

Jettison: Efficient Idle Desktop Consolidation with Partial VM Migration

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Paper Reading Group

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Outline

- 1 Introduction
- 2 Partial VM Migration
- 3 Implementation of Jettison
- 4 Evaluation
- 5 Scalability
- 6 Discussion

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Motivation

- Office computers are left running, even when idle because of applications that require always-on semantics
- Power consumption in idle mode is typically 60% of a fully utilized system
- Straightforward solution requires large network transfers, much memory on the server and long reintegrate times
- Only a small fraction of the memory is needed for running on the server (10% of memory and $< 1\%$ of disk state)

Related work

- Focused on coarse grained migration, i.e. migrating the VM completely
- Live migration: costs significant amount of network traffic, memory and time
- Ballooning: takes considerable amount of time and I/O, although its cheap to finally migrate the shrunken VM
- Remote desktop access: limited, because it doesn't allow seamless access to local devices and has bad performance

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How does it work?

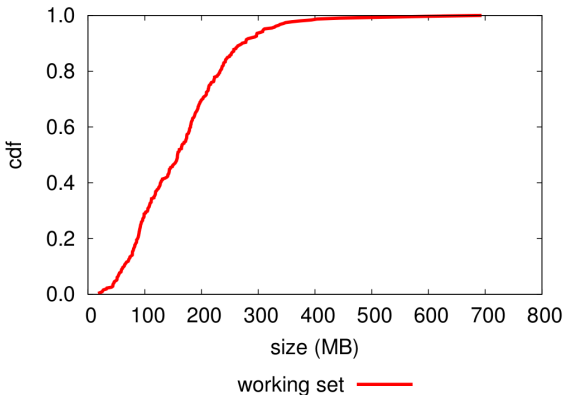
- During consolidation, partial VM migration transfers only a VM descriptor
- As the VM tries to access it's pages, it causes *remote faults* and pages are loaded from the desktop
- The disk state is also fetched on-demand from the desktop (if necessary)
- Between bursts of remote faults, the desktop can *microsleep*
- Reintegration transfers only the dirty state back to the desktop

Find answers to ...

- 1 When to microsleep?
- 2 How can prefetching improve microsleeps?

State access traces

Deployment with 3 users over 7 weeks, using Linux VMs with 4 GiB of memory:



When to microsleep?

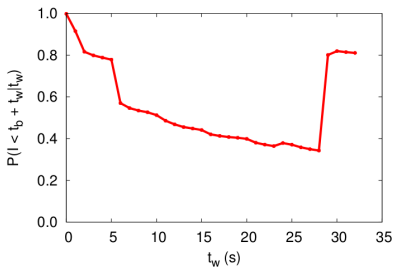
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When to microsleep?

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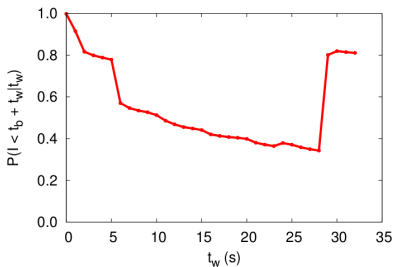
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- 4 Determine optimal wait time by minimizing wasted energy (6s)

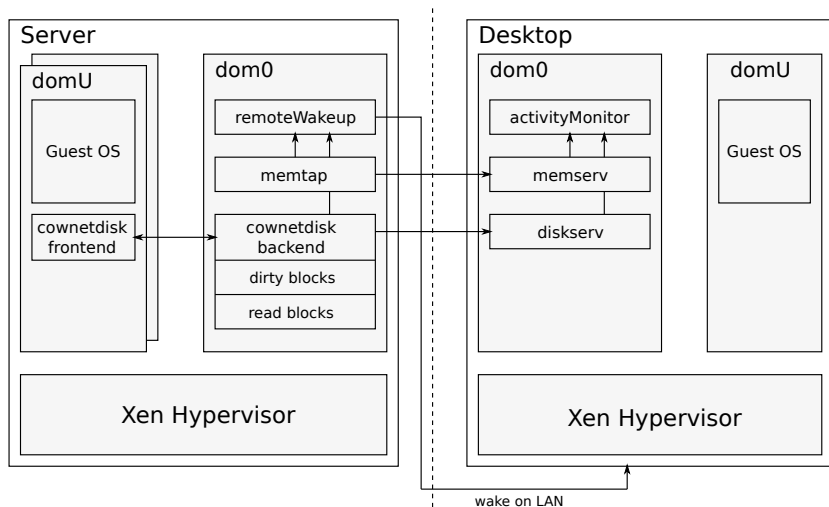
How can prefetching improve microsleeps?

- Prefetching is used to increase the frequency and length of microsleeps
- They tried 2 strategies:
 - ① *hoarding*: at consolidation, transfer a few pages that have been requested in previous migrations of the same VM.
 - ② *on-demand prefetch*: prefetch pages that are near the requested page
- Result: on-demand prefetch reaches longer microsleeps, which results in better energy savings

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Architecture



Consolidation and reintegration

When to consolidate?

- 1 User idles
- 2 Server has enough capacity
- 3 VM can execute sufficiently autonomous on the server

When to reintegrate?

- 1 User becomes active
- 2 Server capacity is exceeded
- 3 VM becomes active (requires a large amount of state from desktop)

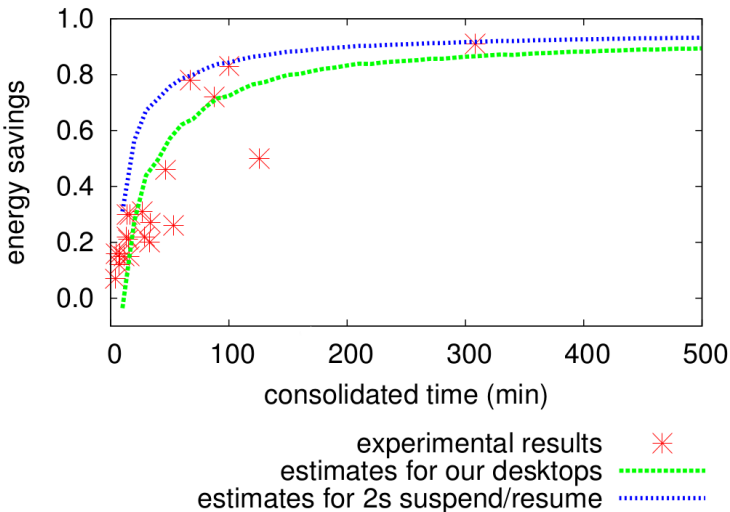
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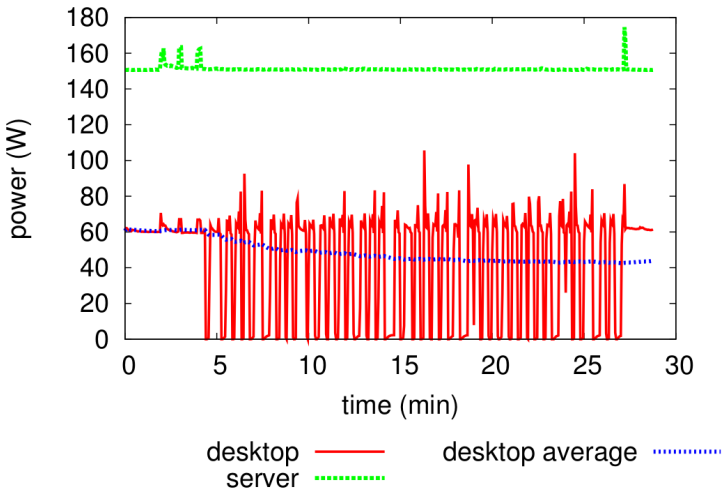
Setup

- Deployment of the prototype with 4 users over 6 days
- Desktops used standard Linux systems with 4 GiB of memory and 12 GiB of disk
- Server had 16 GiB of memory
- Connection was a 1 GiB/s ethernet switch

Energy savings



Power consumption

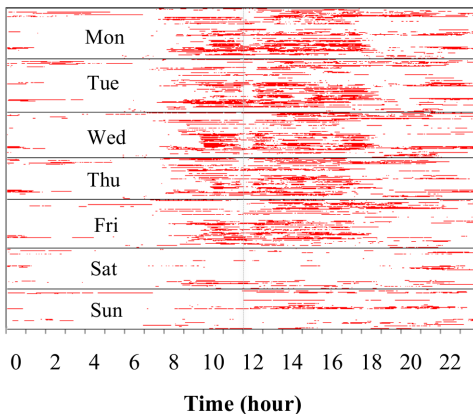


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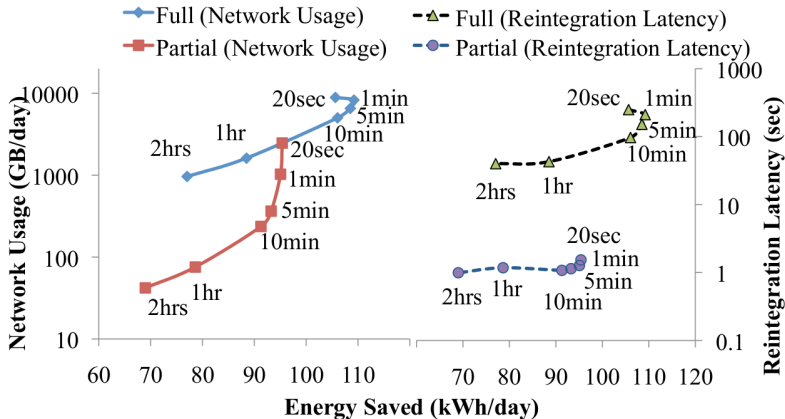
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Approach

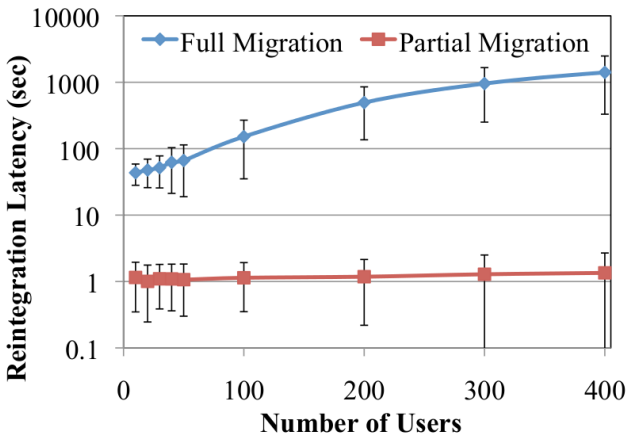
- Collect idleness traces from real users in an office environment
- Idleness was tracked for 4 months at an industrial research lab with 22 researchers



Network usage and reintegration latency



Reintegration latency



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Ideas & Questions

- Do the desktop need to be virtualized?
- The whole concept is based on the assumption that systems in idle mode consume nearly as much energy as under full load. Will that be true for future systems?
- Maybe one could use the same concept to allow people to “take their system home”?