible RW path

- **TG-Path:**



- **Island**

- Maximal tg-connected; subject only subgraph

- **Spans / Bridges**

- v0 subject

- initial span: $\{\overrightarrow{t^*} \overrightarrow{g}\} \cup \{e\}$

- terminal span: $\{\overleftarrow{t^*}\}$

- bridge: $\{\overrightarrow{t^*}, \overleftarrow{t^*}, \overrightarrow{t^*} \overrightarrow{g} \overleftarrow{t^*}, \overleftarrow{t^*} \overleftarrow{g} \overrightarrow{t^*}\}$ and vk is subject

- De Jure Theorem:

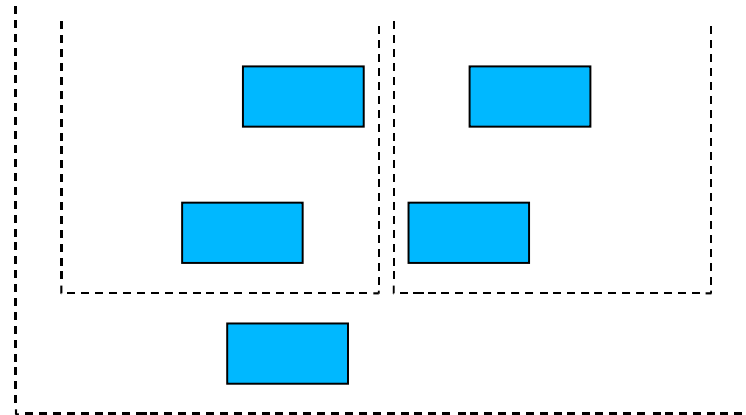
- p **can share** a with q in $G_0 \Leftrightarrow$

- Exists s in G_0 with s-to-q edge labeled a
- Exist subject vertices p' and s' such that
 - p' initially spans to p
 - s' terminally spans to s
- Exists islands $I_1, \dots, I_{v'}$
 - p in I_1 , s' in $I_{v'}$
 - Bridge from I_j to I_{j+1}

- **can share** relies on witness subjects in the Islands
- Can x obtain a even if some subjects don't cooperate?
- x **can steal** a on y in G_0
 - Property:
 - No edge x to y labeled a
 - Sequence G_1, \dots, G_n ; $x \xrightarrow{a} y$ in G_n
 - $G_i \stackrel{\pi_i}{\sim} G_{i+1}$
 - For all vertices v, w in G_{i-1}

Exists v -to- y in G_0 labeled $a \Rightarrow$
 π_i is not: v grants (a on y) to w
 - Theorem:
 - there is no edge from x to y labeled a in G_0
 - Exists subject x' : $x = x'$ or x' initially spans to x
 - Exists vertex s with edge labeled a to y in G_0 and for which **can share** (t, x, s, G_0) holds.

- * property – no writes to lower objects
- Confinement
 - Shapiro: Confined process should not be able to affect any non-authorized entities outside confinement boundary.



- No take outside confinement boundary
(EROS: weak attribute allows take of read only, caps)
- Parent does not grant into confinement

- **I like this stuff:**
 - !!! Weird Math by Drawing Pictures !!!**
- Things to Discuss / Open Issues from my point of view
 - How do Confinement, de Facto Knowledge and Noninterference interact?
 - More on trusted servers
 - How to formulate proper object reuse:
Grant access to C2 after access is being revoked from C1.
 - Security / Resource Policies based on Identity
 - Do we need them?
 - Can we build a pure capability system?
 - No quotas but caps on resources
 - No names
 - Everything is authorized through capabilities