



Analyzing Integrity Protection in the SELinux Example Policy

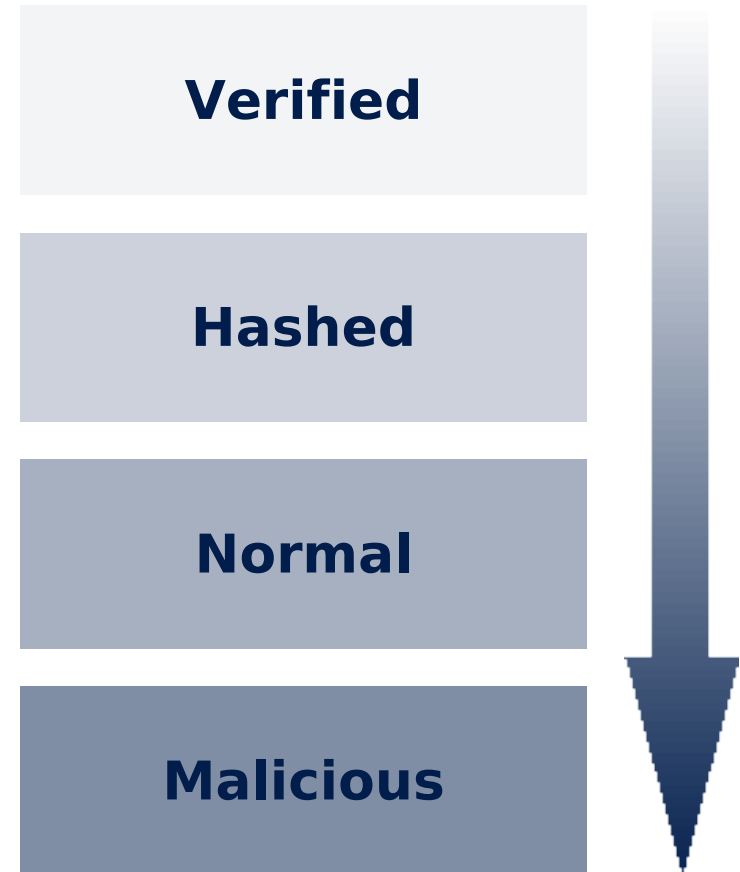
Trent Jaeger et al. (presented by S. Kalkowski)

Dresden, 2008-01-23



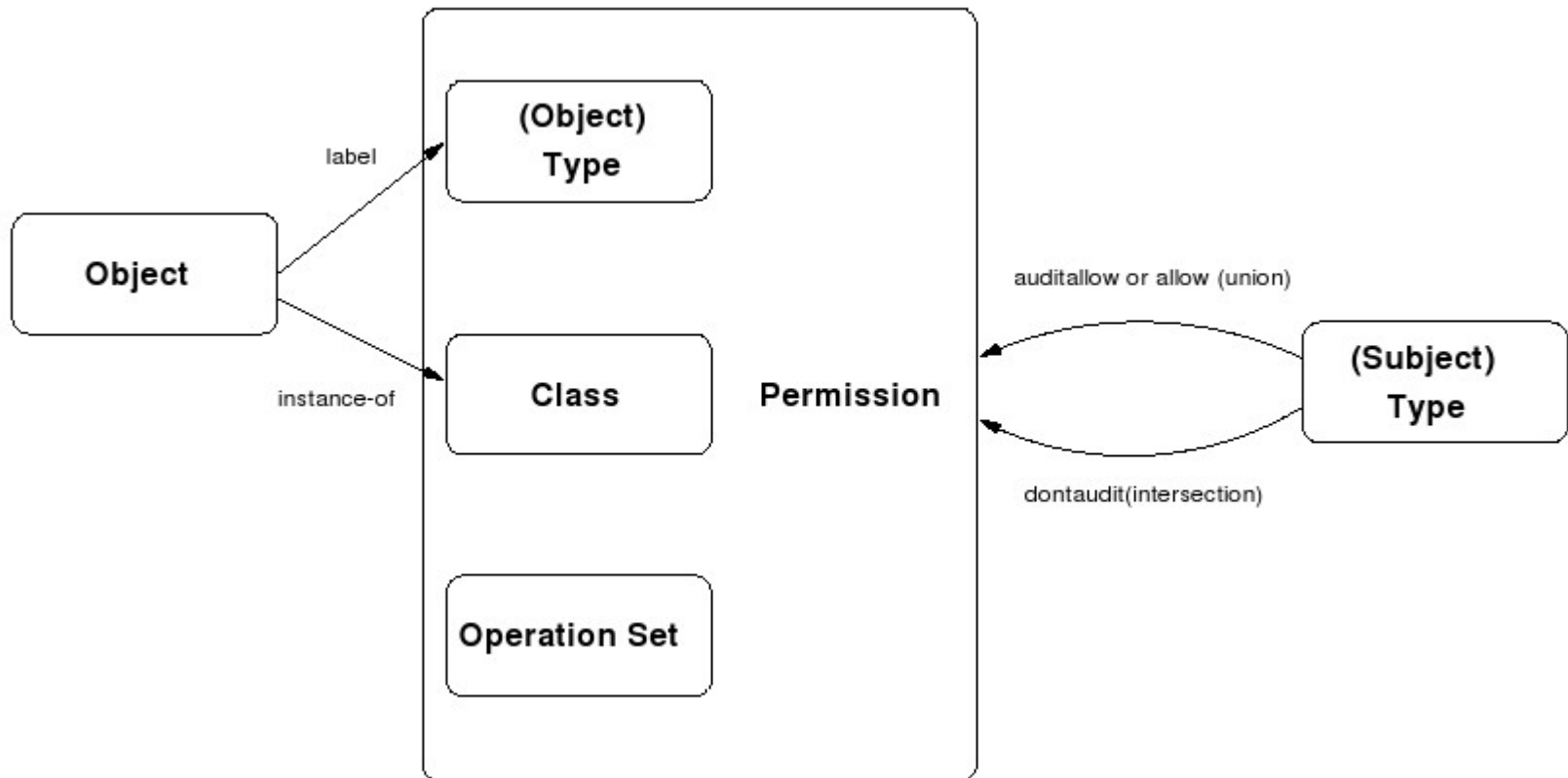
- SELinux example policy doesn't define protection goals, but demands of the different applications
- Define integrity protection goals for the overall TCB
- Handle conflicts with the example policy
- Provide a tool that eases up that work to system administrators and developers

- Easy model
- Goal:
 - higher integrity layers have not to be polluted/dependent by/on lower ones
- Only two rules:
 - no write up
 - no read down
- Very strict and impractical



- Type Enforcement
- Role Based Access Control
- No strict hierarchy
- **Problem:**
 - No explicit security goals anymore
 - The more permissions statement the more tricky is a proof of certain security properties

- Extended **T**ype **E**nforcement model





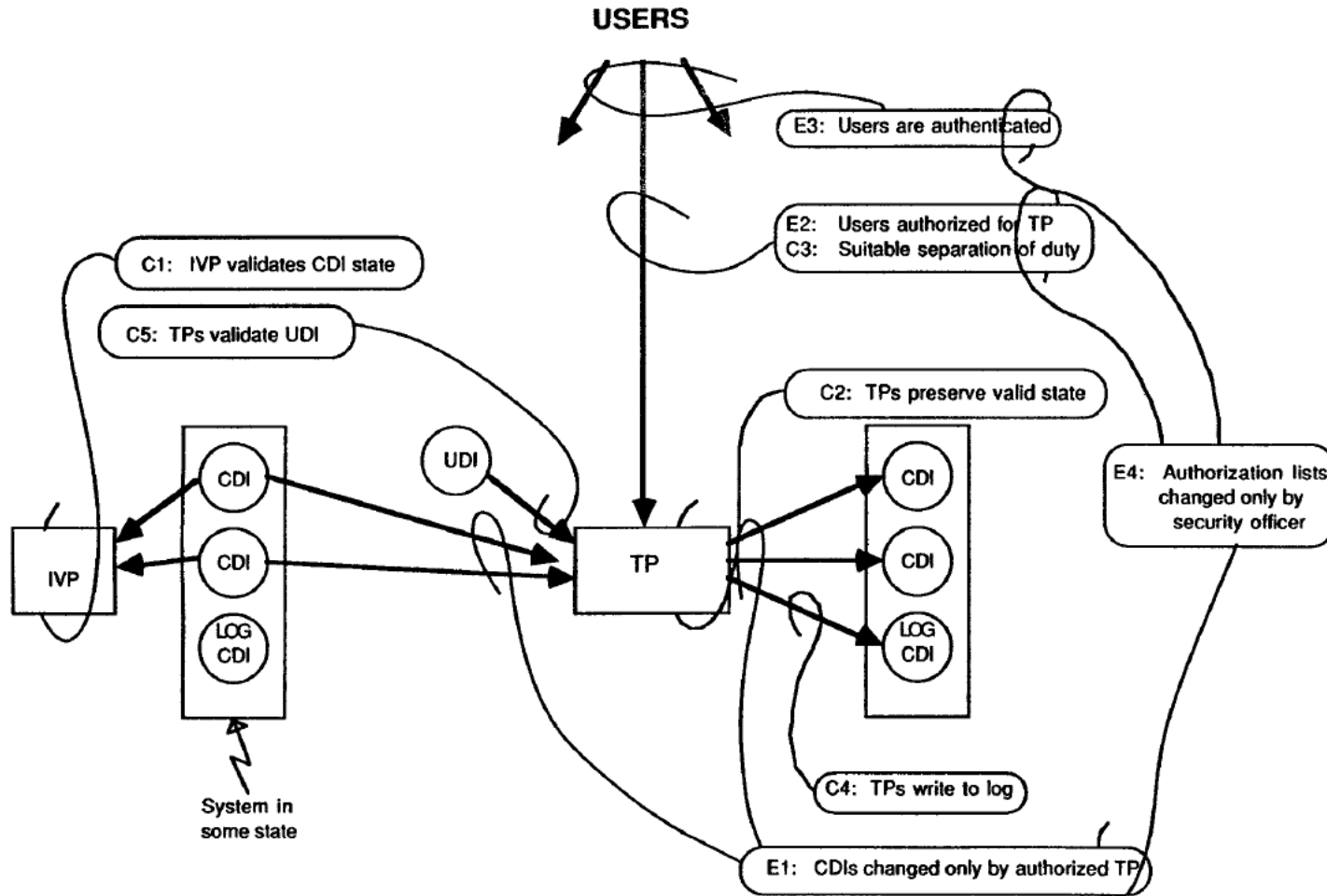
- Beside the type declarations and allow statements, there are:
 - Attribute declarations
 - Type transition statements
 - execution of certain programs
 - creation of files by certain processes
 - Role to types transition statements
 - Role transition statements when executing a new program (e.g.: login)
 - User to roles mappings
- **Complex model!**

```
attribute ssh_server;
type ssh_exec_t;
corecmd_executable_file(ssh_exec_t)
type ssh_keygen_t;
type ssh_keygen_exec_t;
init_system_domain(ssh_keygen_t,ssh_keygen_exec_t)
role system_r types ssh_keygen_t;
type ssh_keysign_exec_t;
corecmd_executable_file(ssh_keysign_exec_t)
type sshd_exec_t;
corecmd_executable_file(sshd_exec_t)
type sshd_key_t;
files_type(sshd_key_t)
ifdef(`targeted_policy',`
unconfined_alias_domain(sshd_t)
init_system_domain(sshd_t,sshd_exec_t)
type sshd_var_run_t;
files_type(sshd_var_run_t)
ifdef(`enable_mcs',`
init_ranged_system_domain(sshd_t,sshd_exec_t,s0 - mcs_systemhigh)
')
`,`
type ssh_agent_exec_t;
files_type(ssh_agent_exec_t)
ssh_server_template(sshd)
ssh_server_template(sshd_extern)
init_daemon_domain(sshd_t,sshd_exec_t)
ifdef(`enable_mcs',`
init_ranged_daemon_domain(sshd_t,sshd_exec_t,s0 - mcs_systemhigh)
')
```

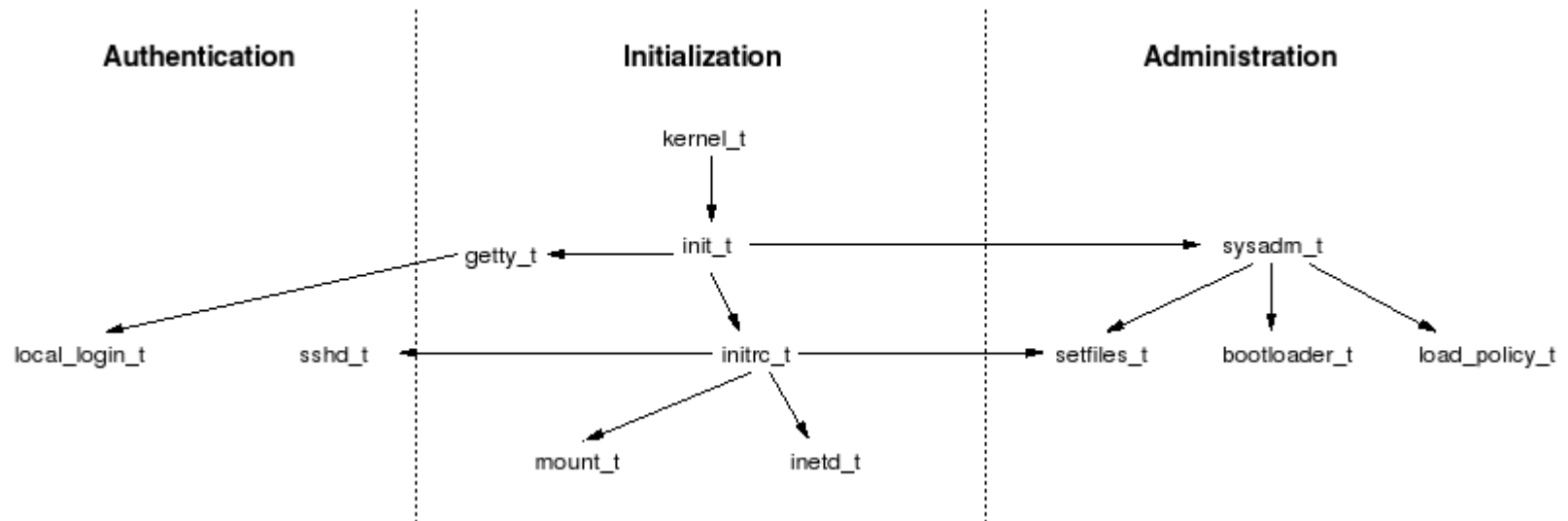
- SELinux policy too complex to simply verify site goals, e.g.: of an organization, user, application developer
- Once some protection goals are verified, we cannot change the policy without verifying it again
- How to ensure integrity of the common TCB

- Identify the TCB common to all components and protect its integrity
- Use Clark-Wilson Model:
 - Defines mainly two levels of integrity UDI and CDI
 - Defines how UDIs can be transformed to CDIs (by TPs) and how their integrity has to be verified (by IVPs)
 - Defines policy enforcement and certification rules
 - Introduced concept of **separation of duty** and **well-formed transaction**

Figure 1: Summary of System Integrity Rules



- Take initial TCB subject set from the SELinux policy (high integrity)
- Analyze read permissions to objects
- Find all non-TCB subjects (low integrity) that write to these object -> **integrity violation**



- Using Gokyo – policy analysis tool:
 - Identify access control spaces of subjects
 - Set of all possible permission assignments
 - Set of prohibitions
 - Set of unknown rights
 - Overlapping regions
 - Define handling routines for whole subspaces
 - Granting or denying subspace
 - Manually changing the policy



- Does the TCB subject only reads or also writes to the concerning object type?
 - read only data might be sanitized
- How many conflicts are caused by a non-TCB subject?
 - > 1 possibly a TCB candidate
- Can we exclude the conflicting subject ?
 - automatically detected by attributes
 - e.g.: httpd_*
 - also include related object types

- Classifications for integrity conflicts

<i>Class</i>	<i>Description</i>
TCB or Candidate	Trusted subject types
Exclude Type	Type can be excluded from secure system with this TCB
Sanitize	A sanitized read may be used to protect TCB
Denial	Denial of conflicting rights can be used to protect TCB
Modify Policy	Policy must be edited to protect TCB

<i>Trusted Type</i>	<i>Conflict Type</i>	<i>Object Type & Op</i>	<i>Class</i>	<i>Resolution</i>
dpkg_t	tmpreaper_t	tmp_dpkg_t:file rw	exclude	exclude
initrc_t	many	file_type:blk/chr/file r	sanitize	sanitize
initrc_t	useradd_t	etc_t:file r	trust	trust
initrc_t	hwclock_t	clock_device_t:chr/blk rw	trust	trust
initrc_t	gpm_t	psaux_t:chr rw	exclude	exclude
initrc_t	sound_t, xdm_t	sound_device_t:chr rw	trust	exclude
initrc_t	httpd_admin_xserver_t	framebuf_device_t:chr rw	deny	exclude
initrc_t	many	initrc_t:fifo rw	deny	sanitize
kernel_t	slapd_t, squid_t, +	*:*_socket r	sanitize	sanitize
kernel_t	dhcpc_t	resolv_conf_t:file r	trust	exclude
kernel_t	dhcpcd_t	var_run_t:file r	trust	exclude
kernel_t	quota_t	file_t:file r	trust	trust
local_login_t	many	proc_t:file r	sanitize	sanitize
local_login_t	insmod_t	local_devpts_t:process r	deny	change
local_login_t	logrotate_t	local_login_t:process r	trust	trust
mount_t	automount_t	autofs_t:dir rw	exclude	trust
mount_t	bootloader_t, fsadm_t	fixed_disk_device_t:* rw	trust	trust
sysadm_t	user_t	misc_device_t:* rw	deny	exclude obj
sysadm_t	many	sysadm_devpts_t/ptyfile:* rw	deny	change
sysadm_t	sysadm_*_t	sysadm_home_t:* rw	deny	change/sanitize one file
sysadm_t	sysadm_*_t	sysadm_tmp_t:file rw	exclude	change
sysadm_t	sysadm_irc_t	sysadm_irc_t:file rw	exclude	change/sanitize
sysadm_t	sysadm_xserver_t	sysadm_xserver_t:shm rw	exclude	exclude
sysadm_t	sysadm_xauth_t	sysadm_home_xauth_t:file rw	exclude	exclude
sysadm_t	admin	kernel_t:system_ave_toggle rw	trust	trust
sshd_t	many	sshd_devpts_t/userpty:* rw	deny	change

**46 %
“wrong classifications”**

- In general, are such techniques (for **automatically** resolving policy conflicts) needed or can we solve the problem of policy complexity (e.g. : by divide & conquer) ?
- What totally different approaches for handling policy complexity can we imagine -> e.g.: (transparent) user-interactive decisions
- Is policy complexity an artificial problem, as we have to implement least privilege beforehand